

# **PIC16C84**

### 8-bit CMOS EEPROM Microcontroller

#### High Performance RISC CPU Features:

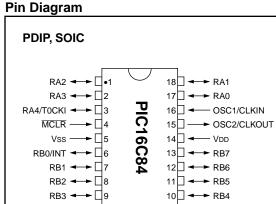
- · Only 35 single word instructions to learn
- All instructions single cycle (400 ns @ 10 MHz) except for program branches which are two-cycle
- Operating speed: DC 10 MHz clock input DC - 400 ns instruction cycle
- 14-bit wide instructions
- 8-bit wide data path
- 1K x 14 EEPROM program memory
- 36 x 8 general purpose registers (SRAM)
- 64 x 8 on-chip EEPROM data memory
- 15 special function hardware registers
- Eight-level deep hardware stack
- Direct, indirect and relative addressing modes
- Four interrupt sources:
- External RB0/INT pin
- TMR0 timer overflow
- PORTB<7:4> interrupt on change
- Data EEPROM write complete
- 1,000,000 data memory EEPROM ERASE/WRITE cycles
- EEPROM Data Retention > 40 years

#### **Peripheral Features:**

- 13 I/O pins with individual direction control
- High current sink/source for direct LED drive
  - 25 mA sink max. per pin
  - 20 mA source max. per pin
- TMR0: 8-bit timer/counter with 8-bit programmable prescaler

#### **Special Microcontroller Features:**

- Power-on Reset (POR)
- Power-up Timer (PWRT)
- Oscillator Start-up Timer (OST)
- Watchdog Timer (WDT) with its own on-chip RC oscillator for reliable operation
- · Code protection
- Power saving SLEEP mode
- · Selectable oscillator options
- · Serial In-System Programming via two pins



#### **CMOS Technology:**

- Low-power, high-speed CMOS EEPROM technology
- · Fully static design
- Wide operating voltage range:
  - Commercial: 2.0V to 6.0V
  - Industrial: 2.0V to 6.0V
- Low power consumption:
  - < 2 mA typical @ 5V, 4 MHz</li>
  - 60  $\mu\text{A}$  typical @ 2V, 32 kHz
  - 26  $\mu A$  typical standby current @ 2V

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#### 1.0 GENERAL DESCRIPTION

The PIC16C84 is a low-cost, high-performance, CMOS, fully-static, 8-bit microcontroller.

All PIC16/17 microcontrollers employ an advanced RISC architecture. PIC16CXX devices have enhanced core features, eight-level deep stack, and multiple internal and external interrupt sources. The separate instruction and data buses of the Harvard architecture allow a 14-bit wide instruction word with a separate 8-bit wide data bus. The two stage instruction pipeline allows all instructions to execute in a single cycle, except for program branches (which require two cycles). A total of 35 instructions (reduced instruction set) are available. Additionally, a large register set is used to achieve a very high performance level.

PIC16CXX microcontrollers typically achieve a 2:1 code compression and up to a 2:1 speed improvement (at 10 MHz) over other 8-bit microcontrollers in their class.

The PIC16C84 has 36 bytes of RAM, 64 bytes of Data EEPROM memory, and 13 I/O pins. A timer/counter is also available.

The PIC16CXX family has special features to reduce external components, thus reducing cost, enhancing system reliability and reducing power consumption. There are four oscillator options, of which the single pin RC oscillator provides a low-cost solution, the LP oscillator minimizes power consumption, XT is a standard crystal, and the HS is for High Speed crystals. The SLEEP (power-down) mode offers power savings. The user can wake the chip from sleep through several external and internal interrupts and resets.

A highly reliable Watchdog Timer with its own on-chip RC oscillator provides protection against software lockup.

The PIC16C84 EEPROM program memory allows the same device package to be used for prototyping and production. In-circuit reprogrammability allows the code to be updated without the device being removed from the end application. This is useful in the development of many applications where the device may not be easily accessible, but the prototypes may require code updates. This is also useful for remote applications where the code may need to be updated (such as rate information).

Table 1-1 lists the features of the PIC16C84. A simplified block diagram of the PIC16C84 is shown in Figure 3-1.

The PIC16C84 fits perfectly in applications ranging from high speed automotive and appliance motor control to low-power remote sensors, electronic locks, security devices and smart cards. The EEPROM technology makes customization of application programs (transmitter codes, motor speeds, receiver frequencies, security codes, etc.) extremely fast and convenient. The small footprint packages make this microcontroller series perfect for all applications with space limitations. Low cost, low power, high performance, ease of use and I/O flexibility make the PIC16C84 very versatile even in areas where no microcontroller use has been considered before (e.g., timer functions, serial communication, capture and compare, PWM functions and co-processor applications).

The serial in-system programming feature (via two pins) offers flexibility of customizing the product after complete assembly and testing. This feature can be used to serialize a product, store calibration data, or program the device with the current firmware before shipping.

#### 1.1 Family and Upward Compatibility

Those users familiar with the PIC16C5X family of microcontrollers will realize that this is an enhanced version of the PIC16C5X architecture. Please refer to Appendix A for a detailed list of enhancements. Code written for PIC16C5X can be easily ported to the PIC16C84 (Appendix B).

#### 1.2 Development Support

The PIC16CXX family is supported by a full-featured macro assembler, a software simulator, an in-circuit emulator, a low-cost development programmer and a full-featured programmer. A "C" compiler and fuzzy logic support tools are also available.

#### TABLE 1-1 PIC16C8X FAMILY OF DEVICES

		PIC16F83	PIC16CR83	PIC16F84	PIC16CR84
Clock	Maximum Frequency of Operation (MHz)	10	10	10	10
	Flash Program Memory	512	—	1K	—
	EEPROM Program Memory	—	—	—	—
Memory	ROM Program Memory	—	512	—	1K
	Data Memory (bytes)	36	36	68	68
	Data EEPROM (bytes)	64	64	64	64
Peripherals	Timer Module(s)	TMR0	TMR0	TMR0	TMR0
	Interrupt Sources	4	4	4	4
	I/O Pins	13	13	13	13
Features	Voltage Range (Volts)	2.0-6.0	2.0-6.0	2.0-6.0	2.0-6.0
	Packages	18-pin DIP, SOIC	18-pin DIP, SOIC	18-pin DIP, SOIC	18-pin DIP, SOIC

All PICmicro<sup>™</sup> Family devices have Power-on Reset, selectable Watchdog Timer, selectable code protect and high I/O current capability. All PIC16C8X Family devices use serial programming with clock pin RB6 and data pin RB7.

#### 2.0 PIC16C84 DEVICE VARIETIES

A variety of frequency ranges and packaging options are available. Depending on application and production requirements the proper device option can be selected using the information in this section. When placing orders, please use the "PIC16C84 Product Identification System" at the back of this data sheet to specify the correct part number.

There are two device "types" as indicated in the device number.

- 1. **C**, as in PIC16**C**84. These devices have EEPROM program memory and operate over the standard voltage range.
- 2. LC, as in PIC16LC84. These devices have EEPROM program memory and operate over an extended voltage range.

When discussing memory maps and other architectural features, the use of **C** also implies the **LC** versions.

#### 2.1 <u>Electrically Erasable Devices</u>

These devices are offered in the lower cost plastic package, even though the device can be erased and reprogrammed. This allows the same device to be used for prototype development and pilot programs as well as production.

A further advantage of the electrically erasable version is that they can be erased and reprogrammed in-circuit, or by device programmers, such as Microchip's PICSTART<sup>®</sup> Plus or PRO MATE<sup>®</sup> II programmers.

NOTES:

#### 3.0 ARCHITECTURAL OVERVIEW

The high performance of the PIC16CXX family can be attributed to a number of architectural features commonly found in RISC microprocessors. To begin with, the PIC16CXX uses a Harvard architecture. This architecture has the program and data accessed from separate memories. So the device has a program memory bus and a data memory bus. This improves bandwidth over traditional von Neumann architecture where program and data are fetched from the same memory (accesses over the same bus). Separating program and data memory further allows instructions to be sized differently than the 8-bit wide data word. PIC16CXX opcodes are 14-bits wide, enabling single word instructions. The full 14-bit wide program memory bus fetches a 14-bit instruction in a single cycle. A twostage pipeline overlaps fetch and execution of instructions (Example 3-1). Consequently, all instructions execute in a single cycle (400 ns @ 10 MHz) except for program branches.

The PIC16C84 addresses 1K x 14 program memory. All program memory is internal.

PIC16CXX devices can directly or indirectly address its register files or data memory. All special function registers including the program counter are mapped in the data memory. An orthogonal (symmetrical) instruction set that makes it possible to carry out any operation on any register using any addressing mode. This symmetrical nature and lack of 'special optimal situations' make programming with the PIC16CXX simple yet efficient. In addition, the learning curve is reduced significantly.

The PIC16C84 has 36 x 8 SRAM and 64 x 8 EEPROM data memory.

PIC16CXX devices contain an 8-bit ALU and working register. The ALU is a general purpose arithmetic unit. It performs arithmetic and Boolean functions between data in the working register and any register file.

The ALU is 8-bits wide and capable of addition, subtraction, shift and logical operations. Unless otherwise mentioned, arithmetic operations are two's complement in nature. In two-operand instructions, typically one operand is the working register (W register), and the other operand is a file register or an immediate constant. In single operand instructions, the operand is either the W register or a file register.

The W register is an 8-bit working register used for ALU operations. It is not an addressable register.

Depending on the instruction executed, the ALU may affect the values of the Carry (C), Digit Carry (DC), and Zero (Z) bits in the STATUS register. The C and DC bits operate as a borrow and digit borrow out bit, respectively, in subtraction. See the SUBLW and SUBWF instructions for examples.

A simplified block diagram for the PIC16C84 is shown in Figure 3-1, its corresponding pin description is shown in Table 3-1.

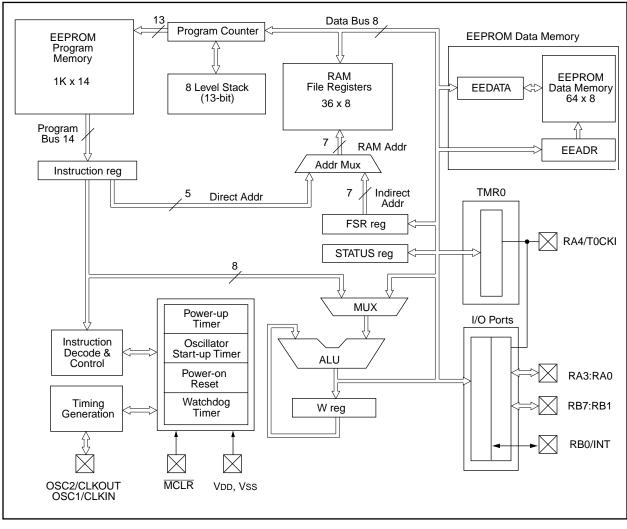


FIGURE 3-1: PIC16C84 BLOCK DIAGRAM

Pin Name	DIP No.	SOIC No.	I/O/P Type	Buffer Type	Description					
OSC1/CLKIN	16	16	I	ST/CMOS <sup>(1)</sup>	Oscillator crystal input/external clock source input.					
OSC2/CLKOUT	15	15	0	_	Oscillator crystal output. Connects to crystal or resonator in crys- tal oscillator mode. In RC mode, OSC2 pin outputs CLKOUT which has 1/4 the frequency of OSC1, and denotes the instruction cycle rate.					
MCLR	4	4	I/P	ST	Master clear (reset) input/programming voltage input. This pin is an active low reset to the device.					
					PORTA is a bi-directional I/O port.					
RA0	17	17	I/O	TTL						
RA1	18	18	I/O	TTL						
RA2	1	1	I/O	TTL						
RA3	2	2	I/O	TTL						
RA4/T0CKI	3	3	I/O	ST	Can also be selected to be the clock input to the TMR0 timer/ counter. Output is open drain type.					
					PORTB is a bi-directional I/O port. PORTB can be software pro- grammed for internal weak pull-up on all inputs.					
RB0/INT	6	6	I/O	TTL	RB0/INT can also be selected as an external interrupt pin.					
RB1	7	7	I/O	TTL						
RB2	8	8	I/O	TTL						
RB3	9	9	I/O	TTL						
RB4	10	10	I/O	TTL	Interrupt on change pin.					
RB5	11	11	I/O	TTL	Interrupt on change pin.					
RB6	12	12	I/O	TTL/ST (2)	Interrupt on change pin. Serial programming clock.					
RB7	13	13	I/O	TTL/ST (2)	Interrupt on change pin. Serial programming data.					
Vss	5	5	Р	—	Ground reference for logic and I/O pins.					
Vdd	14	14	Р	_	Positive supply for logic and I/O pins.					
Legend: I= input	0 = 0 — = N	utput Not used		/O = Input/Out <sub>l</sub> TTL = TTL inpu	Output P = power					

TABLE 3-1 PIC16C8X	<b>PINOUT DESCRIPTION</b>
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Note 1: This buffer is a Schmitt Trigger input when configured in RC oscillator mode and a CMOS input otherwise.

2: This buffer is a Schmitt Trigger input when used in serial programming mode.

#### 3.1 Clocking Scheme/Instruction Cycle

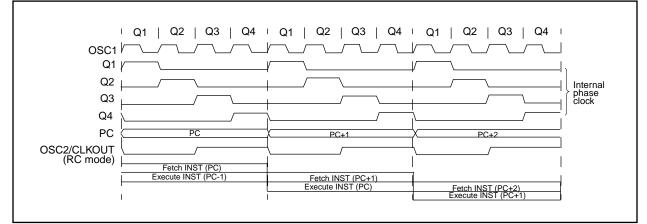
The clock input (from OSC1) is internally divided by four to generate four non-overlapping quadrature clocks namely Q1, Q2, Q3 and Q4. Internally, the program counter (PC) is incremented every Q1, the instruction is fetched from the program memory and latched into the instruction register in Q4. The instruction is decoded and executed during the following Q1 through Q4. The clocks and instruction execution flow is shown in Figure 3-2.

#### 3.2 Instruction Flow/Pipelining

An "Instruction Cycle" consists of four Q cycles (Q1, Q2, Q3 and Q4). The instruction fetch and execute are pipelined such that fetch takes one instruction cycle while decode and execute takes another instruction cycle. However, due to the pipelining, each instruction effectively executes in one cycle. If an instruction causes the program counter to change (e.g., GOTO) then two cycles are required to complete the instruction (Example 3-1).

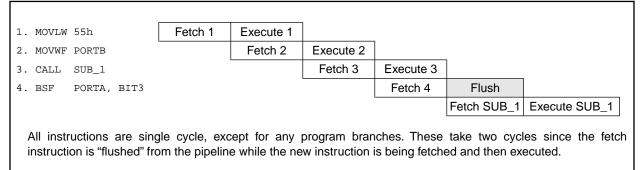
A fetch cycle begins with the Program Counter (PC) incrementing in Q1.

In the execution cycle, the fetched instruction is latched into the "Instruction Register" in cycle Q1. This instruction is then decoded and executed during the Q2, Q3, and Q4 cycles. Data memory is read during Q2 (operand read) and written during Q4 (destination write).



#### FIGURE 3-2: CLOCK/INSTRUCTION CYCLE

#### EXAMPLE 3-1: INSTRUCTION PIPELINE FLOW



#### 4.0 MEMORY ORGANIZATION

There are two memory blocks in the PIC16C84. These are the program memory and the data memory. Each block has its own bus, so that access to each block can occur during the same oscillator cycle.

The data memory can further be broken down into the general purpose RAM and the Special Function Registers (SFRs). The operation of the SFRs that control the "core" are described here. The SFRs used to control the peripheral modules are described in the section discussing each individual peripheral module.

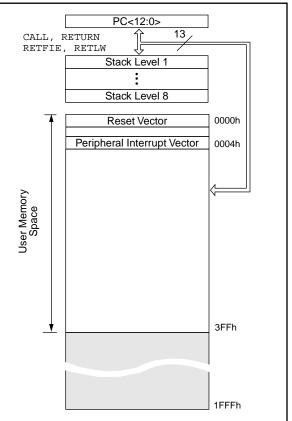
The data memory area also contains the data EEPROM memory. This memory is not directly mapped into the data memory, but is indirectly mapped. That is an indirect address pointer specifies the address of the data EEPROM memory to read/write. The 64 bytes of data EEPROM memory have the address range 0h-3Fh. More details on the EEPROM memory can be found in Section 7.0.

#### 4.1 Program Memory Organization

The PIC16CXX has a 13-bit program counter capable of addressing an 8K x 14 program memory space. For the PIC16C84, only the first 1K x 14 (0000h-03FFh) are physically implemented (Figure 4-1). Accessing a location above the physically implemented address will cause a wraparound. For example, locations 20h, 420h, 820h, C20h, 1020h, 1420h, 1820h, and 1C20h will be the same instruction.

The reset vector is at 0000h and the interrupt vector is at 0004h.

#### FIGURE 4-1: PROGRAM MEMORY MAP AND STACK



#### 4.2 Data Memory Organization

The data memory is partitioned into two areas. The first is the Special Function Registers (SFR) area, while the second is the General Purpose Registers (GPR) area. The SFRs control the operation of the device.

Portions of data memory are banked. This is for both the SFR area and the GPR area. The GPR area is banked to allow greater than 116 bytes of general purpose RAM. The banked areas of the SFR are for the registers that control the peripheral functions. Banking requires the use of control bits for bank selection. These control bits are located in the STATUS Register. Figure 4-2 shows the data memory map organization.

Instructions MOVWF and MOVF can move values from the W register to any location in the register file ("F"), and vice-versa.

The entire data memory can be accessed either directly using the absolute address of each register file or indirectly through the File Select Register (FSR) (Section 4.5). Indirect addressing uses the present value of the RP1:RP0 bits for access into the banked areas of data memory.

Data memory is partitioned into two banks which contain the general purpose registers and the special function registers. Bank 0 is selected by clearing the RP0 bit (STATUS<5>). Setting the RP0 bit selects Bank 1. Each Bank extends up to 7Fh (128 bytes). The first twelve locations of each Bank are reserved for the Special Function Registers. The remainder are General Purpose Registers implemented as static RAM.

#### 4.2.1 GENERAL PURPOSE REGISTER FILE

All devices have some amount of General Purpose Register (GPR) area. Each GPR is 8 bits wide and is accessed either directly or indirectly through the FSR (Section 4.5).

The GPR addresses in bank 1 are mapped to addresses in bank 0. As an example, addressing location 0Ch or 8Ch will access the same GPR.

#### 4.2.2 SPECIAL FUNCTION REGISTERS

The Special Function Registers (Figure 4-2 and Table 4-1) are used by the CPU and Peripheral functions to control the device operation. These registers are static RAM.

The special function registers can be classified into two sets, core and peripheral. Those associated with the core functions are described in this section. Those related to the operation of the peripheral features are described in the section for that specific feature.

#### FIGURE 4-2: REGISTER FILE MAP

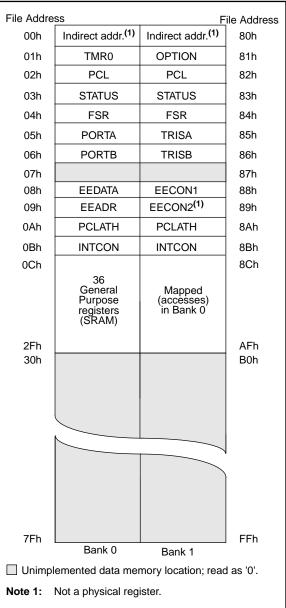


TABLE 4-1	REGISTER FILE SUMMARY
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Address	Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on Power-on Reset	Value on all other resets (Note3)
					E	Bank 0					
00h	INDF		Uses con	tents of FS	R to address da	ta memory	(not a physi	cal registe	r)		
01h	TMR0				8-bit real-time	clock/counte	er			xxxx xxxx	uuuu uuuu
02h	PCL			Low or	ler 8 bits of the I	Program Co	unter (PC)			0000 0000	0000 0000
03h	STATUS (2)	IRP	RP1	RP0	TO	PD	Z	DC	С	0001 1xxx	000q quuu
04h	FSR			Indir	ect data memor	y address p	ointer 0			xxxx xxxx	uuuu uuuu
05h	PORTA	—	—	—	RA4/T0CKI	RA3	RA2	RA1	RA0	x xxxx	u uuuu
06h	PORTB	RB7	RB6	RB5	RB4	RB3	RB2	RB1	RB0/INT	xxxx xxxx	uuuu uuuu
07h		Unimple	mented loc	ation, read	as '0'						
08h	EEDATA				EEPROM da	ata register				xxxx xxxx	uuuu uuuu
09h	EEADR				EEPROM add	lress registe	r			xxxx xxxx	uuuu uuuu
0Ah	PCLATH	_	_		Write	buffer for u	pper 5 bits o	of the PC (	1)	0 0000	0 0000
0Bh	INTCON	GIE	EEIE	TOIE	INTE	RBIE	TOIF	INTF	RBIF	x000 0000x	0000 000u
					E	Bank 1					
80h	INDF		Uses con	tents of FS	R to address da	ta memory	(not a physi	cal registe	r)		
81h	OPTION_ REG	RBPU	INTEDG	TOCS	TOSE	PSA	PS2	PS1	PS0	1111 1111	1111 1111
82h	PCL			Low c	order 8 bits of Pr	ogram Cour	nter (PC)			0000 0000	0000 0000
83h	STATUS (2)	IRP	RP1	RP0	TO	PD	Z	DC	С	0001 1xxx	000q quuu
84h	FSR			Indir	ect data memor	y address p	ointer 0			xxxx xxxx	uuuu uuuu
85h	TRISA	—	—	—	PORTA data d	irection regi	ster			1 1111	1 1111
86h	TRISB				PORTB data dir	rection regis	ter			1111 1111	1111 1111
87h		Unimple	mented loc	ation, read	as '0'						
88h	EECON1	—	—	—	EEIF	WRERR	WREN	WR	RD	0 x000	0 q000
89h	EECON2			EEPROM	control register	2 (not a phy	sical registe	r)			
0Ah	PCLATH	_	—	—	Write	buffer for u	pper 5 bits o	of the PC (	1)	0 0000	0 0000
0Bh	INTCON	GIE	EEIE	TOIE	INTE	RBIE	TOIF	INTF	RBIF	x000 0000	0000 000u

Legend: x = unknown, u = unchanged. - = unimplemented read as '0', q = value depends on condition.

**Note 1:** The upper byte of the program counter is not directly accessible. PCLATH is a slave register for PC<12:8>. The contents of PCLATH can be transferred to the upper byte of the program counter, but the contents of PC<12:8> is never transferred to PCLATH.

**2:** The  $\overline{\text{TO}}$  and  $\overline{\text{PD}}$  status bits in the STATUS register are not affected by a  $\overline{\text{MCLR}}$  reset.

3: Other (non power-up) resets include: external reset through MCLR and the Watchdog Timer Reset.

#### 4.2.2.1 STATUS REGISTER

The STATUS register contains the arithmetic status of the ALU, the RESET status and the bank select bit for data memory.

As with any register, the STATUS register can be the destination for any instruction. If the STATUS register is the destination for an instruction that affects the Z, DC or C bits, then the write to these three bits is disabled. These bits are set or cleared according to device logic. Furthermore, the  $\overline{TO}$  and  $\overline{PD}$  bits are not writable. Therefore, the result of an instruction with the STATUS register as destination may be different than intended.

For example, CLRF STATUS will clear the upper-three bits and set the Z bit. This leaves the STATUS register as 000u uluu (where u = unchanged).

Only the BCF, BSF, SWAPF and MOVWF instructions should be used to alter the STATUS register (Table 9-2) because these instructions do not affect any status bit.

- Note 1: The IRP and RP1 bits (STATUS<7:6>) are not used by the PIC16C84 and should be programmed as cleared. Use of these bits as general purpose R/W bits is NOT recommended, since this may affect upward compatibility with future products.
- Note 2: The C and DC bits operate as a borrow and digit borrow out bit, respectively, in subtraction. See the SUBLW and SUBWF instructions for examples.
- Note 3: When the STATUS register is the destination for an instruction that affects the Z, DC or C bits, then the write to these three bits is disabled. The specified bit(s) will be updated according to device logic

R/W-0	R/W-0	R/W-0	R-1	R-1	R/W-x	R/W-x	R/W-x			
IRP bit7	RP1	RP0	TO	PD	Z	DC	C bit0	R = Readable bit W = Writable bit U = Unimplemented bit, read as '0' - n = Value at POR reset		
bit 7:	<ul> <li>IRP: Register Bank Select bit (used for indirect addressing)</li> <li>0 = Bank 0, 1 (00h - FFh)</li> <li>1 = Bank 2, 3 (100h - 1FFh)</li> <li>The IRP bit is not used by the PIC16C8X. IRP should be maintained clear.</li> </ul>									
bit 6-5:	<ul> <li>-5: RP1:RP0: Register Bank Select bits (used for direct addressing)</li> <li>00 = Bank 0 (00h - 7Fh)</li> <li>01 = Bank 1 (80h - FFh)</li> <li>10 = Bank 2 (100h - 17Fh)</li> <li>11 = Bank 3 (180h - 1FFh)</li> <li>Each bank is 128 bytes. Only bit RP0 is used by the PIC16C8X. RP1 should be maintained clear.</li> </ul>									
bit 4:	<ul> <li>To: Time-out bit</li> <li>1 = After power-up, CLRWDT instruction, or SLEEP instruction</li> <li>0 = A WDT time-out occurred</li> </ul>									
bit 3:	1 = After	er-down bi power-up cecution of	or by the							
bit 2:	1 = The r	<ul> <li>0 = By execution of the SLEEP instruction</li> <li>Z: Zero bit</li> <li>1 = The result of an arithmetic or logic operation is zero</li> <li>0 = The result of an arithmetic or logic operation is not zero</li> </ul>								
bit 1:	<ul> <li>DC: Digit carry/borrow bit (for ADDWF and ADDLW instructions) (For borrow the polarity is reversed)</li> <li>1 = A carry-out from the 4th low order bit of the result occurred</li> <li>0 = No carry-out from the 4th low order bit of the result</li> </ul>									
bit 0:	<ul> <li>C: Carry/borrow bit (for ADDWF and ADDLW instructions)</li> <li>1 = A carry-out from the most significant bit of the result occurred</li> <li>0 = No carry-out from the most significant bit of the result occurred</li> <li>Note: For borrow the polarity is reversed. A subtraction is executed by adding the two's complement of the second operand. For rotate (RRF, RLF) instructions, this bit is loaded with either the high or low order bit of the source register.</li> </ul>									

#### FIGURE 4-3: STATUS REGISTER (ADDRESS 03h, 83h)

When the prescaler is assigned to

the WDT (PSA = '1'), TMR0 has a 1:1

prescaler assignment.

#### 4.2.2.2 OPTION\_REG REGISTER

The OPTION\_REG register is a readable and writable register which contains various control bits to configure the TMR0/WDT prescaler, the external INT interrupt, TMR0, and the weak pull-ups on PORTB.

#### FIGURE 4-4: OPTION\_REG REGISTER (ADDRESS 81h)

R/W-1	R/W-1	R/W-1	R/W-1	R/W-1	R/W-1	R/W-1	R/W-1				
RBPU	INTEDG	TOCS	TOSE	PSA	PS2	PS1	PS0	R = Readable bit			
bit7					1		bit0	W = Writable bit U = Unimplemented bit, read as '0' - n = Value at POR reset			
bit 7:	<ul> <li>7: RBPU: PORTB Pull-up Enable bit</li> <li>1 = PORTB pull-ups are disabled</li> <li>0 = PORTB pull-ups are enabled (by individual port latch values)</li> </ul>										
bit 6:	INTEDG: Interrupt Edge Select bit 1 = Interrupt on rising edge of RB0/INT pin 0 = Interrupt on falling edge of RB0/INT pin										
bit 5:	<b>TOCS</b> : TMR0 Clock Source Select bit 1 = Transition on RA4/T0CKI pin 0 = Internal instruction cycle clock (CLKOUT)										
bit 4:	<b>TOSE</b> : TMR0 Source Edge Select bit 1 = Increment on high-to-low transition on RA4/T0CKI pin 0 = Increment on low-to-high transition on RA4/T0CKI pin										
bit 3:	<b>PSA</b> : Pres 1 = Presca 0 = Presca	aler assig	ned to the	e WDT							
bit 2-0:	PS2:PS0:	Prescale	r Rate Se	lect bits							
	Bit Value	TMR0 R	ate WD	T Rate							
	000 001 010 011 100 101 110 111	1:2 1:4 1:8 1:16 1:32 1:64 1:12	1 2 1 4 1 28 1	: 1 : 2 : 4 : 8 : 16 : 32 : 64 : 128							

Note:

# **PIC16C84**

#### 4.2.2.3 INTCON REGISTER

The INTCON register is a readable and writable register which contains the various enable bits for all interrupt sources.

# Note: Interrupt flag bits get set when an interrupt condition occurs regardless of the state of its corresponding enable bit or the global enable bit, GIE (INTCON<7>).

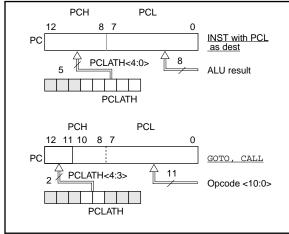
#### FIGURE 4-5: INTCON REGISTER (ADDRESS 0Bh, 8Bh)

R/W-0 GIE	R/W-0 EEIE	R/W-0 T0IE	R/W-0	R/W-0 RBIE	R/W-0 T0IF	R/W-0	R/W-x RBIF	R = Readable bit			
bit7					1		bit0	<ul> <li>W = Writable bit</li> <li>U = Unimplemented bit, read as '0'</li> <li>n = Value at POR reset</li> </ul>			
oit 7:	GIE: Global Interrupt Enable bit 1 = Enables all un-masked interrupts 0 = Disables all interrupts										
	Note: Fo	r the opera	ation of th	e interrupt	structure,	please refe	er to Section	า 8.5.			
bit 6:	<b>EEIE</b> : EE Write Complete Interrupt Enable bit 1 = Enables the EE write complete interrupt 0 = Disables the EE write complete interrupt										
bit 5:	1 = Enab	R0 Overflo les the TM ples the TM	IR0 interro	Jpt	bit						
bit 4:	1 = Enab	80/INT Inte les the RE bles the RI	BO/INT inte	errupt							
bit 3:	<b>RBIE</b> : RB Port Change Interrupt Enable bit 1 = Enables the RB port change interrupt 0 = Disables the RB port change interrupt										
bit 2:	<b>TOIF</b> : TMR0 overflow interrupt flag bit 1 = TMR0 has overflowed (must be cleared in software) 0 = TMR0 did not overflow										
bit 1:	INTF: RB0/INT Interrupt Flag bit 1 = The RB0/INT interrupt occurred 0 = The RB0/INT interrupt did not occur										
bit 0:	<b>RBIF</b> : RB Port Change Interrupt Flag bit 1 = When at least one of the RB7:RB4 pins changed state (must be cleared in software) 0 = None of the RB7:RB4 pins have changed state										

#### 4.3 Program Counter: PCL and PCLATH

The Program Counter (PC) is 13-bits wide. The low byte is the PCL register, which is a readable and writable register. The high byte of the PC (PC<12:8>) is not directly readable nor writable and comes from the PCLATH register. The PCLATH (PC latch high) register is a holding register for PC<12:8>. The contents of PCLATH are transferred to the upper byte of the program counter when the PC is loaded with a new value. This occurs during a CALL, GOTO or a write to PCL. The high bits of PC are loaded from PCLATH as shown in Figure 4-6.

#### FIGURE 4-6: LOADING OF PC IN DIFFERENT SITUATIONS



#### 4.3.1 COMPUTED GOTO

A computed GOTO is accomplished by adding an offset to the program counter (ADDWF PCL). When doing a table read using a computed GOTO method, care should be exercised if the table location crosses a PCL memory boundary (each 256 word block). Refer to the application note *"Implementing a Table Read"* (AN556).

#### 4.3.2 PROGRAM MEMORY PAGING

The PIC16C84 has 1K of program memory. The CALL and GOTO instructions have an 11-bit address range. This 11-bit address range allows a branch within a 2K program memory page size. For future PIC16CXX program memory expansion, there must be another two bits to specify the program memory page. These paging bits come from the PCLATH<4:3> bits (Figure 4-6). When doing a CALL or a GOTO instruction, the user must ensure that these page bits (PCLATH<4:3>) are programmed to the desired program memory page. If a CALL instruction (or interrupt) is executed, the entire 13-bit PC is "pushed" onto the stack (see next section). Therefore, manipulation of the PCLATH<4:3> is not required for the return instructions (which "pops" the PC from the stack).

Note:	The PIC16C84 ignores the PCLATH<4:3>							
	bits, which are used for program memory							
	pages 1, 2 and 3 (0800h - 1FFFh). The use							
	of PCLATH<4:3> as general purpose R/W							
	bits is not recommended since this may							
	affect upward compatibility with future							
	products.							

#### 4.4 Stack

The PIC16C84 has an 8 deep x 13-bit wide hardware stack (Figure 4-1). The stack space is not part of either program or data space and the stack pointer is not readable or writable.

The entire 13-bit PC is "pushed" onto the stack when a CALL instruction is executed or an interrupt is acknowledged. The stack is "popped" in the event of a RETURN, RETLW or a RETFIE instruction execution. PCLATH is not affected by a push or a pop operation.

Note:	There are no instruction mnemonics called push or pop. These are actions that occur						
	from the execution of the CALL, RETURN,						
	RETLW, and RETFIE instructions, or the						
	vectoring to an interrupt address.						

The stack operates as a circular buffer. That is, after the stack has been pushed eight times, the ninth push overwrites the value that was stored from the first push. The tenth push overwrites the second push (and so on).

If the stack is effectively popped nine times, the PC value is the same as the value from the first pop.

**Note:** There are no status bits to indicate stack overflow or stack underflow conditions.

#### 4.5 Indirect Addressing; INDF and FSR Registers

The INDF register is not a physical register. Addressing INDF actually addresses the register whose address is contained in the FSR register (FSR is a *pointer*). This is indirect addressing.

#### EXAMPLE 4-1: INDIRECT ADDRESSING

- Register file 05 contains the value 10h
- Register file 06 contains the value 0Ah
- Load the value 05 into the FSR register
- A read of the INDF register will return the value of 10h
- Increment the value of the FSR register by one (FSR = 06)
- A read of the INDF register now will return the value of 0Ah.

Reading INDF itself indirectly (FSR = 0) will produce 00h. Writing to the INDF register indirectly results in a no-operation (although STATUS bits may be affected).

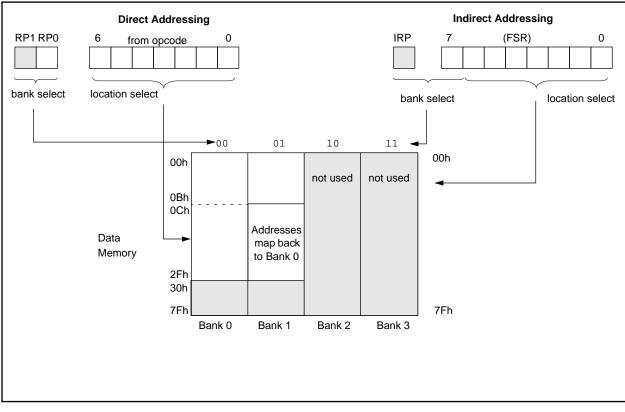
FIGURE 4-7: DIRECT/INDIRECT ADDRESSING

A simple program to clear RAM locations 20h-2Fh using indirect addressing is shown in Example 4-2.

#### EXAMPLE 4-2: HOW TO CLEAR RAM USING INDIRECT ADDRESSING

	movlw	0x20	;initialize pointer
	movwf	FSR	; to RAM
NEXT	clrf	INDF	;clear INDF register
	incf	FSR	;inc pointer
	btfss	FSR,4	;all done?
	goto	NEXT	;NO, clear next
CONTINUE			
	:		;YES, continue

An effective 9-bit address is obtained by concatenating the 8-bit FSR register and the IRP bit (STATUS<7>), as shown in Figure 4-7. However, IRP is not used in the PIC16C84.



#### 5.0 I/O PORTS

The PIC16C84 has two ports, PORTA and PORTB. Some port pins are multiplexed with an alternate function for other features on the device.

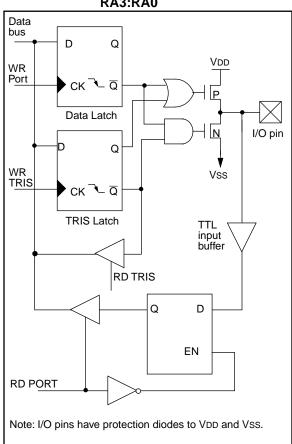
#### 5.1 PORTA and TRISA Registers

PORTA is a 5-bit wide latch. RA4 is a Schmitt Trigger input and an open drain output. All other RA port pins have TTL input levels and full CMOS output drivers. All pins have data direction bits (TRIS registers) which can configure these pins as output or input.

Setting a TRISA bit (=1) will make the corresponding PORTA pin an input, i.e., put the corresponding output driver in a hi-impedance mode. Clearing a TRISA bit (=0) will make the corresponding PORTA pin an output, i.e., put the contents of the output latch on the selected pin.

Reading the PORTA register reads the status of the pins whereas writing to it will write to the port latch. All write operations are read-modify-write operations. So a write to a port implies that the port pins are first read, then this value is modified and written to the port data latch.

The RA4 pin is multiplexed with the TMR0 clock input.

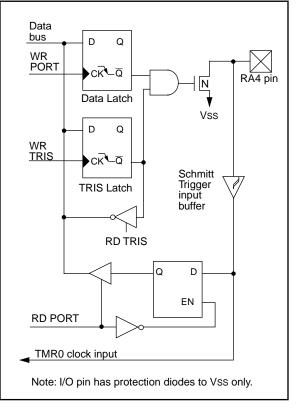


#### FIGURE 5-1: BLOCK DIAGRAM OF PINS RA3:RA0

#### EXAMPLE 5-1: INITIALIZING PORTA

CLRF	PORTA	; Initialize PORTA by
		; setting output
		; data latches
BSF	STATUS, RPO	; Select Bank 1
MOVLW	0x0F	; Value used to
		; initialize data
		; direction
MOVWF	TRISA	; Set RA<3:0> as inputs
		; RA4 as outputs
		; TRISA<7:5> are always
		; read as '0'

#### FIGURE 5-2: BLOCK DIAGRAM OF PIN RA4



Note: For crystal oscillator configurations operating below 500 kHz, the device may generate a spurious internal Q-clock when PORTA<0> switches state. This does not occur with an external clock in RC mode. To avoid this, the RA0 pin should be kept static, i.e. in input/output mode, pin RA0 should not be toggled.

#### TABLE 5-1PORTA FUNCTIONS

Name	Bit0	Buffer Type	Function
RA0	bit0	TTL	Input/output
RA1	bit1	TTL	Input/output
RA2	bit2	TTL	Input/output
RA3	bit3	TTL	Input/output
RA4/T0CKI	bit4	ST	Input/output or external clock input for TMR0. Output is open drain type.

Legend: TTL = TTL input, ST = Schmitt Trigger input

#### TABLE 5-2 SUMMARY OF REGISTERS ASSOCIATED WITH PORTA

Address	Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on Power-on Reset	Value on all other resets
05h	PORTA	—	_	—	RA4/T0CKI	RA3	RA2	RA1	RA0	x xxxx	u uuuu
85h	TRISA	_	_	—	TRISA4	TRISA3	TRISA2	TRISA1	TRISA0	1 1111	1 1111

Legend: x = unknown, u = unchanged, - = unimplemented read as '0'. Shaded cells are unimplemented, read as '0'

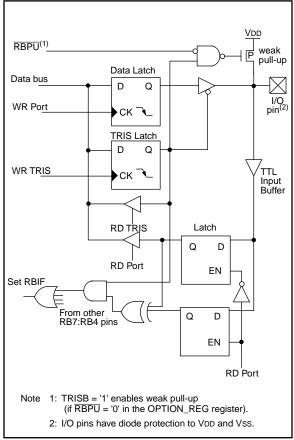
#### 5.2 **PORTB and TRISB Registers**

PORTB is an 8-bit wide bi-directional port. The corresponding data direction register is TRISB. A '1' on any bit in the TRISB register puts the corresponding output driver in a hi-impedance mode. A '0' on any bit in the TRISB register puts the contents of the output latch on the selected pin(s).

Each of the PORTB pins have a weak internal pull-up. A single control bit can turn on all the pull-ups. This is done by clearing the RBPU (OPTION\_REG<7>) bit. The weak pull-up is automatically turned off when the port pin is configured as an output. The pull-ups are disabled on a Power-on Reset.

Four of PORTB's pins, RB7:RB4, have an interrupt on change feature. Only pins configured as inputs can cause this interrupt to occur (i.e., any RB7:RB4 pin configured as an output is excluded from the interrupt on change comparison). The pins value in input mode are compared with the old value latched on the last read of PORTB. The "mismatch" outputs of the pins are OR'ed together to generate the RB port change interrupt.

#### FIGURE 5-3: BLOCK DIAGRAM OF PINS RB7:RB4



This interrupt can wake the device from SLEEP. The user, in the interrupt service routine, can clear the interrupt in the following manner:

- a) Read (or write) PORTB. This will end the mismatch condition.
- b) Clear flag bit RBIF.

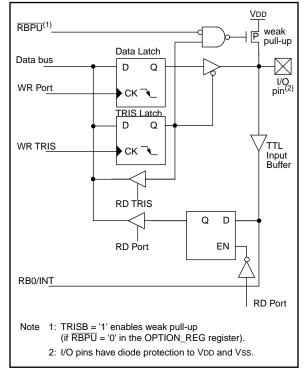
A mismatch condition will continue to set the RBIF bit. Reading PORTB will end the mismatch condition, and allow the RBIF bit to be cleared.

This interrupt on mismatch feature, together with software configurable pull-ups on these four pins allow easy interface to a key pad and make it possible for wake-up on key-depression (see AN552 in the Embedded Control Handbook).

Note 1: If a change on the I/O pin should occur when a read operation of PORTB is being executed (start of the Q2 cycle), the RBIF interrupt flag bit may not be set.

The interrupt on change feature is recommended for wake-up on key depression operation and operations where PORTB is only used for the interrupt on change feature. Polling of PORTB is not recommended while using the interrupt on change feature.

#### FIGURE 5-4: BLOCK DIAGRAM OF PINS RB3:RB0



#### EXAMPLE 5-1: INITIALIZING PORTB

CLRF	PORTB	; Initialize PORTB by
		; setting output
		; data latches
BSF	STATUS, RPO	; Select Bank 1
MOVLW	0xCF	; Value used to
		; initialize data
		; direction
MOVWF	TRISB	; Set RB<3:0> as inputs
		; RB<5:4> as outputs
		; RB<7:6> as inputs

#### TABLE 5-3 PORTB FUNCTIONS

Name	Bit	Buffer Type	I/O Consistency Function
RB0/INT	bit0	TTL	Input/output pin or external interrupt input. Internal software programmable weak pull-up.
RB1	bit1	TTL	Input/output pin. Internal software programmable weak pull-up.
RB2	bit2	TTL	Input/output pin. Internal software programmable weak pull-up.
RB3	bit3	TTL	Input/output pin. Internal software programmable weak pull-up.
RB4	bit4	TTL	Input/output pin (with interrupt on change). Internal software programma- ble weak pull-up.
RB5	bit5	TTL	Input/output pin (with interrupt on change). Internal software programma- ble weak pull-up.
RB6	bit6	TTL/ST <sup>(1)</sup>	Input/output pin (with interrupt on change). Internal software programma- ble weak pull-up. Serial programming clock.
RB7	bit7	TTL/ST <sup>(1)</sup>	Input/output pin (with interrupt on change). Internal software programma- ble weak pull-up. Serial programming data.

Legend: TTL = TTL input, ST = Schmitt Trigger.

**Note 1:** This buffer is a Schmitt Trigger input when used in serial programming mode.

#### TABLE 5-4SUMMARY OF REGISTERS ASSOCIATED WITH PORTB

Address	Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on Power-on Reset	Value on all other resets
06h	PORTB	RB7	RB6	RB5	RB4	RB3	RB2	RB1	RB0/INT	xxxx xxxx	uuuu uuuu
86h	TRISB	TRISB7	TRISB6	TRISB5	TRISB4	TRISB3	TRISB2	TRISB1	TRISB0	1111 1111	1111 1111
81h	OPTION_ REG	RBPU	INTEDG	TOCS	TOSE	PSA	PS2	PS1	PS0	1111 1111	1111 1111

Legend: x = unknown, u = unchanged. Shaded cells are not used by PORTB.

#### 5.3 I/O Programming Considerations

#### 5.3.1 BI-DIRECTIONAL I/O PORTS

Any instruction which writes, operates internally as a read followed by a write operation. The BCF and BSF instructions, for example, read the register into the CPU, execute the bit operation and write the result back to the register. Caution must be used when these instructions are applied to a port with both inputs and outputs defined. For example, a BSF operation on bit5 of PORTB will cause all eight bits of PORTB to be read into the CPU. Then the BSF operation takes place on bit5 and PORTB is written to the output latches. If another bit of PORTB is used as a bi-directional I/O pin (i.e., bit0) and it is defined as an input at this time, the input signal present on the pin itself would be read into the CPU and rewritten to the data latch of this particular pin, overwriting the previous content. As long as the pin stays in the input mode, no problem occurs. However, if bit0 is switched into output mode later on, the content of the data latch is unknown.

Reading the port register, reads the values of the port pins. Writing to the port register writes the value to the port latch. When using read-modify-write instructions (i.e., BCF, BSF, etc.) on a port, the value of the port pins is read, the desired operation is done to this value, and this value is then written to the port latch.

A pin actively outputting a Low or High should not be driven from external devices at the same time in order to change the level on this pin ("wired-or", "wired-and"). The resulting high output current may damage the chip.

#### 5.3.2 SUCCESSIVE OPERATIONS ON I/O PORTS

The actual write to an I/O port happens at the end of an instruction cycle, whereas for reading, the data must be valid at the beginning of the instruction cycle (Figure 5-5). Therefore, care must be exercised if a write followed by a read operation is carried out on the same I/O port. The sequence of instructions should be such that the pin voltage stabilizes (load dependent) before the next instruction which causes that file to be read into the CPU is executed. Otherwise, the previous state of that pin may be read into the CPU rather than the new state. When in doubt, it is better to separate these instructions with a NOP or another instruction not accessing this I/O port.

Example 5-1 shows the effect of two sequential read-modify-write instructions (e.g.,  ${\tt BCF}$ ,  ${\tt BSF}$ , etc.) on an I/O port.

#### EXAMPLE 5-1: READ-MODIFY-WRITE INSTRUCTIONS ON AN I/O PORT

;Initial PORT settings: PORTB<7:4> Inputs
; PORTB<3:0> Outputs
;PORTB<7:6> have external pull-ups and are
;not connected to other circuitry
;

;					PORT	latch	PORT	pins
;								
	BCF	PORTB,	7	;	01pp	ppp	11pp	ppp
	BCF	PORTB,	б	;	10pp	ppp	11pp	ppp
	BSF	STATUS,	RP0	;				
	BCF	TRISB,	7	;	10pp	ppp	11pp	ppp
	BCF	TRISB,	б	;	10pp	ppp	10pp	ppp
;								

;Note that the user may have expected the ;pin values to be 00pp ppp. The 2nd BCF ;caused RB7 to be latched as the pin value ;(high).

PC Instruction	X PC	PC + 1	X PC + 2 X	PC + 3	This example shows a write to POR
fetched	MOVWF PORTB write to PORTB	MOVF PORTB,W	NOP	NOP	followed by a read from PORTB. Note that:
RB7:RB0			X	;	data setup time = (0.25Tcy - TPD)
		1 1 1	Port pin sampled here		where TCY = instruction cycle TPD = propagation delay
Instruction executed		MOVWF PORTB write to PORTB	TPD	NOP	Therefore, at higher clock frequenc a write followed by a read may problematic.
		1			

#### FIGURE 5-5: SUCCESSIVE I/O OPERATION

NOTES:

#### 6.0 TIMER0 MODULE AND TMR0 REGISTER

The Timer0 module timer/counter has the following features:

- 8-bit timer/counter
- Readable and writable
- 8-bit software programmable prescaler
- Internal or external clock select
- Interrupt on overflow from FFh to 00h
- · Edge select for external clock

Timer mode is selected by clearing the TOCS bit (OPTION<5>). In timer mode, the Timer0 module (Figure 6-1) will increment every instruction cycle (without prescaler). If the TMR0 register is written, the increment is inhibited for the following two cycles (Figure 6-2 and Figure 6-3). The user can work around this by writing an adjusted value to the TMR0 register.

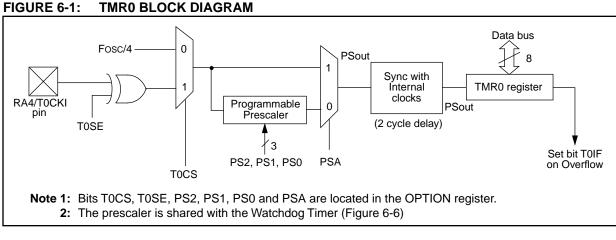
Counter mode is selected by setting the TOCS bit (OPTION<5>). In this mode TMR0 will increment either on every rising or falling edge of pin RA4/T0CKI. The incrementing edge is determined by the T0 source

edge select bit, T0SE (OPTION<4>). Clearing bit T0SE selects the rising edge. Restrictions on the external clock input are discussed in detail in Section 6.2.

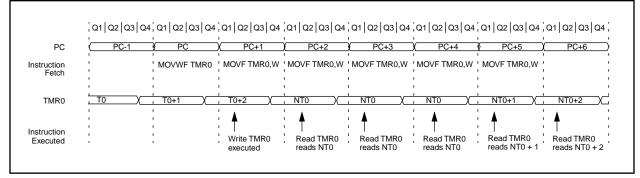
The prescaler is shared between the Timer0 Module and the Watchdog Timer. The prescaler assignment is controlled, in software, by control bit PSA (OPTION<3>). Clearing bit PSA will assign the prescaler to the Timer0 Module. The prescaler is not readable or writable. When the prescaler (Section 6.3) is assigned to the Timer0 Module, the prescale value (1:2, 1:4, ..., 1:256) is software selectable.

#### 6.1 TMR0 Interrupt

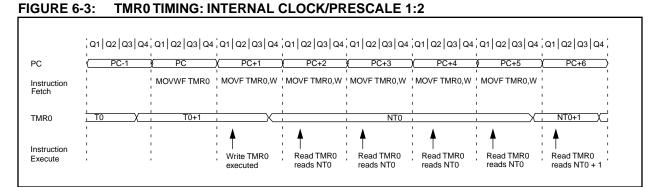
The TMR0 interrupt is generated when the TMR0 register overflows from FFh to 00h. This overflow sets the T0IF bit (INTCON<2>). The interrupt can be masked by clearing enable bit T0IE (INTCON<5>). The T0IF bit must be cleared in software by the Timer0 Module interrupt service routine before re-enabling this interrupt. The TMR0 interrupt (Figure 6-4) cannot wake the processor from SLEEP since the timer is shut off during SLEEP.



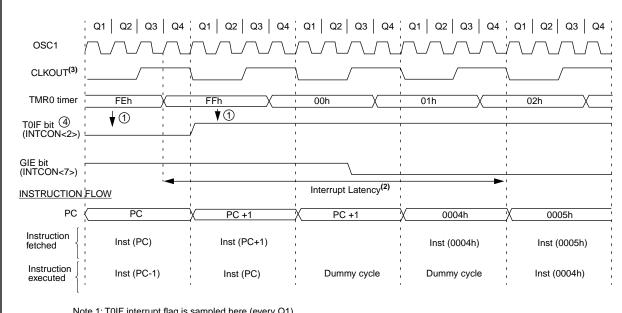




# **PIC16C84**



#### FIGURE 6-4: **TMR0 INTERRUPT TIMING**



Note 1: T0IF interrupt flag is sampled here (every Q1).

2: Interrupt latency = 3.25Tcy, where Tcy = instruction cycle time.

3: CLKOUT is available only in RC oscillator mode. 4: The timer clock (after the synchronizer circuit) which increments the timer from FFh to 00h immediately sets the T0IF bit. The TMR0 register will roll over 3 Tosc cycles later.

#### 6.2 Using TMR0 with External Clock

When an external clock input is used for TMR0, it must meet certain requirements. The external clock requirement is due to internal phase clock (Tosc) synchronization. Also, there is a delay in the actual incrementing of the TMR0 register after synchronization.

#### 6.2.1 EXTERNAL CLOCK SYNCHRONIZATION

When no prescaler is used, the external clock input is the same as the prescaler output. The synchronization of pin RA4/T0CKI with the internal phase clocks is accomplished by sampling the prescaler output on the Q2 and Q4 cycles of the internal phase clocks (Figure 6-5). Therefore, it is necessary for T0CKI to be high for at least 2Tosc (plus a small RC delay) and low for at least 2Tosc (plus a small RC delay). Refer to the electrical specification of the desired device.

When a prescaler is used, the external clock input is divided by an asynchronous ripple counter type prescaler so that the prescaler output is symmetrical. For the external clock to meet the sampling requirement, the ripple counter must be taken into account. Therefore, it is necessary for TOCKI to have a period of at least 4Tosc (plus a small RC delay) divided by the prescaler value. The only requirement on TOCKI high and low time is that they do not violate the minimum pulse width requirement of 10 ns. Refer to parameters 40, 41 and 42 in the AC Electrical Specifications of the desired device.

#### 6.2.2 TMR0 INCREMENT DELAY

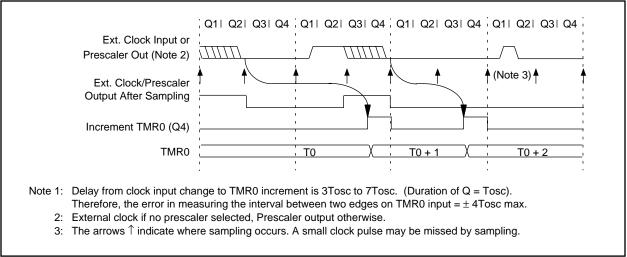
Since the prescaler output is synchronized with the internal clocks, there is a small delay from the time the external clock edge occurs to the time the Timer0 Module is actually incremented. Figure 6-5 shows the delay from the external clock edge to the timer incrementing.

#### 6.3 <u>Prescaler</u>

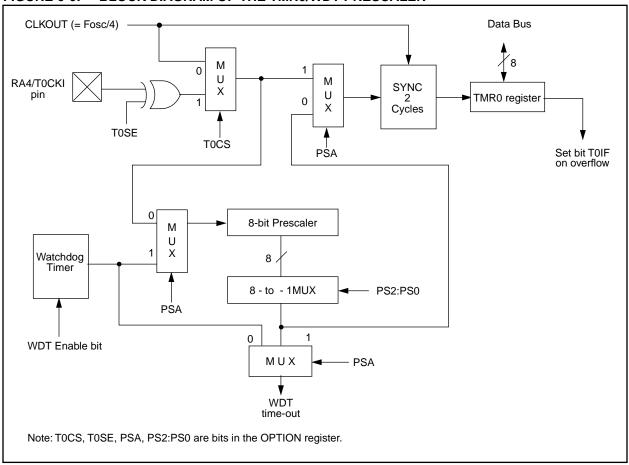
An 8-bit counter is available as a prescaler for the Timer0 Module, or as a postscaler for the Watchdog Timer (Figure 6-6). For simplicity, this counter is being referred to as "prescaler" throughout this data sheet. Note that there is only one prescaler available which is mutually exclusive between the Timer0 Module and the Watchdog Timer. Thus, a prescaler assignment for the Timer0 Module means that there is no prescaler for the Watchdog Timer, and vice-versa.

The PSA and PS2:PS0 bits (OPTION<3:0>) determine the prescaler assignment and prescale ratio.

When assigned to the Timer0 Module, all instructions writing to the Timer0 Module (e.g., CLRF 1, MOVWF 1, BSF 1, x ....etc.) will clear the prescaler. When assigned to WDT, a CLRWDT instruction will clear the prescaler along with the Watchdog Timer. The prescaler is not readable or writable.



#### FIGURE 6-5: TIMER0 TIMING WITH EXTERNAL CLOCK



#### FIGURE 6-6: BLOCK DIAGRAM OF THE TMR0/WDT PRESCALER

#### 6.3.1 SWITCHING PRESCALER ASSIGNMENT

The prescaler assignment is fully under software control (i.e., it can be changed "on the fly" during program execution).

Note: To avoid an unintended device RESET, the following instruction sequence (Example 6-1) must be executed when changing the prescaler assignment from Timer0 to the WDT. This sequence must be taken even if the WDT is disabled. To change prescaler from the WDT to the Timer0 module use the sequence shown in Example 6-2.

## EXAMPLE 6-1: CHANGING PRESCALER (TIMER0 $\rightarrow$ WDT)

BCF	STATUS, RPO	;Bank 0
CLRF	TMR0	;Clear TMR0
		; and Prescaler
BSF	STATUS, RPO	;Bank 1
CLRWDT		;Clears WDT
MOVLW	b'xxxx1xxx'	;Select new
MOVWF	OPTION	; prescale value
BCF	STATUS, RPO	;Bank 0

#### EXAMPLE 6-2: CHANGING PRESCALER (WDT→TIMER0)

CLRWDT		;Clear WDT and
		; prescaler
BSF	STATUS, RPO	;Bank 1
MOVLW	b'xxxx0xxx'	;Select TMR0, new
		; prescale value
		' and clock source
MOVWF	OPTION	;
BCF	STATUS, RPO	;Bank 0

#### TABLE 6-1 REGISTERS ASSOCIATED WITH TIMER0

Address	Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on Power-on Reset	Value on all other resets
01h	TMR0				Timer0 modu	le's register				xxxx xxxx	uuuu uuuu
0Bh	INTCON	GIE	EEIE	TOIE	INTE	RBIE	TOIF	INTF	RBIF	0000 000x	0000 0000
81h	OPTION	RBPU	INTEDG	TOCS	TOSE	PSA	PS2	PS1	PS0	1111 1111	1111 1111
85h	TRISA	—	_	—	TRISA4	TRISA3	TRISA2	TRISA1	TRISA0	1 1111	1 1111

Legend: x = unknown, u = unchanged. - = unimplemented read as '0'. Shaded cells are not associated with Timer0.

NOTES:

#### 7.0 DATA EEPROM MEMORY

The EEPROM data memory is readable and writable during normal operation (full VDD range). This memory is not directly mapped in the register file space. Instead it is indirectly addressed through the Special Function Registers. There are four SFRs used to read and write this memory. These registers are:

- EECON1
- EECON2
- EEDATA
- EEADR

EEDATA holds the 8-bit data for read/write, and EEADR holds the address of the EEPROM location being accessed. PIC16C84 devices have 64 bytes of data EEPROM with an address range from 0h to 3Fh.

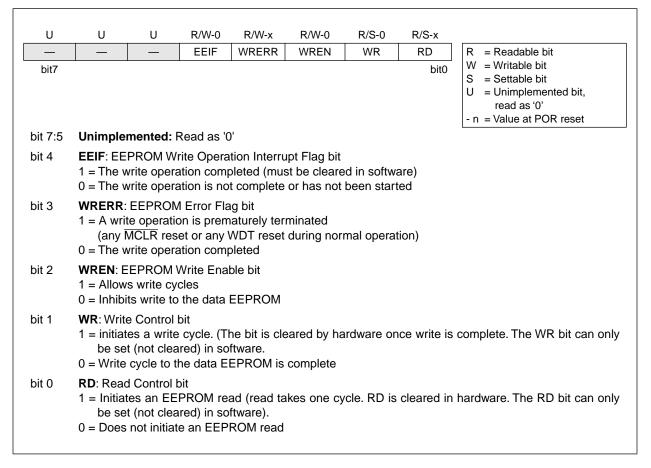
The EEPROM data memory allows byte read and write. A byte write automatically erases the location and writes the new data (erase before write). The EEPROM data memory is rated for high erase/write cycles. The write time is controlled by an on-chip timer. The writetime will vary with voltage and temperature as well as from chip to chip. Please refer to AC specifications for exact limits. When the device is code protected, the CPU may continue to read and write the data EEPROM memory. The device programmer can no longer access this memory.

#### 7.1 <u>EEADR</u>

The EEADR register can address up to a maximum of 256 bytes of data EEPROM. Only the first 64 bytes of data EEPROM are implemented.

The upper two bits are address decoded. This means that these two bits must always be '0' to ensure that the address is in the 64 byte memory space.

#### FIGURE 7-1: EECON1 REGISTER (ADDRESS 88h)



#### 7.2 EECON1 and EECON2 Registers

EECON1 is the control register with five low order bits physically implemented. The upper-three bits are nonexistent and read as '0's.

Control bits RD and WR initiate read and write, respectively. These bits cannot be cleared, only set, in software. They are cleared in hardware at completion of the read or write operation. The inability to clear the WR bit in software prevents the accidental, premature termination of a write operation.

The WREN bit, when set, will allow a write operation. On power-up, the WREN bit is clear. The WRERR bit is set when a write operation is interrupted by a  $\overline{\text{MCLR}}$  reset or a WDT time-out reset during normal operation. In these situations, following reset, the user can check the WRERR bit and rewrite the location. The data and address will be unchanged in the EEDATA and EEADR registers.

Interrupt flag bit EEIF is set when write is complete. It must be cleared in software.

EECON2 is not a physical register. Reading EECON2 will read all '0's. The EECON2 register is used exclusively in the Data EEPROM write sequence.

#### 7.3 Reading the EEPROM Data Memory

To read a data memory location, the user must write the address to the EEADR register and then set control bit RD (EECON1<0>). The data is available, in the very next cycle, in the EEDATA register; therefore it can be read in the next instruction. EEDATA will hold this value until another read or until it is written to by the user (during a write operation).

#### EXAMPLE 7-1: DATA EEPROM READ

BCF	STATUS, RPO	; Bank 0
MOVLW	CONFIG_ADDR	;
MOVWF	EEADR	; Address to read
BSF	STATUS, RPO	; Bank 1
BSF	EECON1, RD	; EE Read
BCF	STATUS, RPO	; Bank 0
MOVF	EEDATA, W	; W = EEDATA

#### 7.4 Writing to the EEPROM Data Memory

To write an EEPROM data location, the user must first write the address to the EEADR register and the data to the EEDATA register. Then the user must follow a specific sequence to initiate the write for each byte.

#### EXAMPLE 7-1: DATA EEPROM WRITE

	BSF	STATUS, RPO	;	Bank 1
	BCF	INTCON, GIE	;	Disable INTs.
	BSF	EECON1, WREN	;	Enable Write
	MOVLW	55h	;	
	MOVWF	EECON2	;	Write 55h
မ ရ	MOVLW	AAh	;	
kequired sequence	MOVWF	EECON2	;	Write AAh
b g g	BSF	EECON1,WR	;	Set WR bit
a s			;	begin write
	BSF	INTCON, GIE	;	Enable INTs.

The write will not initiate if the above sequence is not exactly followed (write 55h to EECON2, write AAh to EECON2, then set WR bit) for each byte. We strongly recommend that interrupts be disabled during this code segment.

Additionally, the WREN bit in EECON1 must be set to enable write. This mechanism prevents accidental writes to data EEPROM due to errant (unexpected) code execution (i.e., lost programs). The user should keep the WREN bit clear at all times, except when updating EEPROM. The WREN bit is not cleared by hardware

After a write sequence has been initiated, clearing the WREN bit will not affect this write cycle. The WR bit will be inhibited from being set unless the WREN bit is set.

At the completion of the write cycle, the WR bit is cleared in hardware and the EE Write Complete Interrupt Flag bit (EEIF) is set. The user can either enable this interrupt or poll this bit. EEIF must be cleared by software.

Note:	The data EEPROM memory E/W cycle
	time may occasionally exceed the 10 ms
	specification (typical). To ensure that the
	write cycle is complete, use the EE
	interrupt or poll the WR bit (EECON1<1>).
	Both these events signify the completion of
	the write cycle.

#### 7.5 Write Verify

Depending on the application, good programming practice may dictate that the value written to the Data EEPROM should be verified (Example 7-1) to the desired value to be written. This should be used in applications where an EEPROM bit will be stressed near the specification limit. The Total Endurance disk will help determine your comfort level.

Generally the EEPROM write failure will be a bit which was written as a '1', but reads back as a '0' (due to leakage off the bit).

#### EXAMPLE 7-1: WRITE VERIFY

```
BCF
          STATUS, RPO ; Bank 0
                     ; Any code can go here
   :
    :
                      ;
   MOVF
        EEDATA, W ; Must be in Bank 0
         STATUS, RP0 ; Bank 1
   BSF
READ
   BSF
         EECON1, RD ; YES, Read the
                     ; value written
   BCF
          STATUS, RPO ; Bank 0
;
; Is the value written (in W reg) and
   read (in EEDATA) the same?
;
;
   SUBWF EEDATA, W
                     ;
                    ; Is difference 0?
   BTFSS STATUS, Z
   GOTO WRITE_ERR ; NO, Write error
    :
                     ; YES, Good write
                     ; Continue program
    :
```

#### 7.6 Protection Against Spurious Writes

There are conditions when the device may not want to write to the data EEPROM memory. To protect against spurious EEPROM writes, various mechanisms have been built in. On power-up, WREN is cleared. Also, the Power-up Timer (72 ms duration) prevents EEPROM write.

The write initiate sequence and the WREN bit together help prevent an accidental write during brown-out, power glitch, or software malfunction.

#### 7.7 Data EEPROM Operation during Code Protect

When the device is code protected, the CPU is able to read and write unscrambled data to the Data EEPROM.

For ROM devices, there are two code protection bits (Section 8.1). One for the ROM program memory and one for the Data EEPROM memory.

#### 7.8 Power Consumption Considerations

Note: It is recommended that the EEADR<7:6> bits be cleared. When either of these bits is set, the maximum IDD for the device is higher than when both are cleared. The specification is 400 μA. With EEADR<7:6> cleared, the maximum is approximately 150 μA.

#### TABLE 7-1 REGISTERS/BITS ASSOCIATED WITH DATA EEPROM

Address	Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on Power-on Reset	Value on all other resets
08h	EEDATA	EEPROM	EPROM data register								uuuu uuuu
09h	EEADR	EEPROM	EEPROM address register								uuuu uuuu
88h	EECON1	—	_	_	EEIF	WRERR	WREN	WR	RD	0 x000	0 d000
89h	EECON2	EEPROM	EEPROM control register 2								

Legend: x = unknown, u = unchanged, - = unimplemented read as '0', q = value depends upon condition. Shaded cells are not used by Data EEPROM.

NOTES:

# 8.0 SPECIAL FEATURES OF THE CPU

What sets a microcontroller apart from other processors are special circuits to deal with the needs of real time applications. The PIC16C84 has a host of such features intended to maximize system reliability, minimize cost through elimination of external components, provide power saving operating modes and offer code protection. These features are:

- OSC selection
- Reset
  - Power-on Reset (POR)
  - Power-up Timer (PWRT)
  - Oscillator Start-up Timer (OST)
- Interrupts
- Watchdog Timer (WDT)
- SLEEP
- Code protection
- ID locations
- In-circuit serial programming

The PIC16C84 has a Watchdog Timer which can be shut off only through configuration bits. It runs off its own RC oscillator for added reliability. There are two timers that offer necessary delays on power-up. One is the Oscillator Start-up Timer (OST), intended to keep

#### FIGURE 8-1: CONFIGURATION WORD

the chip in reset until the crystal oscillator is stable. The other is the Power-up Timer (PWRT), which provides a fixed delay of 72 ms (nominal) on power-up only. This design keeps the device in reset while the power supply stabilizes. With these two timers on-chip, most applications need no external reset circuitry.

SLEEP mode offers a very low current power-down mode. The user can wake-up from SLEEP through external reset, Watchdog Timer time-out or through an interrupt. Several oscillator options are provided to allow the part to fit the application. The RC oscillator option saves system cost while the LP crystal option saves power. A set of configuration bits are used to select the various options.

#### 8.1 Configuration Bits

The configuration bits can be programmed (read as '0') or left unprogrammed (read as '1') to select various device configurations. These bits are mapped in program memory location 2007h.

Address 2007h is beyond the user program memory space and it belongs to the special test/configuration memory space (2000h - 3FFFh). This space can only be accessed during programming.

To find out how to program the PIC16C84, refer to *PIC16C84 EEPROM Memory Programming Specifica-tion* (DS30189).

U-1	U-1 U-1	U-1	U-1	U-1	U-1	U-1	U-1	R/P-u	R/P-u	R/P-u	R/P-u	R/P-u	
_		_						CP	PWRTE	WDTE	FOSC1		
bit13												bit0	
											P = Pro U = Uni rea - n = Val	adable bit ogrammable bit implemented b ad as '1' ue at POR rese unchanged	it,
bit 13:5	Unimpleme	ented: F	Read as	s '1'									
bit 4	<b>CP</b> : Code F 1 = Code p 0 = All mem	rotection	n off	otecte	d								
bit 3	<b>PWRTE</b> : Po 1 = Power-u 0 = Power-u	up timer	is enal	bled	e bit								
bit 2	WDTE: Watchdog Timer Enable bit 1 = WDT enabled 0 = WDT disabled												
bit 1:0	<b>FOSC1:FO</b> 11 =RC os 10 = HS os 01 = XT os 00 = LP os	cillator cillator cillator	scillato	r Sele	ction t	oits							

#### 8.2 Oscillator Configurations

#### 8.2.1 OSCILLATOR TYPES

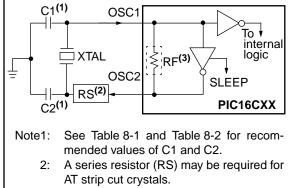
The PIC16C84 can be operated in four different oscillator modes. The user can program two configuration bits (FOSC1 and FOSC0) to select one of these four modes:

- LP Low Power Crystal
- XT Crystal/Resonator
- HS High Speed Crystal/Resonator
- RC Resistor/Capacitor

### 8.2.2 CRYSTAL OSCILLATOR / CERAMIC RESONATORS

In XT, LP or HS modes a crystal or ceramic resonator is connected to the OSC1/CLKIN and OSC2/CLKOUT pins to establish oscillation (Figure 8-2).

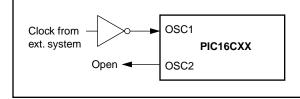
#### FIGURE 8-2: CRYSTAL/CERAMIC RESONATOR OPERATION (HS, XT OR LP OSC CONFIGURATION)



3: RF varies with the crystal chosen.

The PIC16C84 oscillator design requires the use of a parallel cut crystal. Use of a series cut crystal may give a frequency out of the crystal manufacturers specifications. When in XT, LP or HS modes, the device can have an external clock source to drive the OSC1/ CLKIN pin (Figure 8-3).

#### FIGURE 8-3: EXTERNAL CLOCK INPUT OPERATION (HS, XT OR LP OSC CONFIGURATION)



# TABLE 8-1CAPACITOR SELECTION<br/>FOR CERAMIC RESONATORS

Ranges Tested:								
Mode	Freq	OSC1/C1	OSC2/C2					
ХТ	455 kHz	47 - 100 pF	47 - 100 pF					
	2.0 MHz	15 - 33 pF	15 - 33 pF					
	4.0 MHz	15 - 33 pF	15 - 33 pF					
HS	8.0 MHz	15 - 33 pF	15 - 33 pF					
	10.0 MHz	15 - 33 pF	15 - 33 pF					
Note: Recommended values of C1 and C2 are identical to								

the ranges tested table. Higher capacitance increases the stability of the oscillator but also increases the start-up time. These values are for design guidance only. Since each resonator has its own characteristics, the user should consult the resonator manufacturer for the appropriate values of external components.

#### **Resonators Tested:**

4.0 MHz

10.0 MHz

455 kHz	Panasonic EFO-A455K04B	± 0.3%					
2.0 MHz	Murata Erie CSA2.00MG	± 0.5%					
4.0 MHz	Murata Erie CSA4.00MG	± 0.5%					
8.0 MHz	Murata Erie CSA8.00MT	± 0.5%					
10.0 MHz Murata Erie CSA10.00MTZ ± 0.5%							
None of the resonators had built-in capacitors.							

## TABLE 8-2 CAPACITOR SELECTION FOR CRYSTAL OSCILLATOR

Mode	Freq	OSC1/C1	OSC2/C2				
LP	32 kHz	68 - 100 pF	68 - 100 pF				
	200 kHz	15 - 33 pF	15 - 33 pF				
XT	100 kHz	100 - 150 pF	100 - 150 pF				
	2 MHz	15 - 33 pF	15 - 33 pF				
	4 MHz	15 - 33 pF	15 - 33 pF				
HS	4 MHz	15 - 33 pF	15 - 33 pF				
	10 MHz	15 - 33 pF	15 - 33 pF				
, v		e increases the s					
		increases the sta					
		for design guidan					
		mode as well as					
		crystals with low c ch crystal has its o					
		e user should con					
		ppropriate values					
	nponents.	FF F					
For	VDD > 4.5V, C	1 = C2 ≈ 30 pF is	recommended.				
Crystals Tested:							
32.768 kH	Iz Epson C-0	Epson C-001R32.768K-A					
100 kHz	Epson C-2	Epson C-2 100.00 KC-P					
200 kHz	STD XTL	200.000 KHz	± 20 PPM				
1.0 MHz	ECS ECS	-10-13-2	± 50 PPM				
2.0 MHz	ECS ECS	-20-S-2	± 50 PPM				

ECS ECS-40-S-4

ECS ECS-100-S-4

± 50 PPM

± 50 PPM

#### 8.2.3 EXTERNAL CRYSTAL OSCILLATOR CIRCUIT

Either a prepackaged oscillator can be used or a simple oscillator circuit with TTL gates can be built. Prepackaged oscillators provide a wide operating range and better stability. A well-designed crystal oscillator will provide good performance with TTL gates. Two types of crystal oscillator circuits are available; one with series resonance, and one with parallel resonance.

Figure 8-4 shows a parallel resonant oscillator circuit. The circuit is designed to use the fundamental frequency of the crystal. The 74AS04 inverter performs the 180-degree phase shift that a parallel oscillator requires. The 4.7 k $\Omega$  resistor provides negative feedback for stability. The 10 k $\Omega$  potentiometer biases the 74AS04 in the linear region. This could be used for external oscillator designs.

#### FIGURE 8-4: EXTERNAL PARALLEL RESONANT CRYSTAL OSCILLATOR CIRCUIT

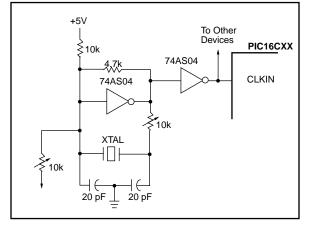
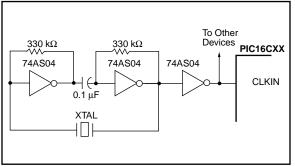


Figure 8-5 shows a series resonant oscillator circuit. This circuit is also designed to use the fundamental frequency of the crystal. The inverter performs a 180-degree phase shift. The 330 k $\Omega$  resistors provide the negative feedback to bias the inverters in their linear region.

## FIGURE 8-5: EXTERNAL SERIES RESONANT CRYSTAL OSCILLATOR CIRCUIT



## 8.2.4 RC OSCILLATOR

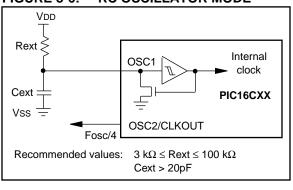
For timing insensitive applications the RC device option offers additional cost savings. The RC oscillator frequency is a function of the supply voltage, the resistor (Rext) values, capacitor (Cext) values, and the operating temperature. In addition to this, the oscillator frequency will vary from unit to unit due to normal process parameter variation. Furthermore. the difference in lead frame capacitance between package types also affects the oscillation frequency, especially for low Cext values. The user needs to take into account variation due to tolerance of the external R and C components. Figure 8-6 shows how an R/C combination is connected to the PIC16C84. For Rext values below 2.2 k $\Omega$ , the oscillator operation may become unstable, or stop completely. For very high Rext values (e.g.,  $1 M\Omega$ ), the oscillator becomes sensitive to noise, humidity and leakage. Thus, we recommend keeping Rext between 3 k $\Omega$  and 100 k $\Omega$ .

Although the oscillator will operate with no external capacitor (Cext = 0 pF), we recommend using values above 20 pF for noise and stability reasons. With little or no external capacitance, the oscillation frequency can vary dramatically due to changes in external capacitances, such as PCB trace capacitance or package lead frame capacitance.

See the electrical specification section for RC frequency variation from part to part due to normal process variation. The variation is larger for larger R (since leakage current variation will affect RC frequency more for large R) and for smaller C (since variation of input capacitance has a greater affect on RC frequency).

See the electrical specification section for variation of oscillator frequency due to VDD for given Rext/Cext values as well as frequency variation due to operating temperature.

The oscillator frequency, divided by 4, is available on the OSC2/CLKOUT pin, and can be used for test purposes or to synchronize other logic (see Figure 3-2 for waveform).



## FIGURE 8-6: RC OSCILLATOR MODE

**Note:** When the device oscillator is in RC mode, do not drive the OSC1 pin with an external clock or you may damage the device.

## 8.3 <u>Reset</u>

The PIC16C84 differentiates between various kinds of reset:

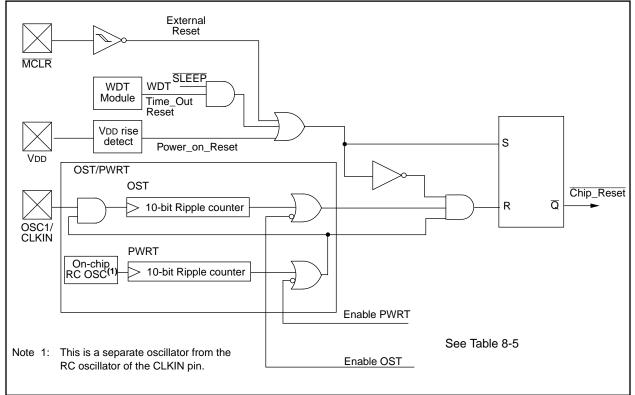
- Power-on Reset (POR)
- MCLR reset during normal operation
- MCLR reset during SLEEP
- WDT Reset (during normal operation)
- WDT Wake-up (during SLEEP)

Figure 8-7 shows a simplified block diagram of the onchip reset circuit. The electrical specifications state the pulse width requirements for the  $\overline{\text{MCLR}}$  pin. Some registers are not affected in any reset condition; their status is unknown on a POR reset and unchanged in any other reset. Most other registers are reset to a "reset state" on POR, MCLR or WDT reset during normal operation and on MCLR reset during SLEEP. They are not affected by a WDT reset during SLEEP, since this reset is viewed as the resumption of normal operation.

Table 8-3 gives a description of reset conditions for the program counter (PC) and the STATUS register. Table 8-4 gives a full description of reset states for all registers.

The  $\overline{\text{TO}}$  and  $\overline{\text{PD}}$  bits are set or cleared differently in different reset situations (Section 8.7). These bits are used in software to determine the nature of the reset.

#### FIGURE 8-7: SIMPLIFIED BLOCK DIAGRAM OF ON-CHIP RESET CIRCUIT



## TABLE 8-3 RESET CONDITION FOR PROGRAM COUNTER AND THE STATUS REGISTER

Condition	Program Counter	STATUS Register
Power-on Reset	000h	0001 1xxx
MCLR Reset during normal operation	000h	000u uuuu
MCLR Reset during SLEEP	000h	0001 0uuu
WDT Reset (during normal operation)	000h	0000 luuu
WDT Wake-up	PC + 1	սսս0 Օսսս
Interrupt wake-up from SLEEP	PC + 1 <sup>(1)</sup>	uuul 0uuu

Legend: u = unchanged, x = unknown.

Note 1: When the wake-up is due to an interrupt and the GIE bit is set, the PC is loaded with the interrupt vector (0004h).

Register	Address	Power-on Reset	MCLR Reset during: – normal operation – SLEEP WDT Reset during nor- mal operation	Wake-up from SLEEP: – through interrupt – through WDT time-out		
W	_	xxxx xxxx	นนนน นนนน	นนนน นนนน		
INDF	00h					
TMR0	01h	xxxx xxxx	นนนน นนนน	นนนน นนนน		
PCL	02h	0000h	0000h	PC + 1 <sup>(2)</sup>		
STATUS	03h	0001 1xxx	000q quuu <sup>(3)</sup>	uuuq quuu <sup>(3)</sup>		
FSR	04h	xxxx xxxx	นนนน นนนน	นนนน นนนน		
PORTA	05h	x xxxx	u uuuu	u uuuu		
PORTB	06h	xxxx xxxx	นนนน นนนน	นนนน นนนน		
EEDATA	08h	XXXX XXXX	นนนน นนนน	นนนน นนนน		
EEADR	09h	xxxx xxxx	นนนน นนนน	นนนน นนนน		
PCLATH	0Ah	0 0000	0 0000	u uuuu		
INTCON	0Bh	0000 000x	0000 000u	uuuu uuuu <sup>(1)</sup>		
INDF	80h					
OPTION_REG	81h	1111 1111	1111 1111	นนนน นนนน		
PCL	82h	0000h	0000h	PC + 1		
STATUS	83h	0001 1xxx	000q quuu <sup>(3)</sup>	uuuq quuu <sup>(3)</sup>		
FSR	84h	xxxx xxxx	นนนน นนนน	นนนน นนนน		
TRISA	85h	1 1111	1 1111	u uuuu		
TRISB	86h	1111 1111	1111 1111	นนนน นนนน		
EECON1	88h	0 x000	0 q000	0 uuuu		
EECON2	89h					
PCLATH	8Ah	0 0000	0 0000	u uuuu		
INTCON	8Bh	0000 000x	0000 000u uuuu uuuu <sup>(1)</sup>			

## TABLE 8-4 RESET CONDITIONS FOR ALL REGISTERS

Legend: u = unchanged, x = unknown, -= unimplemented bit read as '0',

q = value depends on condition.

Note 1: One or more bits in INTCON will be affected (to cause wake-up).

2: When the wake-up is due to an interrupt and the GIE bit is set, the PC is loaded with the interrupt vector (0004h).

3: Table 8-3 lists the reset value for each specific condition.

## 8.4 Power-on Reset (POR)

A Power-on Reset pulse is generated on-chip when VDD rise is detected (in the range of 1.2V - 1.7V). To take advantage of the POR, just tie the  $\overline{\text{MCLR}}$  pin directly (or through a resistor) to VDD. This will eliminate external RC components usually needed to create Power-on Reset. A minimum rise time for VDD must be met for this to operate properly. See Electrical Specifications for details.

When the device starts normal operation (exits the reset condition), device operating parameters (voltage, frequency, temperature, ...) must be meet to ensure operation. If these conditions are not met, the device must be held in reset until the operating conditions are met.

For additional information, refer to Application Note AN607, *Power-up Trouble Shooting*.

The POR circuit does not produce an internal reset when VDD declines.

### 8.5 <u>Power-up Timer (PWRT)</u>

The Power-up Timer (PWRT) provides a fixed 72 ms nominal time-out (TPWRT) from POR (Figure 8-9, Figure 8-10, Figure 8-11 and Figure 8-12). The Powerup Timer operates on an internal RC oscillator. The chip is kept in reset as long as the PWRT is active. The PWRT delay allows the VDD to rise to an acceptable level (Possible exception shown in Figure 8-12).

A configuration bit, PWRTE, can enable/disable the PWRT (Figure 8-1).

The power-up time delay TPWRT will vary from chip to chip due to VDD, temperature, and process variation. See DC parameters for details.

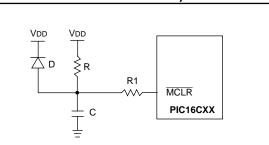
### 8.6 Oscillator Start-up Timer (OST)

The Oscillator Start-up Timer (OST) provides a 1024 oscillator cycle delay (from OSC1 input) after the PWRT delay ends (Figure 8-9, Figure 8-10, Figure 8-11 and Figure 8-12). This ensures the crystal oscillator or resonator has started and stabilized.

The OST time-out (TOST) is invoked only for XT, LP and HS modes and only on Power-on Reset or wake-up from SLEEP.

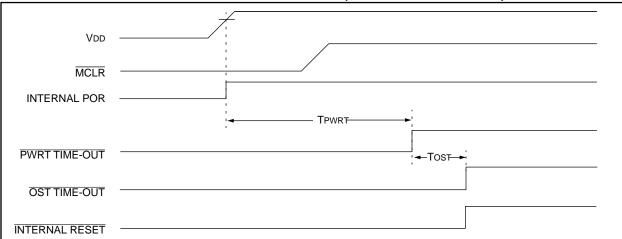
When VDD rises very slowly, it is possible that the TPWRT time-out and TOST time-out will expire before VDD has reached its final value. In this case (Figure 8-12), an external power-on reset circuit may be necessary (Figure 8-8).

#### FIGURE 8-8: EXTERNAL POWER-ON RESET CIRCUIT (FOR SLOW VDD POWER-UP)

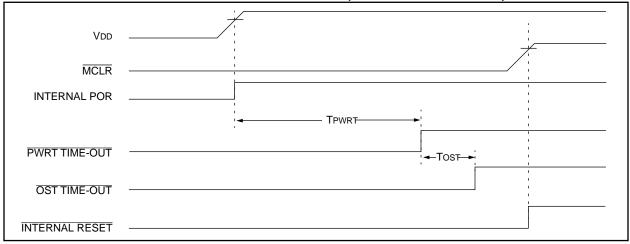


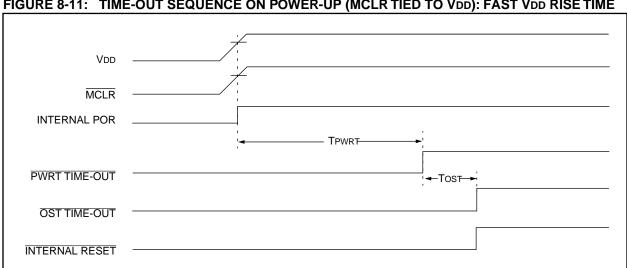
- Note 1: External Power-on Reset circuit is required only if VDD power-up rate is too slow. The diode D helps discharge the capacitor quickly when VDD powers down.
  - 2:  $R < 40 \text{ k}\Omega$  is recommended to make sure that voltage drop across R does not exceed 0.2V (max leakage current spec on MCLR pin is 5 µA). A larger voltage drop will degrade VIH level on the MCLR pin.
  - 3:  $R1 = 100\Omega$  to 1 k $\Omega$  will limit any current flowing into  $\overline{MCLR}$  from external capacitor C in the event of an  $\overline{MCLR}$  pin breakdown due to ESD or EOS.

FIGURE 8-9: TIME-OUT SEQUENCE ON POWER-UP (MCLR NOT TIED TO VDD): CASE 1



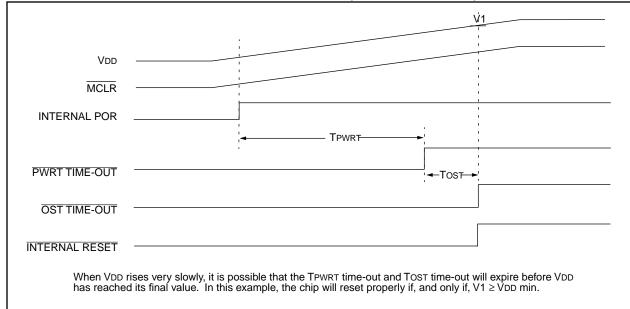
## FIGURE 8-10: TIME-OUT SEQUENCE ON POWER-UP (MCLR NOT TIED TO VDD): CASE 2





## FIGURE 8-11: TIME-OUT SEQUENCE ON POWER-UP (MCLR TIED TO VDD): FAST VDD RISE TIME

## FIGURE 8-12: TIME-OUT SEQUENCE ON POWER-UP (MCLR TIED TO VDD): SLOW VDD RISE TIME



### 8.7 <u>Time-out Sequence and Power Down</u> <u>Status Bits (TO/PD)</u>

On power-up (Figure 8-9, Figure 8-10, Figure 8-11 and Figure 8-12) the time-out sequence is as follows: First PWRT time-out is invoked after a POR has expired. Then the OST is activated. The total time-out will vary based on oscillator configuration and PWRTE configuration bit status. For example, in RC mode with the PWRT disabled, there will be no time-out at all.

## TABLE 8-5TIME-OUT IN VARIOUSSITUATIONS

Oscillator	Powe	Wake-up		
Configuration	PWRT Enabled	PWRT Disabled	from SLEEP	
XT, HS, LP	72 ms + 1024Tosc	1024Tosc	1024Tosc	
RC	72 ms	_	_	

Since the time-outs occur from the POR reset pulse, if  $\overline{\text{MCLR}}$  is kept low long enough, the time-outs will expire. Then bringing  $\overline{\text{MCLR}}$  high, execution will begin immediately (Figure 8-9). This is useful for testing purposes or to synchronize more than one PIC16CXX device when operating in parallel.

Table 8-6 shows the significance of the  $\overline{\text{TO}}$  and  $\overline{\text{PD}}$  bits. Table 8-3 lists the reset conditions for some special registers, while Table 8-4 lists the reset conditions for all the registers.

## TABLE 8-6STATUS BITS AND THEIRSIGNIFICANCE

то	PD	Condition
1	1	Power-on Reset
0	x	Illegal, TO is set on POR
x	0	Illegal, PD is set on POR
0	1	WDT Reset (during normal operation)
0	0	WDT Wake-up
1	1	MCLR Reset during normal operation
1	0	MCLR Reset during SLEEP or interrupt
		wake-up from SLEEP

## 8.8 Reset on Brown-Out

A brown-out is a condition where device power (VDD) dips below its minimum value, but not to zero, and then recovers. The device should be reset in the event of a brown-out.

To reset PIC16C84 devices when a brown-out occurs, external brown-out protection circuits may be built, as shown in Figure 8-13 and Figure 8-14.

<b>FIGURE 8-13:</b>	<b>BROWN-OUT PROTECTION</b>
	CIRCUIT 1

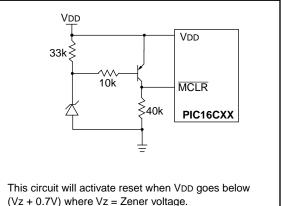
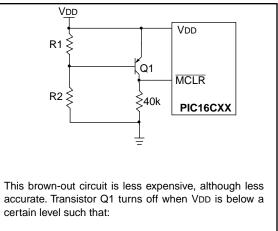


FIGURE 8-14: BROWN-OUT PROTECTION





$$V_{DD} \bullet \frac{R1}{R1 + R2} = 0.7V$$

### 8.9 Interrupts

The PIC16C84 has 4 sources of interrupt:

- External interrupt RB0/INT pin
- TMR0 overflow interrupt
- PORTB change interrupts (pins RB7:RB4)
- EEPROM write complete interrupt

The interrupt control register (INTCON) records individual interrupt requests in flag bits. It also contains the individual and global interrupt enable bits.

The global interrupt enable bit, GIE (INTCON<7>) enables (if set) all un-masked interrupts or disables (if cleared) all interrupts. Individual interrupts can be disabled through their corresponding enable bits in INTCON register. Bit GIE is cleared on reset.

The "return from interrupt" instruction, RETFIE, exits interrupt routine as well as sets the GIE bit, which reenable interrupts.

The RB0/INT pin interrupt, the RB port change interrupt and the TMR0 overflow interrupt flags are contained in the INTCON register. When an interrupt is responded to; the GIE bit is cleared to disable any further interrupt, the return address is pushed onto the stack and the PC is loaded with 0004h. For external interrupt events, such as the RB0/INT pin or PORTB change interrupt, the interrupt latency will be three to four instruction cycles. The exact latency depends when the interrupt event occurs (Figure 8-16). The latency is the same for both one and two cycle instructions. Once in the interrupt service routine the source(s) of the interrupt can be determined by polling the interrupt flag bits. The interrupt flag bit(s) must be cleared in software before re-enabling interrupts to avoid infinite interrupt requests.

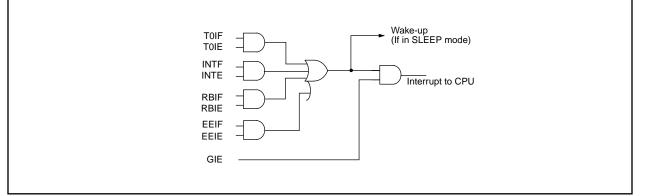
- Note 1: Individual interrupt flag bits are set regardless of the status of their corresponding mask bit or the GIE bit.
- Note 2: If an interrupt occurs while the Global Interrupt Enable (GIE) bit is being cleared, the GIE bit may unintentionally be re-enabled by the user's Interrupt Service Routine (the RETFIE instruction). The events that would cause this to occur are:
  - 1. An instruction clears the GIE bit while an interrupt is acknowledged
  - 2. The program branches to the Interrupt vector and executes the Interrupt Service Routine.
  - 3. The Interrupt Service Routine completes with the execution of the RETFIE instruction. This causes the GIE bit to be set (enables interrupts), and the program returns to the instruction after the one which was meant to disable interrupts.

The method to ensure that interrupts are globally disabled is:

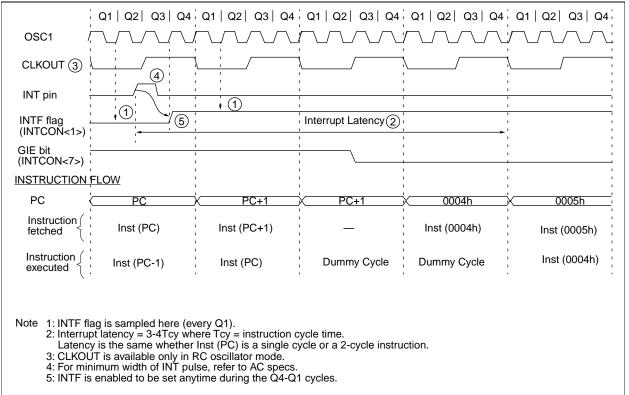
 Ensure that the GIE bit is cleared by the instruction, as shown in the following code:

LOOP	BCF	INTCON,GIE	;Disable All
			; Interrupts
	BTFSC	INTCON,GIE	;All Interrupts
			; Disabled?
	GOTO	LOOP	;NO, try again
			; Yes, continue
			; with program
			; flow

## FIGURE 8-15: INTERRUPT LOGIC



#### FIGURE 8-16: INT PIN INTERRUPT TIMING



#### 8.9.1 INT INTERRUPT

External interrupt on RB0/INT pin is edge triggered: either rising if INTEDG bit (OPTION\_REG<6>) is set, or falling, if INTEDG bit is clear. When a valid edge appears on the RB0/INT pin, the INTF bit (INTCON<1>) is set. This interrupt can be disabled by clearing control bit INTE (INTCON<4>). Flag bit INTF must be cleared in software via the interrupt service routine before re-enabling this interrupt. The INT interrupt can wake the processor from SLEEP (Section 8.12) only if the INTE bit was set prior to going into SLEEP. The status of the GIE bit decides whether the processor branches to the interrupt vector following wake-up.

#### 8.9.2 TMR0 INTERRUPT

An overflow (FFh  $\rightarrow$  00h) in TMR0 will set flag bit T0IF (INTCON<2>). The interrupt can be enabled/disabled by setting/clearing enable bit T0IE (INTCON<5>) (Section 6.0).

#### 8.9.3 PORT RB INTERRUPT

An input change on PORTB<7:4> sets flag bit RBIF (INTCON<0>). The interrupt can be enabled/disabled by setting/clearing enable bit RBIE (INTCON<3>) (Section 5.2).

**Note 1:** If a change on an I/O pin should occur when a read operation of PORTB is being executed (start of the Q2 cycle), the RBIF interrupt flag bit may not get set.

### EXAMPLE 8-1: SAVING STATUS AND W REGISTERS IN RAM

PUSH	MOVWF	W_TEMP	; Copy W to TEMP register,
	SWAPF	STATUS, W	; Swap status to be saved into W
	MOVWF	STATUS_TEMP	; Save status to STATUS_TEMP register
ISR	:		:
	:		; Interrupt Service Routine
	:		; should configure Bank as required
	:		;
POP	SWAPF	STATUS_TEMP, W	; Swap nibbles in STATUS_TEMP register
			; and place result into W
	MOVWF	STATUS	; Move W into STATUS register
			; (sets bank to original state)
	SWAPF	W_TEMP, F	; Swap nibbles in W_TEMP and place result in W_TEMP
	SWAPF	W_TEMP, W	; Swap nibbles in W_TEMP and place result into W

### 8.10 <u>Context Saving During Interrupts</u>

During an interrupt, only the return PC value is saved on the stack. Typically, users wish to save key register values during an interrupt (e.g., W register and STATUS register). This is implemented in software.

Example 8-1 stores and restores the STATUS and W register's values. The User defined registers, W\_TEMP and STATUS\_TEMP are the temporary storage locations for the W and STATUS registers values.

Example 8-1 does the following:

- a) Stores the W register.
- b) Stores the STATUS register in STATUS\_TEMP.
- c) Executes the Interrupt Service Routine code.
- d) Restores the STATUS (and bank select bit) register.
- e) Restores the W register.

### 8.11 Watchdog Timer (WDT)

The Watchdog Timer is a free running on-chip RC oscillator which does not require any external components. This RC oscillator is separate from the RC oscillator of the OSC1/CLKIN pin. That means that the WDT will run even if the clock on the OSC1/CLKIN and OSC2/CLKOUT pins of the device has been stopped, for example, by execution of a SLEEP instruction. During normal operation a WDT time-out generates a device RESET. If the device is in SLEEP mode, a WDT Wake-up causes the device to wake-up and continue with normal operation. The WDT can be permanently disabled by programming configuration bit WDTE as a '0' (Section 8.1).

#### 8.11.1 WDT PERIOD

The WDT has a nominal time-out period of 18 ms, (with no prescaler). The time-out periods vary with temperature, VDD and process variations from part to

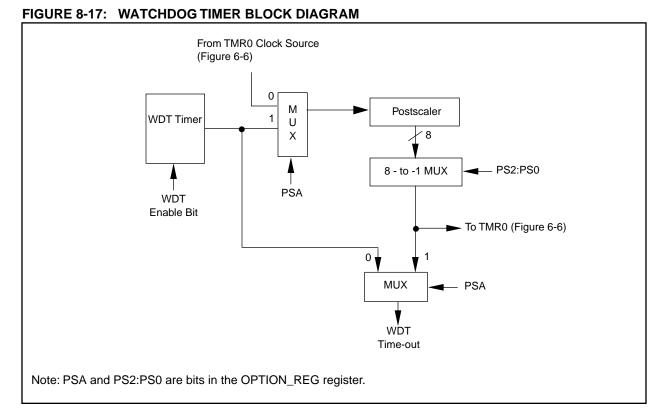
part (see DC specs). If longer time-out periods are desired, a prescaler with a division ratio of up to 1:128 can be assigned to the WDT under software control by writing to the OPTION\_REG register. Thus, time-out periods up to 2.3 seconds can be realized.

The CLRWDT and SLEEP instructions clear the WDT and the postscaler (if assigned to the WDT) and prevent it from timing out and generating a device RESET condition.

The  $\overline{\text{TO}}$  bit in the STATUS register will be cleared upon a WDT time-out.

8.11.2 WDT PROGRAMMING CONSIDERATIONS

It should also be taken into account that under worst case conditions (VDD = Min., Temperature = Max., max. WDT prescaler) it may take several seconds before a WDT time-out occurs.



## TABLE 8-7 SUMMARY OF REGISTERS ASSOCIATED WITH THE WATCHDOG TIMER

Address	Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on Power-on Reset	Value on all other resets
2007h	Config. bits	_	—	—	CP	PWRTE	WDTE	FOSC1	FOSC0		
81h	OPTION_ REG	RBPU	INTEDG	TOCS	TOSE	PSA	PS2	PS1	PS0	1111 1111	1111 1111

Legend: x = unknown. Shaded cells are not used by the WDT.

### 8.12 Power-down Mode (SLEEP)

A device may be powered down (SLEEP) and later powered up (Wake-up from SLEEP).

#### 8.12.1 SLEEP

The Power-down mode is entered by executing the  $\ensuremath{\mathtt{SLEEP}}$  instruction.

If enabled, the Watchdog Timer is cleared (but keeps running), the  $\overline{PD}$  bit (STATUS<3>) is cleared, the  $\overline{TO}$  bit (STATUS<4>) is set, and the oscillator driver is turned off. The I/O ports maintain the status they had before the SLEEP instruction was executed (driving high, low, or hi-impedance).

For the lowest current consumption in SLEEP mode, place all I/O pins at either at VDD or VSS, with no external circuitry drawing current from the I/O pins, and disable external clocks. I/O pins that are hi-impedance inputs should be pulled high or low externally to avoid switching currents caused by floating inputs. The TOCKI input should also be at VDD or VSS. The contribution from on-chip pull-ups on PORTB should be considered.

The  $\overline{\text{MCLR}}$  pin must be at a logic high level (VIHMC).

It should be noted that a RESET generated by a WDT time-out does not drive the  $\overline{\text{MCLR}}$  pin low.

#### 8.12.2 WAKE-UP FROM SLEEP

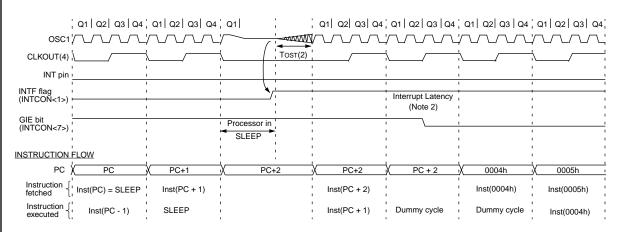
The device can wake-up from SLEEP through one of the following events:

- 1. External reset input on MCLR pin.
- 2. WDT Wake-up (if WDT was enabled).
- 3. Interrupt from RB0/INT pin, RB port change, or data EEPROM write complete.

Peripherals cannot generate interrupts during SLEEP, since no on-chip Q clocks are present.

The first event ( $\overline{\text{MCLR}}$  reset) will cause a device reset. The two latter events are considered a continuation of program execution. The  $\overline{\text{TO}}$  and  $\overline{\text{PD}}$  bits can be used to determine the cause of a device reset. The  $\overline{\text{PD}}$  bit, which is set on power-up, is cleared when SLEEP is invoked. The  $\overline{\text{TO}}$  bit is cleared if a WDT time-out occurred (and caused wake-up).

While the SLEEP instruction is being executed, the next instruction (PC + 1) is pre-fetched. For the device to wake-up through an interrupt event, the corresponding interrupt enable bit must be set (enabled). Wake-up occurs regardless of the state of the GIE bit. If the GIE bit is clear (disabled), the device continues execution at the instruction after the SLEEP instruction. If the GIE bit is set (enabled), the device executes the instruction after the SLEEP instruction after the sLEEP instruction after the subject (0004h). In cases where the execution of the instruction following SLEEP is not desirable, the user should have a NOP after the SLEEP instruction.



## FIGURE 8-18: WAKE-UP FROM SLEEP THROUGH INTERRUPT

Note 1: XT, HS or LP oscillator mode assumed.

2: TOST = 1024TOSC (drawing not to scale) This delay will not be there for RC osc mode.

3: GIE = '1' assumed. In this case after wake- up, the processor jumps to the interrupt routine. If GIE = '0', execution will continue in-line.
4: CLKOUT is not available in these osc modes, but shown here for timing reference.

#### 8.12.3 WAKE-UP USING INTERRUPTS

When global interrupts are disabled (GIE cleared) and any interrupt source has both its interrupt enable bit and interrupt flag bit set, one of the following will occur:

- If the interrupt occurs before the execution of a SLEEP instruction, the SLEEP instruction will complete as a NOP. Therefore, the WDT and WDT postscaler will not be cleared, the TO bit will not be set and PD bits will not be cleared.
- If the interrupt occurs during or after the execution of a SLEEP instruction, the device will immediately wake up from sleep . The SLEEP instruction will be completely executed before the wake-up. Therefore, the WDT and WDT postscaler will be cleared, the TO bit will be set and the PD bit will be cleared.

Even if the flag bits were checked before executing a SLEEP instruction, it may be possible for flag bits to become set before the SLEEP instruction completes. To determine whether a SLEEP instruction executed, test the  $\overline{PD}$  bit. If the  $\overline{PD}$  bit is set, the SLEEP instruction was executed as a NOP.

To ensure that the WDT is cleared, a CLRWDT instruction should be executed before a SLEEP instruction.

## 8.13 Program Verification/Code Protection

If the code protection bit(s) have not been programmed, the on-chip program memory can be read out for verification purposes.

Note:	Microchip does not recommend code pro-
	tecting windowed devices

### 8.14 ID Locations

Four memory locations (2000h - 2003h) are designated as ID locations to store checksum or other code identification numbers. These locations are not accessible during normal execution but are readable and writable only during program/verify. Only the 4 least significant bits of ID location are usable.

For ROM devices, these values are submitted along with the ROM code.

## 8.15 In-Circuit Serial Programming

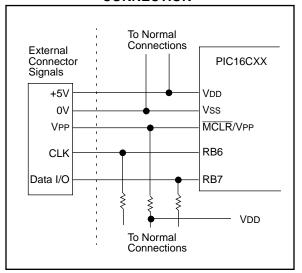
PIC16C84 microcontrollers can be serially programmed while in the end application circuit. This is simply done with two lines for clock and data, and three other lines for power, ground, and the programming voltage. Customers can manufacture boards with unprogrammed devices, and then program the microcontroller just before shipping the product, allowing the most recent firmware or custom firmware to be programmed.

The device is placed into a program/verify mode by holding the RB6 and RB7 pins low, while raising the MCLR pin from VIL to VIHH (see *PIC16C84 EEPROM Memory Programming Specification* (DS30189)). RB6 becomes the programming clock and RB7 becomes the programming data. Both RB6 and RB7 are Schmitt Trigger inputs in this mode.

After reset, to place the device into programming/verify mode, the program counter (PC) points to location 00h. A 6-bit command is then supplied to the device, 14-bits of program data is then supplied to or from the device, using load or read-type instructions. For complete details of serial programming, please refer to the *In-Circuit Serial Programming Guide* (DS30277).

For ROM devices, both the program memory and Data EEPROM memory may be read, but only the Data EEPROM memory may be programmed.

#### FIGURE 8-19: TYPICAL IN-SYSTEM SERIAL PROGRAMMING CONNECTION



NOTES:

## 9.0 INSTRUCTION SET SUMMARY

Each PIC16CXX instruction is a 14-bit word divided into an OPCODE which specifies the instruction type and one or more operands which further specify the operation of the instruction. The PIC16CXX instruction set summary in Table 9-2 lists **byte-oriented**, **bit-oriented**, and **literal and control** operations. Table 9-1 shows the opcode field descriptions.

For **byte-oriented** instructions, 'f' represents a file register designator and 'd' represents a destination designator. The file register designator specifies which file register is to be used by the instruction.

The destination designator specifies where the result of the operation is to be placed. If 'd' is zero, the result is placed in the W register. If 'd' is one, the result is placed in the file register specified in the instruction.

For **bit-oriented** instructions, 'b' represents a bit field designator which selects the number of the bit affected by the operation, while 'f' represents the number of the file in which the bit is located.

For **literal and control** operations, 'k' represents an eight or eleven bit constant or literal value.

## TABLE 9-1OPCODE FIELDDESCRIPTIONS

Field	Description
f	Register file address (0x00 to 0x7F)
W	Working register (accumulator)
b	Bit address within an 8-bit file register
k	Literal field, constant data or label
x	Don't care location (= 0 or 1) The assembler will generate code with $x = 0$ . It is the recommended form of use for compatibility with all Microchip software tools.
d	Destination select; d = 0: store result in W, d = 1: store result in file register f. Default is d = 1
label	Label name
TOS	Top of Stack
PC	Program Counter
PCLATH	Program Counter High Latch
GIE	Global Interrupt Enable bit
WDT	Watchdog Timer/Counter
TO	Time-out bit
PD	Power-down bit
dest	Destination either the W register or the specified register file location
[]	Options
()	Contents
$\rightarrow$	Assigned to
<>	Register bit field
E	In the set of
italics	User defined term (font is courier)

The instruction set is highly orthogonal and is grouped into three basic categories:

- Byte-oriented operations
- Bit-oriented operations
- Literal and control operations

All instructions are executed within one single instruction cycle, unless a conditional test is true or the program counter is changed as a result of an instruction. In this case, the execution takes two instruction cycles with the second cycle executed as a NOP. One instruction cycle consists of four oscillator periods. Thus, for an oscillator frequency of 4 MHz, the normal instruction execution time is 1  $\mu$ s. If a conditional test is true or the program counter is changed as a result of an instruction, the instruction execution time is 2  $\mu$ s.

Table 9-2 lists the instructions recognized by the MPASM assembler.

Figure 9-1 shows the general formats that the instructions can have.

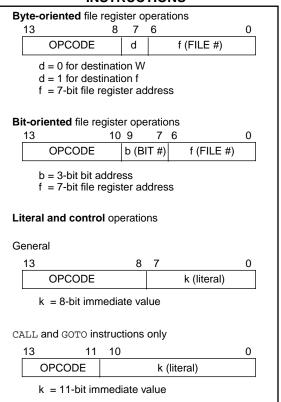
**Note:** To maintain upward compatibility with future PIC16CXX products, <u>do not use</u> the OPTION and TRIS instructions.

All examples use the following format to represent a hexadecimal number:

0xhh

where h signifies a hexadecimal digit.

## FIGURE 9-1: GENERAL FORMAT FOR INSTRUCTIONS



Mnemonic,		Description	Cycles		14-Bit	Opcode	Status	Notes	
Operan	ds			MSb			LSb	Affected	
		BYTE-ORIENTED FILE REGIS	TER OPE	RATIC	ONS				
ADDWF	f, d	Add W and f	1	00	0111	dfff	ffff	C,DC,Z	1,2
ANDWF	f, d	AND W with f	1	00	0101	dfff	ffff	Z	1,2
CLRF	f	Clear f	1	00	0001	lfff	ffff	Z	2
CLRW	-	Clear W	1	00	0001	0xxx	xxxx	Z	
COMF	f, d	Complement f	1	00	1001	dfff	ffff	Z	1,2
DECF	f, d	Decrement f	1	00	0011	dfff	ffff	Z	1,2
DECFSZ	f, d	Decrement f, Skip if 0	1(2)	00	1011	dfff	ffff		1,2,3
INCF	f, d	Increment f	1	00	1010	dfff	ffff	Z	1,2
INCFSZ	f, d	Increment f, Skip if 0	1(2)	00	1111	dfff	ffff		1,2,3
IORWF	f, d	Inclusive OR W with f	1	00	0100	dfff	ffff	Z	1,2
MOVF	f, d	Move f	1	00	1000	dfff	ffff	Z	1,2
MOVWF	f	Move W to f	1	00	0000	lfff	ffff		
NOP	-	No Operation	1	00	0000	0xx0	0000		
RLF	f, d	Rotate Left f through Carry	1	00	1101	dfff	ffff	С	1,2
RRF	f, d	Rotate Right f through Carry	1	00	1100	dfff	ffff	С	1,2
SUBWF	f, d	Subtract W from f	1	00	0010	dfff	ffff	C,DC,Z	1,2
SWAPF	f, d	Swap nibbles in f	1	00	1110	dfff	ffff	-, -,	1,2
XORWF	f, d	Exclusive OR W with f	1	00	0110	dfff	ffff	Z	1,2
		BIT-ORIENTED FILE REGIST		RATIO	NS				
BCF	f, b	Bit Clear f	1	01	00bb	bfff	ffff		1,2
BSF	f, b	Bit Set f	1	01	01bb	bfff	ffff		1,2
BTFSC	f, b	Bit Test f, Skip if Clear	1 (2)	01	10bb	bfff	ffff		3
BTFSS	f, b	Bit Test f, Skip if Set	1 (2)	01	11bb	bfff	ffff		3
		LITERAL AND CONTROL	OPERAT	IONS					
ADDLW	k	Add literal and W	1	11	111x	kkkk	kkkk	C,DC,Z	
ANDLW	k	AND literal with W	1	11	1001	kkkk	kkkk	Z	
CALL	k	Call subroutine	2	10	0kkk	kkkk	kkkk		
CLRWDT	-	Clear Watchdog Timer	1	00	0000	0110	0100	TO,PD	
GOTO	k	Go to address	2	10	1kkk	kkkk	kkkk		
IORLW	k	Inclusive OR literal with W	1	11	1000	kkkk	kkkk	Z	
MOVLW	k	Move literal to W	1	11	00xx	kkkk	kkkk		
RETFIE	-	Return from interrupt	2	00	0000	0000	1001		
RETLW	k	Return with literal in W	2	11	01xx	kkkk	kkkk		
RETURN	-	Return from Subroutine	2	00	0000	0000	1000		
SLEEP	-	Go into standby mode	1	00	0000	0110	0011	TO,PD	
SUBLW	k	Subtract W from literal	1	11		kkkk		C,DC,Z	
	k	Exclusive OR literal with W	1					Z	

### TABLE 9-2 PIC16CXX INSTRUCTION SET

Note 1: When an I/O register is modified as a function of itself (e.g., MOVF PORTB, 1), the value used will be that value present on the pins themselves. For example, if the data latch is '1' for a pin configured as input and is driven low by an external device, the data will be written back with a '0'.

2: If this instruction is executed on the TMR0 register (and, where applicable, d = 1), the prescaler will be cleared if assigned to the Timer0 Module.

3: If Program Counter (PC) is modified or a conditional test is true, the instruction requires two cycles. The second cycle is executed as a NOP.

## 9.1 Instruction Descriptions

ADDLW	Add Lite	ral and \	N				
Syntax:	[ <i>label</i> ] A	DDLW	k				
Operands:	$0 \le k \le 255$						
Operation:	$(W) + k \rightarrow (W)$						
Status Affected:	C, DC, Z						
Encoding:	11 111x kkkk kkkk						
Description:	The contents of the W register are added to the eight bit literal 'k' and the result is placed in the W register.						
Words:	1						
Cycles:	1						
Q Cycle Activity:	Q1	Q2	Q3	Q4			
	Decode	Read literal 'k'	Process data	Write to W			
Example:	ADDLW	0x15					
	Before In	W =	0x10				
	/	W =	0x25				

ADDWF	Add W a	nd f		
Syntax:	[ <i>label</i> ] A	DDWF	f,d	
Operands:	$\begin{array}{l} 0 \leq f \leq 12 \\ d \in \left[0,1\right] \end{array}$	.7		
Operation:	(W) + (f)	ightarrow (desti	nation)	
Status Affected:	C, DC, Z			
Encoding:	00	0111	dfff	ffff
Description:	Add the co register 'f'. in the W re stored bac	If 'd' is 0 egister. If '	the result i d' is 1 the	s stored
Words:	1			
Cycles:	1			
Q Cycle Activity:	Q1	Q2	Q3	Q4
	Decode	Read register 'f'	Process data	Write to destination
Example	ADDWF	FSR,	0	
	Before In	struction	1	
		W = FSR =	0x17 0xC2	
	After Inst		0x02	
		W =	0xD9	
		FSR =	0xC2	

AND'ed w	ith the eig	ht bit litera	l 'k'. The	
1				
1				
Q1	Q2	Q3	Q4	
Decode	Read literal "k"	Process data	Write to W	
ANDLW	0x5F			
		1		
		0xA3		
		0x03		
AND W v				
ANDWv [ <i>label</i> ] A		f,d		
[ <i>label</i> ] A 0 ≤ f ≤ 12	NDWF	f,d		
[ <i>label</i> ] A $0 \le f \le 12$ $d \in [0,1]$	NDWF 27		n)	
[ <i>label</i> ] A 0 ≤ f ≤ 12 d ∈ [0,1] (W) .ANE	NDWF 27		n)	
[ <i>label</i> ] A 0 ≤ f ≤ 12 d ∈ [0,1] (W) .ANE Z	NDWF 27 0. (f) → (d	destinatio		
[ <i>label</i> ] A $0 \le f \le 12$ $d \in [0,1]$ (W) .ANE Z 00	NDWF 27 D. (f) $\rightarrow$ (c) 0101	destinatio	ffff	
[ <i>label</i> ] A $0 \le f \le 12$ $d \in [0,1]$ (W) .ANE Z 00 AND the V is 0 the re- ter. If 'd' is	NDWF 27 D. (f) $\rightarrow$ (c 0101 V register sult is stor 1 the resi	destinatio	ffff er 'f'. If 'd' <i>N</i> regis-	
[ <i>label</i> ] A $0 \le f \le 12$ $d \in [0,1]$ (W) .ANE Z 00 AND the V is 0 the read	NDWF 27 D. (f) $\rightarrow$ (c 0101 V register sult is stor 1 the resi	destinatio	ffff er 'f'. If 'd' <i>N</i> regis-	
[ <i>label</i> ] A $0 \le f \le 12$ $d \in [0,1]$ (W) .ANE Z 00 AND the V is 0 the re- ter. If 'd' is register 'f'.	NDWF 27 D. (f) $\rightarrow$ (c 0101 V register sult is stor 1 the resi	destinatio	ffff er 'f'. If 'd' <i>N</i> regis-	
[ <i>label</i> ] A $0 \le f \le 12$ $d \in [0,1]$ (W) .ANE Z 00 AND the V is 0 the re- ter. If 'd' is register 'f'. 1	NDWF 27 D. (f) $\rightarrow$ (c 0101 V register sult is stor 1 the resi	destinatio	ffff er 'f'. If 'd' <i>N</i> regis-	
[ <i>label</i> ] A $0 \le f \le 12$ $d \in [0,1]$ (W) .ANE Z 00 AND the V is 0 the re- ter. If 'd' is register 'f'. 1 1	NDWF 27 D. (f) $\rightarrow$ (c 0101 V register sult is sto 1 the rest Q2 Read	destinatio	ffff er 'f'. If 'd' W regis- d back in Q4 Write to	
[ <i>label</i> ] A $0 \le f \le 12$ $d \in [0,1]$ (W) .ANE Z 00 AND the V is 0 the re- ter. If 'd' is register 'f'. 1 1 Q1	NDWF 2.7 D. (f) $\rightarrow$ (c 0101 V register sult is sto 1 the rest Q2	destinatio dfff with regist red in the V ult is store Q3	ffff er 'f'. If 'd' W regis- d back in Q4 Write to	
$\begin{bmatrix} label \end{bmatrix} A$ $0 \le f \le 12$ $d \in [0,1]$ $(W) .ANE$ $Z$ $\boxed{00}$ AND the V is 0 the reter. If 'd' is register 'f'. 1 1 Q1 $\boxed{Decode}$	NDWF 27 D. (f) $\rightarrow$ (c 0101 V register sult is stor 1 the rest Q2 Read register 'f'	destinatio	ffff er 'f'. If 'd' W regis- d back in Q4 Write to	
$[label] A$ $0 \le f \le 12$ $d \in [0,1]$ $(W) .ANE$ $Z$ $\boxed{00}$ AND the V is 0 the re- ter. If 'd' is register 'f'. 1 1 Q1 $\boxed{Decode}$ ANDWF	NDWF 27 27 2. (f) $\rightarrow$ (d 0101 V register sult is stor 1 the result Q2 Read register 'f' FSR,	destinatio	ffff er 'f'. If 'd' W regis- d back in Q4 Write to	
$\begin{bmatrix} label \end{bmatrix} A$ $0 \le f \le 12$ $d \in [0,1]$ $(W) .ANE$ $Z$ $\boxed{00}$ AND the V is 0 the reter. If 'd' is register 'f'. 1 1 Q1 $\boxed{Decode}$	NDWF 27 27 2. (f) $\rightarrow$ (d 0101 V register sult is stor 1 the result Q2 Read register 'f' FSR,	destinatio	ffff er 'f'. If 'd' W regis- d back in Q4 Write to	
$[label] A$ $0 \le f \le 12$ $d \in [0,1]$ $(W) .ANE$ $Z$ $\boxed{00}$ AND the V is 0 the re- ter. If 'd' is register 'f'. 1 1 Q1 $\boxed{Decode}$ ANDWF Before In	NDWF 27 D. (f) $\rightarrow$ (d 0101 V register sult is ston 1 the result Q2 Read register 'f' FSR, struction W = FSR =	destinatio	ffff er 'f'. If 'd' W regis- d back in Q4	
	[ <i>label</i> ] A $0 \le k \le 24$ (W) .ANE Z 11 The conte AND'ed w result is pl 1 1 Q1 Decode ANDLW Before In After Inst	[ <i>label</i> ] ANDLW $0 \le k \le 255$ (W) .AND. (k) $\rightarrow$ ( Z 11 1001 The contents of W r AND'ed with the eig result is placed in th 1 2 Q1 Q2 Decode Read literal "k" ANDLW 0x5F Before Instruction W = After Instruction	$\begin{array}{c} 0 \leq k \leq 255 \\ (W) \ .AND. \ (k) \rightarrow (W) \\ Z \\ \hline 11 \ 1001 \ kkkk \\ \hline The contents of W register are \\ AND'ed with the eight bit literal result is placed in the W register are \\ AND'ed with the eight bit literal \\ 1 \\ 1 \\ \hline 2 \\ Q1 \ Q2 \ Q3 \\ \hline \hline Decode \ Read \\ literal \ 'k'' \ Process \\ data \\ \hline \hline ANDLW \ 0x5F \\ \hline Before \ Instruction \\ W = \ 0xA3 \\ \hline After \ Instruction \\ \hline \end{array}$	

W =

FSR =

0x17

0x02

BCF	Bit Clear	f			
Syntax:	[ <i>label</i> ] BC	CF f,b			
Operands:	$\begin{array}{l} 0 \leq f \leq 127 \\ 0 \leq b \leq 7 \end{array}$				
Operation:	$0 \rightarrow (f < b;$	>)			
Status Affected:	None				
Encoding:	01	00bb	bfff	ffff	
Description:	Bit 'b' in register 'f' is cleared.				
Words:	1				
Cycles:	1				
Q Cycle Activity:	Q1	Q2	Q3	Q4	
	Decode	Read register 'f'	Process data	Write register 'f'	
Example	BCF	FLAG_	REG, 7		
	Before Instruction FLAG_REG = 0xC7				
	After Inst		EG = 0x47		

BTFSC	Bit Test,	Skip if Cl	ear			
Syntax:	[ <i>label</i> ] BT	FSC f,b				
Operands:	$0 \le f \le 12$ $0 \le b \le 7$	$0 \le f \le 127$ $0 \le b \le 7$				
Operation:	skip if (f<	b>) = 0				
Status Affected:	None					
Encoding:	01	10bb	bfff	ffff		
Description:	If bit 'b' in register 'f' is '1' then the next instruction is executed. If bit 'b', in register 'f', is '0' then the next instruction is discarded, and a NOP is executed instead, making this a 2Tcy instruction.					
Words:	1					
Cycles:	1(2)					
Q Cycle Activity:	Q1	Q2	Q3	Q4		
	Decode	Read register 'f'	Process data	No-Operat ion		
If Skip:	(2nd Cyc	le)				
	Q1	Q2	Q3	Q4		
	No-Operat ion	No-Operati on	No-Opera tion	No-Operat ion		
Example	HERE FALSE TRUE		FLAG,1 PROCESS_	_CODE		
	Before In	struction				
			ddress H	ERE		
	After Inst	ruction if FLAG<1>	= 0,			
			address T	RUE		

BSF	Bit Set f
Syntax:	[ <i>label</i> ] BSF f,b
Operands:	$\begin{array}{l} 0 \leq f \leq 127 \\ 0 \leq b \leq 7 \end{array}$
Operation:	$1 \rightarrow (f < b >)$
Status Affected:	None
Encoding:	01 01bb bfff fff
Description:	Bit 'b' in register 'f' is set.
Words:	1
Cycles:	1
Q Cycle Activity:	Q1 Q2 Q3 Q4
	Decode Read register 'f' 'data 'register 'f'
Example	BSF FLAG_REG, 7
	Before Instruction FLAG_REG = 0x0A After Instruction FLAG_REG = 0x8A

if FLAG<1>=1, PC = addre

address FALSE

BTFSS	Bit Test f, Skip if Set	CALL	Call Subroutine
Syntax:	[ <i>label</i> ] BTFSS f,b	Syntax:	[ <i>label</i> ] CALL k
Operands:	$0 \le f \le 127$	Operands:	$0 \le k \le 2047$
Operation:	0 ≤ b < 7 skip if (f <b>) = 1</b>	Operation:	$(PC)$ + 1 $\rightarrow$ TOS, k $\rightarrow$ PC<10:0>,
Status Affected:	None		$(PCLATH{<}4{:}3{>}) \rightarrow PC{<}12{:}11{>}$
Encoding:	01 11bb bfff ffff	Status Affected:	None
Description:	If bit 'b' in register 'f' is '0' then the next	Encoding:	10 0kkk kkkk kkkk
Words:	instruction is executed. If bit 'b' is '1', then the next instruction is discarded and a NOP is executed instead, making this a 2Tcy instruction.	Description:	Call Subroutine. First, return address (PC+1) is pushed onto the stack. The eleven bit immediate address is loaded into PC bits <10:0>. The upper bits of the PC are loaded from PCLATH. CALL is a two cycle instruction.
Cycles:	1(2)	Words:	1
Q Cycle Activity:	Q1 Q2 Q3 Q4	Cycles:	2
	Decode Read Process No-Operat register 'f' data ion	Q Cycle Activity:	Q1 Q2 Q3 Q4
If Skip:	(2nd Cycle)	1st Cycle	literal 'k', data PC
	Q1 Q2 Q3 Q4		Push PC to Stack
	No-Operati ion No-Operati No-Operat ion No-Operati	2nd Cycle	No-Opera No-Opera No-Operat tion tion
Example	HERE BTFSC FLAG,1 FALSE GOTO PROCESS CODE	Example	HERE CALL THERE
	TRUE •		Before Instruction
	•		PC = Address HERE After Instruction
	Before Instruction		PC = Address THERE
	PC = address HERE After Instruction if FLAG<1> = 0, PC = address FALSE if FLAG<1> = 1, PC = address TRUE		TOS = Address HERE+1

CLRF	Clear f				
Syntax:	[ <i>label</i> ] CLRF f				
Operands:	$0 \le f \le 12$	27			
Operation:	$\begin{array}{l} 00h \rightarrow (f) \\ 1 \rightarrow Z \end{array}$				
Status Affected:	Z				
Encoding:	00	0001	lfff	ffff	
Description:	The contents of register 'f' are cleared and the Z bit is set.				
Words:	1				
Cycles:	1				
Q Cycle Activity:	Q1	Q2	Q3	Q4	
	Decode	Read register 'f'	Process data	Write register 'f'	
Example	CLRF	FLAG	G_REG		
	Before In After Inst	FLAG_RE		0x5A	
		FLAG_RE Z	EG = =	0x00 1	

CLRW	Clear W					
Syntax:	[ label ]	CLRW				
Operands:	None					
Operation:	$\begin{array}{l} 00h \rightarrow (V \\ 1 \rightarrow Z \end{array}$	$\begin{array}{l} 00h \rightarrow (W) \\ 1 \rightarrow Z \end{array}$				
Status Affected:	Z					
Encoding:	00	0001	0xxx	xxxx		
Description:	W register set.	is cleared	. Zero bit (	Z) is		
Words:	1					
Cycles:	1					
Q Cycle Activity:	Q1	Q2	Q3	Q4		
	Decode	No-Opera tion	Process data	Write to W		
Example	CLRW					
	Before In					
	After Inst		0x5A			
			0x00			
		Z =	1			
CLRWDT	Clear Wa	atchdog 1	imer			
Syntax:	[ label ]	CLRWD	-			
Operands:	None					
Operation:	$00h \rightarrow W$					
	$0 \rightarrow WD^{-1}$ $1 \rightarrow \overline{TO}$	F prescale	ər,			
	$1 \rightarrow \overline{PD}$					
Status Affected:	TO, PD					
Encoding:	00	0000	0110	0100		
Description:		struction r				
		. It also res T. Status b				
Words:	1					
Cycles:	1					
Q Cycle Activity:	Q1	Q2	Q3	Q4		
	Decode	No-Opera tion	Process data	Clear WDT Counter		
	L					
Example	CLRWDT					
Example	Before In		4	2		
Example	Before In	WDT cour	iter =	?		
Example	Before In After Inst	WDT cour ruction WDT cour	iter =	? 0x00		
Example	Before In After Inst	WDT cour ruction WDT cour WDT pres	ter = caler=	0x00 0		
Example	Before In After Inst	WDT cour ruction WDT cour	ter = caler= =	0x00		
Example	Before In After Inst	WDT cour ruction WDT cour WDT pres TO	ter = caler= =	0x00 0 1		

Complement f	DECFSZ	Decrement f, Skip if 0
[ <i>label</i> ] COMF f,d	Syntax:	[label] DECFSZ f,d
$\begin{array}{l} 0 \leq f \leq 127 \\ d \in \left[0,1\right] \end{array}$	Operands:	0 ≤ f ≤ 127 d ∈ [0,1]
$(\overline{f}) \rightarrow$ (destination)	Operation:	(f) - 1 $\rightarrow$ (destination);
Z		skip if result = 0
00 1001 dfff ffff	Status Affected:	None
The contents of register 'f' are comple-	Encoding:	00 1011 dfff ffff
W. If 'd' is 1 the result is stored back in	Description:	The contents of register 'f' are decre- mented. If 'd' is 0 the result is placed in the
		W register. If 'd' is 1 the result is placed back in register 'f'.
		If the result is 1, the next instruction, is executed. If the result is 0, then a NOP is
		executed instead making it a 2TCY instruc- tion.
Decode Read Process Write to	Words:	1
register data destination	Cycles:	1(2)
	Q Cycle Activity:	Q1 Q2 Q3 Q4
COMF REG1,0		Decode Read Process Write to
Before Instruction		register 'f' data destination
After Instruction	If Skip:	(2nd Cycle)
REG1 = 0x13		Q1 Q2 Q3 Q4
VV = UXEC		No-Operat No-Operat No-Operati No-Operat tion ion on
Decrement f		ion
[ <i>label</i> ] DECF f,d	Example	HERE DECFSZ CNT, 1
0 < 1 < 107		GOTO LOOP
$d \in [0,1]$		CONTINUE •
		CONTINUE • • •
d ∈ [0,1]		CONTINUE • • • Before Instruction
$d \in [0,1]$ (f) - 1 $\rightarrow$ (destination)		CONTINUE • • •
$d \in [0,1]$ (f) - 1 $\rightarrow$ (destination) Z 00 0011 dfff ffff		CONTINUE • • • Before Instruction PC = address HERE After Instruction CNT = CNT - 1
$d \in [0,1]$ (f) - 1 $\rightarrow$ (destination) Z		CONTINUE • • • Before Instruction PC = address HERE After Instruction CNT = CNT - 1 if CNT = 0, PC = address CONTINUE
$d \in [0,1]$ (f) - 1 $\rightarrow$ (destination) Z 00 0011 dfff ffff		CONTINUE • • • Before Instruction PC = address HERE After Instruction CNT = CNT - 1 if CNT = 0, PC = address CONTINUE if CNT $\neq$ 0,
$\begin{array}{l} d \in [0,1] \\ (f) - 1 \rightarrow (destination) \\ \hline \\ \hline \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ Decrement register 'f'. If 'd' is 0 the result is stored in the W register. If 'd' is 1 the result is stored back in register 'f'. \\ 1 \\ 1 \end{array}$		CONTINUE • • • Before Instruction PC = address HERE After Instruction CNT = CNT - 1 if CNT = 0, PC = address CONTINUE
$d \in [0,1]$ (f) - 1 $\rightarrow$ (destination) Z $00  0011  dfff  ffff$ Decrement register 'f'. If 'd' is 0 the result is stored in the W register. If 'd' is 1 the result is stored back in register 'f'. 1 1 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2		CONTINUE • • • Before Instruction PC = address HERE After Instruction CNT = CNT - 1 if CNT = 0, PC = address CONTINUE if CNT $\neq$ 0,
$\begin{array}{l} d \in [0,1] \\ (f) - 1 \rightarrow (destination) \\ \hline \\ \hline \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ Decrement register 'f'. If 'd' is 0 the result is stored in the W register. If 'd' is 1 the result is stored back in register 'f'. \\ 1 \\ 1 \end{array}$		CONTINUE • • • Before Instruction PC = address HERE After Instruction CNT = CNT - 1 if CNT = 0, PC = address CONTINUE if CNT $\neq$ 0,
$\begin{array}{c c} d \in [0,1] \\ (f) - 1 \rightarrow (destination) \\ \hline Z \\ \hline \hline 00 & 0011 & dfff & ffff \\ \hline \\ Decrement register 'f'. If 'd' is 0 the result is stored in the W register. If 'd' is 1 the result is stored back in register 'f'. \\ 1 \\ 1 \\ \hline Q1 & Q2 & Q3 & Q4 \\ \hline \hline \\ Decode & Read & Process & Write to destination \\ \hline \end{array}$		CONTINUE • • • Before Instruction PC = address HERE After Instruction CNT = CNT - 1 if CNT = 0, PC = address CONTINUE if CNT $\neq$ 0,
$\begin{array}{c c} d \in [0,1] \\ (f) - 1 \rightarrow (destination) \\ \hline Z \\ \hline \hline 00 & 0011 & dfff & ffff \\ \hline Decrement register 'f'. If 'd' is 0 the result is stored in the W register. If 'd' is 1 the result is stored back in register 'f'. \\ 1 \\ 1 \\ \hline Q1 & Q2 & Q3 & Q4 \\ \hline \hline Decode & Read \\ register & 'f' & data & destination \\ \hline \end{array}$		CONTINUE • • • Before Instruction PC = address HERE After Instruction CNT = CNT - 1 if CNT = 0, PC = address CONTINUE if CNT $\neq$ 0,
$\begin{array}{l} d \in [0,1] \\ (f) - 1 \rightarrow (destination) \\ \hline Z \\ \hline \hline 00 & 0011 & dfff & ffff \\ \hline \\ \hline 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff \\ \hline \\ 00 & 001 & dff \\ \hline \\ \hline \\ 00 & 001 & dff \\ \hline \\ 00 & 00 & d$		CONTINUE • • • Before Instruction PC = address HERE After Instruction CNT = CNT - 1 if CNT = 0, PC = address CONTINUE if CNT $\neq$ 0,
$\begin{array}{l} d \in [0,1] \\ (f) - 1 \rightarrow (destination) \\ \hline Z \\ \hline \hline 00 & 0011 & dfff & ffff \\ \hline \\ \hline 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ Decrement register 'f'. If 'd' is 0 the result is stored in the W register. If 'd' is 1 the result is stored back in register 'f'. 1 \\ 1 \\ 1 \\ \hline \\ Q1 & Q2 & Q3 & Q4 \\ \hline \hline \\ Decode & Read & Process & Write to destination \\ \hline \\ DECF & CNT, 1 \\ \hline \\ Before Instruction \\ \hline \\ CNT & = & 0x01 \\ \hline \\ Z & = & 0 \\ \hline \end{array}$		CONTINUE • • • Before Instruction PC = address HERE After Instruction CNT = CNT - 1 if CNT = 0, PC = address CONTINUE if CNT $\neq$ 0,
$\begin{array}{l} d \in [0,1] \\ (f) - 1 \rightarrow (destination) \\ \hline Z \\ \hline \hline 00 & 0011 & dfff & ffff \\ \hline \\ \hline 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff & ffff \\ \hline \\ \hline \\ 00 & 0011 & dfff \\ \hline \\ 00 & 001 & dff \\ \hline \\ \hline \\ 00 & 001 & dff \\ \hline \\ 00 & 00 & d$		CONTINUE • • • Before Instruction PC = address HERE After Instruction CNT = CNT - 1 if CNT = 0, PC = address CONTINUE if CNT $\neq$ 0,
-	$ \begin{bmatrix} label \end{bmatrix} COMF  f,d \\ 0 \leq f \leq 127 \\ d \in [0,1] \\ (\tilde{f}) \rightarrow (destination) \\ Z \\ \hline 00  1001  dfff  ffff \\ \hline The contents of register 'f' are complemented. If 'd' is 0 the result is stored in W. If 'd' is 1 the result is stored back in register 'f'. \\ 1 \\ 1 \\ Q1  Q2  Q3  Q4 \\ \hline Decode  Read \\ register \\ 'f'  register \\ data  destination \\ \hline COMF  REG1, 0 \\ \hline Before Instruction \\ REG1 = 0x13 \\ W = 0xEC \\ \hline Decoment f \\ \hline \end{tabular} $	[ label ] COMF f,dSyntax: $0 \le f \le 127$ $d \in [0,1]$ Operands: $(f) \rightarrow$ (destination)Operation:Z $00  1001  dfff  ffff$ The contents of register 'f' are complemented. If 'd' is 0 the result is stored in W. If 'd' is 1 the result is stored back in register 'f'.Status Affected: Encoding: Description:11Q1 Q2 Q3 Q4DecodeRead register'' dataWrite to destination11COMFREG1,0Before Instruction REG1 = 0x13 W = 0xECIf Skip:Decrement f[label]DECF f,dLabel JDECF f,dExample

GOTO	Unconditional Branch	INCF	Increment f
Syntax:	[ <i>label</i> ] GOTO k	Syntax:	[label] INCF f,d
Operands:	$0 \le k \le 2047$	Operands:	0 ≤ f ≤ 127 d ∈ [0,1]
Operation:	$k \rightarrow PC < 10:0>$ PCLATH<4:3> $\rightarrow$ PC<12:11>	Operation:	(f) + 1 $\rightarrow$ (destination)
Status Affected:	None	Status Affected:	Z
Encoding:	10 1kkk kkkk kkkk	Encoding:	00 1010 dfff ffff
Description:	GOTO is an unconditional branch. The eleven bit immediate value is loaded into PC bits <10:0>. The upper bits of PC are loaded from PCLATH<4:3>. GOTO is a two cycle instruction.	Description:	The contents of register 'f' are incre- mented. If 'd' is 0 the result is placed in the W register. If 'd' is 1 the result is placed back in register 'f'.
Words:	1	Words:	1
Cycles:	2	Cycles:	1
Q Cycle Activity:	Q1 Q2 Q3 Q4	Q Cycle Activity:	Q1 Q2 Q3 Q4
1st Cycle	Decode Read Process Write to literal 'k' data PC		Decode Read register data Vite to destination
2nd Cycle	No-Operat No-Operat No-Operat ion ion	Example	INCF CNT, 1
Example	GOTO THERE		Before Instruction CNT = 0xFF Z = 0
	After Instruction PC = Address THERE		After Instruction CNT = 0x00 Z = 1

INCFSZ	Increment f, Skip if 0	IORLW	Inclusive OR Literal with W
Syntax:	[label] INCFSZ f,d	Syntax:	[ <i>label</i> ] IORLW k
Operands:	$0 \le f \le 127$	Operands:	$0 \le k \le 255$
	d ∈ [0,1]	Operation:	(W) .OR. $k \rightarrow$ (W)
Operation:	(f) + 1 $\rightarrow$ (destination),	Status Affected:	Z
Status Affected	skip if result = 0	Encoding:	11 1000 kkkk kkkk
Status Affected:	None	Description:	The contents of the W register is
Encoding:			OR'ed with the eight bit literal 'k'. The result is placed in the W register.
Description:	The contents of register 'f' are incre- mented. If 'd' is 0 the result is placed in	Words:	1
	the W register. If 'd' is 1 the result is placed back in register 'f'.	Cycles:	1
	If the result is 1, the next instruction is executed. If the result is 0, a NOP is exe- cuted instead making it a 2Tcy instruc-	Q Cycle Activity:	' Q1 Q2 Q3 Q4
	cuted instead making it a 21 CY instruc- tion.	Q Oycle Activity.	Decode Read Process Write to
Words:	1		literal 'k' data W
Cycles:	1(2)	Evenale	TODIN 0-25
Q Cycle Activity:	Q1 Q2 Q3 Q4	Example	IORLW 0x35 Before Instruction
	Decode Read Process Write to		W = 0x9A
	register 'f' data destination		After Instruction
If Skip:	(2nd Cycle)		W = 0xBF Z = 1
	Q1 Q2 Q3 Q4		
	No-Operation No-		
Example	HERE INCFSZ CNT, 1 GOTO LOOP CONTINUE •		
	Before Instruction PC = address HERE After Instruction CNT = CNT + 1 if CNT= 0, PC = address CONTINUE if CNT $\neq$ 0, PC = address HERE +1		

IORWF	Inclusive	e OR W v	with f	
Syntax:	[ label ]	IORWF	f,d	
Operands:	$\begin{array}{l} 0 \leq f \leq 12 \\ d \in \ [0,1] \end{array}$	27		
Operation:	(W) .OR.	$(f) \rightarrow (de)$	estination	i)
Status Affected:	Z			
Encoding:	00	0100	dfff	ffff
Description:	Inclusive C ter 'f'. If 'd' W register back in reg	is 0 the re . If 'd' is 1	sult is place	ced in the
Words:	1			
Cycles:	1			
Q Cycle Activity:	Q1	Q2	Q3	Q4
	Decode	Read register 'f'	Process data	Write to destination
Example	IORWF		RESULT,	0
	Before In			
		RESULT W	= 0x13 = 0x91	-
		ruction RESULT W Z	= 0x13 = 0x93 = 1	-

MOVLW	Move Lit	eral to V	v	
Syntax:	[ label ]	MOVLW	/ k	
Operands:	$0 \le k \le 25$	55		
Operation:	$k \to (W)$			
Status Affected:	None			
Encoding:	11	00xx	kkkk	kkkk
Description:	The eight l register. Th as 0's.			
Words:	1			
Cycles:	1			
Q Cycle Activity:	Q1	Q2	Q3	Q4
	Decode	Read literal 'k'	Process data	Write to W
Example	MOVLW After Inst	<sup>0x5A</sup> ruction W =	0x5A	

MOVF	Move f					
Syntax:	[ label ]	[label] MOVF f,d				
Operands:	$\begin{array}{l} 0 \leq f \leq 12 \\ d \in \left[0,1\right] \end{array}$	.7				
Operation:	$(f) \rightarrow (de)$	stination	)			
Status Affected:	Z					
Encoding:	00	1000	dfff	ffff		
Description:	The contents of register f is moved to a destination dependant upon the status of d. If $d = 0$ , destination is W register. If $d = 1$ , the destination is file register f itself. $d = 1$ is useful to test a file register since status flag Z is affected.					
Words:	1					
Cycles:	1					
Q Cycle Activity:	Q1	Q2	Q3	Q4		
	Decode	Read register 'f'	Process data	Write to destination		
Example		ruction	0 le in FSR I	register		

MOVWF	Move W	to f		
Syntax:	[ label ]	MOVW	= f	
Operands:	$0 \le f \le 12$	27		
Operation:	$(W) \rightarrow (f)$			
Status Affected:	None			
Encoding:	00	0000	lfff	ffff
Description:	Move data 'f'.	from W r	egister to	register
Words:	1			
Cycles:	1			
Q Cycle Activity:	Q1	Q2	Q3	Q4
, , , , , , , , , , , , , , , , , , ,	Decode	Read register 'f'	Process data	Write register 'f'
Example	Decode	register 'f'		
	MOVWF Before In After Inst	OPTIC OPTIC struction OPTION W	data DN_REG = 0xFF = 0x4F	register 'f'

NOP	No Oper	ation		
Syntax:	[ label ]	NOP		
Operands:	None			
Operation:	No opera	ition		
Status Affected:	None			
Encoding:	00	0000	0xx0	0000
Description:	No operat	ion.		
Words:	1			
Cycles:	1			
Q Cycle Activity:	Q1	Q2	Q3	Q4
	Decode	No-Opera tion	No-Opera tion	No-Operat ion
Example	NOP			

RETFIE	Return fi	rom Inter	rupt	
Syntax:	[ label ]	RETFIE		
Operands:	None			
Operation:	$\begin{array}{l} TOS \rightarrow F \\ 1 \rightarrow GIE \end{array}$	PC,		
Status Affected:	None			
Encoding:	00	0000	0000	1001
Description:	Return from Interrupt. Stack is POPed and Top of Stack (TOS) is loaded in the PC. Interrupts are enabled by setting Global Interrupt Enable bit, GIE (INTCON<7>). This is a two cycle instruction.			
Words:	1			
Cycles:	2			
Q Cycle Activity:	Q1	Q2	Q3	Q4
1st Cycle	Decode	No-Opera tion	Set the GIE bit	Pop from the Stack
2nd Cycle	No-Operat ion	No-Opera tion	No-Opera tion	No-Operat ion
Example	RETFIE			

After Interrupt PC = TOS GIE = 1

OPTION	Load Option Register				
Syntax:	[ label ]	OPTION	٧		
Operands:	None				
Operation:	$(W)\toO$	PTION			
Status Affected:	None				
Encoding:	00	0000	0110	0010	
Description:	The conte loaded in t instruction patibility w Since OPT register, th it.	he ΟΡΤΙΟ is suppo ith ΡΙC16 ΓΙΟΝ is a	DN register rted for coo C5X produ readable/v	r. This de com- ucts. vritable	
Words:	1				
Cycles:	1				
Example					
	To maintain upward compatibility with future PIC16CXX products, do not use this instruction.				

RETLW	Return w	vith Liter	al in W		RETURN	Return f	rom Sub	routine	
Syntax:	[ label ]	RETLW	k		Syntax:	[ label ]	RETUR	N	
Operands:	$0 \le k \le 2$	55			Operands:	None			
Operation:	$k \rightarrow (W);$				Operation:	$TOS\toF$	с		
	$TOS \rightarrow F$	°C			Status Affected:	None			
Status Affected:	None				Encoding:	00	0000	0000	1000
Encoding:	11	01xx	kkkk	kkkk	Description:	Return fro	m subrout	ine. The st	ack is
Description:	The W reg bit literal 'k loaded fro	d. The prog	gram coun	ter is		POPed an is loaded i is a two cy	nto the pro	gram cour	```
	return add		s is a two c	ycle	Words:	1			
Words:	1				Cycles:	2			
Cycles:	2				Q Cycle Activity:	Q1	Q2	Q3	Q4
Q Cycle Activity:	2 Q1	Q2	Q3	Q4	1st Cycle	Decode	No-Opera tion	No-Opera tion	Pop from the Stack
1st Cycle	Decode	Read literal 'k'	No-Opera tion	Write to W, Pop from the Stack	2nd Cycle	No-Operat ion	No-Opera tion	No-Opera tion	No-Opera tion
2nd Cycle	No-Operat ion	No-Opera tion	No-Opera tion	No-Operat ion	Example	RETURN			
						After Inte			
Example	CALL TABL	;offset	ains tabl value has table				PC =	TOS	
TABLE	ADDWF PC RETLW k1 RETLW k2	;W = off ;Begin t ;							
	• RETLW kn	; End of	table						
		W =	0x07						
	After Inst		value of k8	3					

RLF	Rotate Left f through Carry	RRF	Rotate Right f through Carry
Syntax:	[label] RLF f,d	Syntax:	[ <i>label</i> ] RRF f,d
Operands:	$\begin{array}{l} 0 \leq f \leq 127 \\ d \in \left[0,1\right] \end{array}$	Operands:	$0 \le f \le 127$ $d \in [0,1]$
Operation:	See description below	Operation:	See description below
Status Affected:	С	Status Affected:	С
Encoding:	00 1101 dfff ffff	Encoding:	00 1100 dfff ffff
Description:	The contents of register 'f' are rotated one bit to the left through the Carry Flag. If 'd' is 0 the result is placed in the W register. If 'd' is 1 the result is stored back in register 'f'.	Description:	The contents of register 'f' are rotated one bit to the right through the Carry Flag. If 'd' is 0 the result is placed in the W register. If 'd' is 1 the result is placed back in register 'f'.
Words:	1	Words:	1
Cycles:	1	Cycles:	1
Q Cycle Activity:	Q1 Q2 Q3 Q4	Q Cycle Activity:	Q1 Q2 Q3 Q4
	Decode Read register 'f' Process Write to data destination		Decode Read register data Vite to destination
Example	RLF REG1,0	Example	RRF REG1,0
	$\begin{array}{rrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrr$		$\begin{array}{rrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrr$

## SLEEP

Syntax:	[ label ]	SLEEP			
Operands:	None				
Operation:	$\begin{array}{l} 00h \rightarrow WDT, \\ 0 \rightarrow WDT \ prescaler, \\ 1 \rightarrow \overline{TO}, \\ 0 \rightarrow \overline{PD} \end{array}$				
Status Affected:	$\overline{TO}, \overline{PD}$				
Encoding:	00	0000	0110	0011	
Description:	The power-down status bit, PD is cleared. Time-out status bit, TO is set. Watchdog Timer and its prescaler are cleared. The processor is put into SLEEP mode with the oscillator stopped. See Section 14.8 for more details.				
Words:	1				
Cycles:	1				
Q Cycle Activity:	Q1	Q2	Q3	Q4	
	Decode	No-Opera tion	No-Opera tion	Go to Sleep	
Example:	SLEEP				

SUBLW	Subtract	W from	Literal			
Syntax:	[ label ]	SUBLV	/ k			
Operands:	$0 \le k \le 25$	$0 \le k \le 255$				
Operation:	k - (W) $\rightarrow$	$k \text{ - } (W) \to (W)$				
Status Affected:	C, DC, Z					
Encoding:	11	110x	kkkk	kkkk		
Description:	ment meth	od) from t	otracted (2's he eight bit in the W reg	literal 'k'.		
Words:	1					
Cycles:	1					
Q Cycle Activity:	Q1	Q2	Q3	Q4		
	Decode	Read literal 'k'	Process data	Write to W		
Example 1:	SUBLW		0x02			
	Before In:	struction				
		W = C = Z =	1 ? ?			
	After Inst	ruction				
		W = C = Z =	1 1; result is 0	positive		
Example 2:	Before In:	struction				
		W = C = Z =	2 ? ?			
	After Inst	ruction				
		W = C = Z =	0 1; result i 1	s zero		
Example 3:	Before In:	struction				
		W = C = Z =	3 ? ?			
	After Inst	ruction				
		W = C = tive	0xFF 0; result is	s nega-		
		Z =	0			

SUBWF	Subtract	W from f		
Syntax:	[ label ]	SUBWF	f,d	
Operands:	$\begin{array}{l} 0 \leq f \leq 12 \\ d \in \left[0,1\right] \end{array}$	7		
Operation:	(f) - (W) -	→ (destina	ition)	
Status Affected:	C, DC, Z			
Encoding:	00	0010	dfff	ffff
Description:	ister from r stored in th	egister 'f'. I ne W regist	nent metho f 'd' is 0 the er. If 'd' is 1 n register 'f	result is the
Words:	1			
Cycles:	1			
Q Cycle Activity:	Q1	Q2	Q3	Q4
	Decode	Read register 'f'	Process data	Write to destination
Example 1:	SUBWF		reg1,1	
	Before Ins	struction		
	REG1 W	=	3 2	
	С	=	?	
	Z	=	?	
	After Instr			
	REG1 W	=	1 2	
	C	=	1; result is	positive
European la Or	Z	=	0	
Example 2:	Before Ins		0	
	REG1 W	=	2 2	
	ç	=	?	
	Z After Instr	=	?	
	REG1		0	
	W	=	2	
	C Z	=	1; result is 1	zero
Example 3:	Z Before Ins		1	
	REG1		1	
	W	=	2	
	C Z	=	? ?	
	After Instr	ruction		
	REG1	=	0xFF	
	W	=	2 O: requit is	nogotivo
	C Z	=	0; result is 0	negative

SWAPF	Swap Ni	bbles in	f			
Syntax:	[ label ]	[label] SWAPF f,d				
Operands:	$0 \le f \le 12$ $d \in [0,1]$	$0 \le f \le 127$ $d \in [0,1]$				
Operation:	· · · ·	ightarrow (destin $ ightarrow$ (destin				
Status Affected:	None					
Encoding:	00	1110	dfff	ffff		
Description:	The upper and lower nibbles of register 'f' are exchanged. If 'd' is 0 the result is placed in W register. If 'd' is 1 the result is placed in register 'f'.					
Words:	1					
Cycles:	1					
Q Cycle Activity:	Q1	Q2	Q3	Q4		
	Decode	Read register 'f'	Process data	Write to destination		
Example	SWAPF	REG,	0			
	Before In	struction				
		REG1	= 0x/	45		
	After Inst	truction				
		REG1 W	= 0x/ = 0x5			

TRIS	Load TRIS Register		
Syntax:	[ <i>label</i> ] TRIS f		
Operands:	$5 \le f \le 7$		
Operation:	(W) $\rightarrow$ TRIS register f;		
Status Affected:	None		
Encoding:	00 0000 0110 0fff		
Description:	The instruction is supported for code compatibility with the PIC16C5X prod- ucts. Since TRIS registers are read- able and writable, the user can directly address them.		
Words:	1		
Cycles:	1		
Example			
	To maintain upward compatibility with future PIC16CXX products, do not use this instruction.		

XORLW	Exclusive OR Literal with W	XORWF	Exclusive OR W with f
Syntax:	[ <i>label</i> ] XORLW k	Syntax:	[ <i>label</i> ] XORWF f,d
Operands:	$0 \le k \le 255$	Operands:	0 ≤ f ≤ 127 d ∈ [0,1]
Operation: Status Affected: Encoding: Description:	(W) .XOR. $k \rightarrow$ (W) Z 11 1010 kkkk kkkk The contents of the W register are XOR'ed with the eight bit literal 'k'.	Operation: Status Affected: Encoding: Description:	(W) .XOR. (f) $\rightarrow$ (destination) Z Exclusive OR the contents of the W
Words:	The result is placed in the W regis- ter. 1	Words:	register with register 'f'. If 'd' is 0 the result is stored in the W register. If 'd' is 1 the result is stored back in register 'f'. 1
Cycles:	1	Cycles:	1
Q Cycle Activity:	Q1Q2Q3Q4DecodeRead literal 'k'Process dataWrite to W	Q Cycle Activity:	Q1         Q2         Q3         Q4           Decode         Read register         Process data         Write to destination
Example:	XORLW 0xAF		'f'
	Before Instruction W = 0xB5	Example	XORWF REG 1 Before Instruction
	After Instruction W = 0x1A		REG = 0xAF $W = 0xB5$ After Instruction $REG = 0x1A$
			W = 0xB5

## 10.0 DEVELOPMENT SUPPORT

## 10.1 Development Tools

The PICmicro<sup>™</sup> microcontrollers are supported with a full range of hardware and software development tools:

- PICMASTER<sup>®</sup>/PICMASTER CE Real-Time In-Circuit Emulator
- ICEPIC Low-Cost PIC16C5X and PIC16CXXX In-Circuit Emulator
- PRO MATE<sup>®</sup> II Universal Programmer
- PICSTART<sup>®</sup> Plus Entry-Level Prototype Programmer
- PICDEM-1 Low-Cost Demonstration Board
- PICDEM-2 Low-Cost Demonstration Board
- PICDEM-3 Low-Cost Demonstration Board
- MPASM Assembler
- MPLAB<sup>™</sup> SIM Software Simulator
- MPLAB-C (C Compiler)
- Fuzzy Logic Development System (*fuzzy*TECH<sup>®</sup>–MP)

## 10.2 <u>PICMASTER: High Performance</u> <u>Universal In-Circuit Emulator with</u> <u>MPLAB IDE</u>

The PICMASTER Universal In-Circuit Emulator is intended to provide the product development engineer with a complete microcontroller design tool set for all microcontrollers in the SX, PIC14C000, PIC16C5X, PIC16CXXX and PIC17CXX families. PICMASTER is supplied with the MPLAB<sup>™</sup> Integrated Development Environment (IDE), which allows editing, "make" and download, and source debugging from a single environment.

Interchangeable target probes allow the system to be easily reconfigured for emulation of different processors. The universal architecture of the PICMASTER allows expansion to support all new Microchip microcontrollers.

The PICMASTER Emulator System has been designed as a real-time emulation system with advanced features that are generally found on more expensive development tools. The PC compatible 386 (and higher) machine platform and Microsoft Windows<sup>®</sup> 3.x environment were chosen to best make these features available to you, the end user.

A CE compliant version of PICMASTER is available for European Union (EU) countries.

## 10.3 <u>ICEPIC: Low-Cost PICmicro™</u> In-Circuit Emulator

ICEPIC is a low-cost in-circuit emulator solution for the Microchip PIC12CXXX, PIC16C5X and PIC16CXXX families of 8-bit OTP microcontrollers.

ICEPIC is designed to operate on PC-compatible machines ranging from 286-AT<sup>®</sup> through Pentium<sup>™</sup> based machines under Windows 3.x environment. ICEPIC features real time, non-intrusive emulation.

## 10.4 PRO MATE II: Universal Programmer

The PRO MATE II Universal Programmer is a full-featured programmer capable of operating in stand-alone mode as well as PC-hosted mode.

The PRO MATE II has programmable VDD and VPP supplies which allows it to verify programmed memory at VDD min and VDD max for maximum reliability. It has an LCD display for displaying error messages, keys to enter commands and a modular detachable socket assembly to support various package types. In standalone mode the PRO MATE II can read, verify or program PIC12CXXX, PIC14C000, PIC16C5X, PIC16CXXX and PIC17CXX devices. It can also set configuration and code-protect bits in this mode.

## 10.5 <u>PICSTART Plus Entry Level</u> <u>Development System</u>

The PICSTART programmer is an easy-to-use, lowcost prototype programmer. It connects to the PC via one of the COM (RS-232) ports. MPLAB Integrated Development Environment software makes using the programmer simple and efficient. PICSTART Plus is not recommended for production programming.

PICSTART Plus supports all PIC12CXXX, PIC14C000, PIC16C5X, PIC16CXXX and PIC17CXX devices with up to 40 pins. Larger pin count devices such as the PIC16C923 and PIC16C924 may be supported with an adapter socket.

## 10.6 <u>PICDEM-1 Low-Cost PICmicro</u> <u>Demonstration Board</u>

The PICDEM-1 is a simple board which demonstrates the capabilities of several of Microchip's microcontrollers. The microcontrollers supported are: PIC16C5X (PIC16C54 to PIC16C58A), PIC16C61, PIC16C62X, PIC16C71, PIC16C8X, PIC17C42, PIC17C43 and PIC17C44. All necessary hardware and software is included to run basic demo programs. The users can program the sample microcontrollers provided with the PICDEM-1 board, on a PRO MATE II or PICSTART-Plus programmer, and easily test firmware. The user can also connect the PICDEM-1 board to the PICMASTER emulator and download the firmware to the emulator for testing. Additional prototype area is available for the user to build some additional hardware and connect it to the microcontroller socket(s). Some of the features include an RS-232 interface, a potentiometer for simulated analog input, push-button switches and eight LEDs connected to PORTB.

## 10.7 <u>PICDEM-2 Low-Cost PIC16CXX</u> Demonstration Board

The PICDEM-2 is a simple demonstration board that supports the PIC16C62, PIC16C64, PIC16C65, PIC16C73 and PIC16C74 microcontrollers. All the necessary hardware and software is included to run the basic demonstration programs. The user can program the sample microcontrollers provided with the PICDEM-2 board, on a PRO MATE II programmer or PICSTART-Plus, and easily test firmware. The PICMASTER emulator may also be used with the PICDEM-2 board to test firmware. Additional prototype area has been provided to the user for adding additional hardware and connecting it to the microcontroller socket(s). Some of the features include a RS-232 interface, push-button switches, a potentiometer for simulated analog input, a Serial EEPROM to demonstrate usage of the I<sup>2</sup>C bus and separate headers for connection to an LCD module and a keypad.

## 10.8 PICDEM-3 Low-Cost PIC16CXXX Demonstration Board

The PICDEM-3 is a simple demonstration board that supports the PIC16C923 and PIC16C924 in the PLCC package. It will also support future 44-pin PLCC microcontrollers with a LCD Module. All the necessary hardware and software is included to run the basic demonstration programs. The user can program the sample microcontrollers provided with the PICDEM-3 board, on a PRO MATE II programmer or PICSTART Plus with an adapter socket, and easily test firmware. The PICMASTER emulator may also be used with the PICDEM-3 board to test firmware. Additional prototype area has been provided to the user for adding hardware and connecting it to the microcontroller socket(s). Some of the features include an RS-232 interface, push-button switches, a potentiometer for simulated analog input, a thermistor and separate headers for connection to an external LCD module and a keypad. Also provided on the PICDEM-3 board is an LCD panel, with 4 commons and 12 segments, that is capable of displaying time, temperature and day of the week. The PICDEM-3 provides an additional RS-232 interface and Windows 3.1 software for showing the demultiplexed LCD signals on a PC. A simple serial interface allows the user to construct a hardware demultiplexer for the LCD signals.

## 10.9 <u>MPLAB™ Integrated Development</u> <u>Environment Software</u>

The MPLAB IDE Software brings an ease of software development previously unseen in the 8-bit microcontroller market. MPLAB is a windows based application which contains:

- A full featured editor
- Three operating modes
  - editor
  - emulator
  - simulator
- A project manager
- Customizable tool bar and key mapping
- · A status bar with project information
- Extensive on-line help

#### MPLAB allows you to:

- Edit your source files (either assembly or 'C')
- One touch assemble (or compile) and download to PICmicro tools (automatically updates all project information)
- Debug using:
- source files
- absolute listing file
- Transfer data dynamically via DDE (soon to be replaced by OLE)
- Run up to four emulators on the same PC

The ability to use MPLAB with Microchip's simulator allows a consistent platform and the ability to easily switch from the low cost simulator to the full featured emulator with minimal retraining due to development tools.

### 10.10 Assembler (MPASM)

The MPASM Universal Macro Assembler is a PChosted symbolic assembler. It supports all microcontroller series including the PIC12C5XX, PIC14000, PIC16C5X, PIC16CXXX, and PIC17CXX families.

MPASM offers full featured Macro capabilities, conditional assembly, and several source and listing formats. It generates various object code formats to support Microchip's development tools as well as third party programmers.

MPASM allows full symbolic debugging from PICMASTER, Microchip's Universal Emulator System.

MPASM has the following features to assist in developing software for specific use applications.

- Provides translation of Assembler source code to object code for all Microchip microcontrollers.
- Macro assembly capability.
- Produces all the files (Object, Listing, Symbol, and special) required for symbolic debug with Microchip's emulator systems.
- Supports Hex (default), Decimal and Octal source and listing formats.

MPASM provides a rich directive language to support programming of the PICmicro. Directives are helpful in making the development of your assemble source code shorter and more maintainable.

## 10.11 Software Simulator (MPLAB-SIM)

The MPLAB-SIM Software Simulator allows code development in a PC host environment. It allows the user to simulate the PICmicro series microcontrollers on an instruction level. On any given instruction, the user may examine or modify any of the data areas or provide external stimulus to any of the pins. The input/ output radix can be set by the user and the execution can be performed in; single step, execute until break, or in a trace mode.

MPLAB-SIM fully supports symbolic debugging using MPLAB-C and MPASM. The Software Simulator offers the low cost flexibility to develop and debug code outside of the laboratory environment making it an excellent multi-project software development tool.

## 10.12 C Compiler (MPLAB-C)

The MPLAB-C Code Development System is a complete 'C' compiler and integrated development environment for Microchip's PICmicro<sup>™</sup> family of microcontrollers. The compiler provides powerful integration capabilities and ease of use not found with other compilers.

For easier source level debugging, the compiler provides symbol information that is compatible with the MPLAB IDE memory display.

## 10.13 <u>Fuzzy Logic Development System</u> (*fuzzy*TECH-MP)

*fuzzy*TECH-MP fuzzy logic development tool is available in two versions - a low cost introductory version, MP Explorer, for designers to gain a comprehensive working knowledge of fuzzy logic system design; and a full-featured version, *fuzzy*TECH-MP, edition for implementing more complex systems.

Both versions include Microchip's *fuzzy*LAB<sup>™</sup> demonstration board for hands-on experience with fuzzy logic systems implementation.

## 10.14 <u>MP-DriveWay™ – Application Code</u> <u>Generator</u>

MP-DriveWay is an easy-to-use Windows-based Application Code Generator. With MP-DriveWay you can visually configure all the peripherals in a PICmicro device and, with a click of the mouse, generate all the initialization and many functional code modules in C language. The output is fully compatible with Microchip's MPLAB-C C compiler. The code produced is highly modular and allows easy integration of your own code. MP-DriveWay is intelligent enough to maintain your code through subsequent code generation.

## 10.15 <u>SEEVAL<sup>®</sup> Evaluation and</u> <u>Programming System</u>

The SEEVAL SEEPROM Designer's Kit supports all Microchip 2-wire and 3-wire Serial EEPROMs. The kit includes everything necessary to read, write, erase or program special features of any Microchip SEEPROM product including Smart Serials<sup>™</sup> and secure serials. The Total Endurance<sup>™</sup> Disk is included to aid in tradeoff analysis and reliability calculations. The total kit can significantly reduce time-to-market and result in an optimized system.

### 10.16 <u>KEELOQ<sup>®</sup> Evaluation and</u> <u>Programming Tools</u>

KEELOQ evaluation and programming tools support Microchips HCS Secure Data Products. The HCS evaluation kit includes an LCD display to show changing codes, a decoder to decode transmissions, and a programming interface to program test transmitters.

## TABLE 10-1: DEVELOPMENT TOOLS FROM MICROCHIP

		PIC12C5XX	PIC14000	PIC16C5X	PIC16CXXX	PIC16C6X	PIC16C7XX	PIC16C8X	PIC16C9XX	PIC17C4X	PIC17C75X	24CXX 25CXX 93CXX	HCS200 HCS300 HCS301
stoubor	PICMASTER®/ PICMASTER-CE In-Circuit Emulator	2	7	7	7	7	7	2	>	>	Available 3Q97		
Emulator P	ICEPIC Low-Cost In-Circuit Emulator	7		7	7	7	7	7					
	MPLAB™ Integrated Development Environment	>	7	7	7	7	7	7	>	7	7		
slo	MPLAB <sup>TM</sup> C Compiler	7	7	7	7	7	7	7	7	7	7		
oT əıswitoö	<i>fuzzy</i> TECH®-MP Explorer/Edition Fuzzy Logic Dev. Tool	7	7	7	7	7	7	7	7	7			
\$	MP-DriveWay™ Applications Code Generator			7	7	7	7	2		7			
	Total Endurance™ Software Model											7	
	PICSTART <sup>®</sup> Lite Ultra Low-Cost Dev. Kit			7		7	7	2					
ສາມາຍເຮ		>	7	7	7	7	7	2	7	7	7		
Progr	PRO MATE <sup>®</sup> II Universal Programmer	7	7	7	7	7	7	7	7	7	7	7	7
	KEELOQ <sup>®</sup> Programmer												7
	SEEVAL <sup>®</sup> Designers Kit											7	
spie	PICDEM-1			7	7			7		7			
o Bo	PICDEM-2					7	7						
Dem	PICDEM-3								7				
	KEELOQ <sup>®</sup> Evaluation Kit												7

## 11.0 ELECTRICAL CHARACTERISTICS FOR PIC16C84

## Absolute Maximum Ratings †

Ambient temperature under bias	55°C to +125°C	
Storage temperature	65°C to +150°C	
Voltage on VDD with respect to Vss	0.3 to +7.5V	
Voltage on MCLR with respect to Vss <sup>(2)</sup>		
Voltage on all other pins with respect to Vss	0.6V to (VDD + 0.6V)	
Total power dissipation <sup>(1)</sup>		
Maximum current out of Vss pin	150 mA	
Maximum current into Vod pin	100 mA	
Input clamp current, Iк (Vi < 0 or Vi > VDD)	± 20 mA	
Output clamp current, loк (Vo < 0 or Vo >VDD)	±20 mA	
Maximum output current sunk by any I/O pin	25 mA	
Maximum output current sourced by any I/O pin	20 mA	
Maximum current sunk by PORTA	80 mA	
Maximum current sourced by PORTA	50 mA	
Maximum current sunk by PORTB	150 mA	
Maximum current sourced by PORTB	100 mA	
<b>Note 1:</b> Power dissipation is calculated as follows: Pdis = VDD x {IDD - $\Sigma$ IOH} + $\Sigma$ {(VDD-VOH) x IOH} + $\Sigma$ (VOI x IOL)		
Note 2: Voltage spikes below Vss at the MCLR pin, inducing currents greater than 80 mA, ma	ay cause latch-up. Thus,	

Note 2: Voltage spikes below Vss at the MCLR pin, inducing currents greater than 80 mA, may cause latch-up. Thus, a series resistor of 50-100Ω should be used when applying a "low" level to the MCLR pin rather than pulling this pin directly to Vss.

† NOTICE: Stresses above those listed under "Absolute Maximum Ratings" may cause permanent damage to the device. This is a stress rating only and functional operation of the device at those or any other conditions above those indicated in the operation listings of this specification is not implied. Exposure to maximum rating conditions for extended periods may affect device reliability.

## TABLE 11-1CROSS REFERENCE OF DEVICE SPECS FOR OSCILLATOR CONFIGURATIONS<br/>AND FREQUENCIES OF OPERATION (COMMERCIAL DEVICES)

OSC	PIC16C84-04	PIC16C84-10	PIC16LC84-04
RC	VDD: 4.0V to 6.0V IDD: 4.5 mA max. at 5.5V IPD: 100 μA max. at 4.0V WDT dis Freq: 4.0 MHz max.	VDD: 4.5V to 5.5V IDD: 1.8 mA typ. at 5.5V IPD: 40.0 μA typ. at 4.5V WDT dis Freq: 4.0 MHz max.	VDD: 2.0V to 6.0V IDD: 4.5 mA max. at 5.5V IPD: 100 μA max. at 4V WDT dis Freq: 2.0 MHz max.
ХТ	VDD:         4.0V to 6.0V           IDD:         4.5 mA max. at 5.5V           IPD:         100 μA max. at 4.0V WDT dis           Freq:         4.0 MHz max.	VDD: 4.5V to 5.5V IDD: 1.8 mA typ. at 5.5V IPD: 40.0 μA typ. at 4.5V WDT dis Freq: 4.0 MHz max.	VDD: 2.0V to 6.0V IDD: 4.5 mA max. at 5.5V IPD: 100 μA max. at 4V WDT dis Freq: 2.0 MHz max.
HS	VDD:         4.5V to 5.5V           IDD:         4.5 mA typ. at 5.5V           IPD:         40.0 μA typ. at 4.5V WDT dis           Freq:         4.0 MHz max.	VDD: 4.5V to 5.5V IDD: 10 mA max. at 5.5V typ. IPD: 40.0 μA typ. at 4.5V WDT dis Freq: 10 MHz max.	Do not use in HS mode
LP	VDD:         4.0V to 6.0V           IDD:         60 μA typ. at 32 kHz, 2.0V           IPD:         26 μA typ. at 2.0V WDT dis           Freq:         200 kHz max.	Do not use in LP mode	VDD: 2.0V to 6.0V IDD: 400 μA max. at 32 kHz, 2.0V IPD: 100 μA max. at 4.0V WDT dis Freq: 200 kHz max.

The shaded sections indicate oscillator selections which are tested for functionality, but not for MIN/MAX specifications. It is recommended that the user select the device type that ensures the specifications required.

#### 11.1 DC CHARACTERISTICS: PIC16C84-04 (Commercial, Industrial) PIC16C84-10 (Commercial, Industrial)

DC Chara Power Su						ure 0°	ditions (unless otherwise stated) $C \leq TA \leq +70^{\circ}C$ (commercial)
Parame- ter No.	Sym	Characteristic	Min	Тур†	Max	-40° Units	°C ≤ TA ≤ +85°C (industrial) Conditions
D001 D001A	Vdd	Supply Voltage	4.0 4.5	_	6.0 5.5	V V	XT, RC and LP osc configuration HS osc configuration
D002	Vdr	RAM Data Retention Voltage <sup>(1)</sup>	1.5*	—	_	V	Device in SLEEP mode
D003	VPOR	VDD start voltage to ensure internal Power-on Reset signal		Vss		V	See section on Power-on Reset for details
D004	SVDD	VDD rise rate to ensure internal Power-on Reset signal	0.05*		—	V/ms	See section on Power-on Reset for details
D010 D010A	IDD	Supply Current <sup>(2)</sup>		1.8 7.3	4.5 10	mA mA	RC and XT osc configuration <sup>(4)</sup> Fosc = 4 MHz, VDD = 5.5V Fosc = 4 MHz, VDD = 5.5V (During EEPROM programming) HS osc configuration (PIC16C84-10)
D013				5.0	10	mA	Fosc = 10  MHz,  VDD = 5.5  V
D020 D021 D021A	IPD	Power-down Current <sup>(3)</sup>		40 38 38	100 100 100	μΑ μΑ μΑ	VDD = 4.0V, WDT enabled, industrial VDD = 4.0V, WDT disabled, commercial VDD = 4.0V, WDT disabled, industrial

These parameters are characterized but not tested.

Data in "Typ" column is at 5.0V, 25°C unless otherwise stated. These parameters are for design guidance only † and are not tested.

Note 1: This is the limit to which VDD can be lowered in SLEEP mode without losing RAM data.

2: The supply current is mainly a function of the operating voltage and frequency. Other factors such as I/O pin loading and switching rate, oscillator type, internal code execution pattern, and temperature also have an impact on current consumption. The test conditions for all IDD measurements in active operation mode are:

OSC1 = external square wave, from rail to rail; all I/O pins tristated, pulled to VDD, T0CKI = VDD,  $\overline{MCLR} = VDD$ ; WDT enabled/disabled as specified.

3: The power down current in SLEEP mode does not depend on the oscillator type. Power-down current is measured with the part in SLEEP mode, with all I/O pins in hi-impedance state and tied to VDD or Vss.

4: For RC osc configuration, current through Rext is not included. The current through the resistor can be estimated by the formula IR = VDD/2Rext (mA) with Rext in kOhm.

## 11.2 DC CHARACTERISTICS PIC16LC84-04 (Commercial, Industrial)

DC Charac Power Sup		Standa Operat	-		ure 0	itions (unless otherwise stated) $^{\circ}C \leq TA \leq +70^{\circ}C$ (commercial) $^{\circ}C \leq TA \leq +85^{\circ}C$ (industrial)	
Parameter No.	Sym	Characteristic	Min	Тур†	Max	Units	Conditions
D001	Vdd	Supply Voltage	2.0	_	6.0	V	XT, RC, and LP osc configuration
D002	Vdr	RAM Data Retention Voltage <sup>(1)</sup>	1.5 *		_	V	Device in SLEEP mode
D003	VPOR	VDD start voltage to ensure internal Power-on Reset signal		Vss	_	V	See section on Power-on Reset for details
D004	Svdd	VDD rise rate to ensure internal Power-on Reset signal	0.05*	_	_	V/ms	See section on Power-on Reset for details
	IDD	Supply Current <sup>(2)</sup>					RC and XT osc configuration <sup>(4)</sup>
D010			—	1	4	mA	FOSC = 2 MHz, VDD = 5.5V
D010A				7.3	10	mA	Fosc = 2 MHz, VDD = 5.5V (During EEPROM programming) LP osc configuration
D014			_	60	400	μA	Fosc = $32 \text{ kHz}$ , VDD = 2.0V, WDT disabled
D020	IPD	Power-down Current <sup>(3)</sup>	—	26	100	μA	VDD = 2.0V, WDT enabled, industrial
D021			—	26	100	μA	VDD = 2.0V, WDT disabled, commercial
D021A			—	26	100	μA	VDD = 2.0V, WDT disabled, industrial

\* These parameters are characterized but not tested.

† Data in "Typ" column is at 5.0V, 25°C unless otherwise stated. These parameters are for design guidance only and are not tested.

Note 1: This is the limit to which VDD can be lowered in SLEEP mode without losing RAM data.

2: The supply current is mainly a function of the operating voltage and frequency. Other factors such as I/O pin loading and switching rate, oscillator type, internal code execution pattern, and temperature also have an impact on the current consumption.

The test conditions for all IDD measurements in active operation mode are:

OSC1=external square wave, from rail to rail; all I/O pins tristated, pulled to VDD, T0CKI = VDD,  $\overline{MCLR}$  = VDD; WDT enabled/disabled as specified.

3: The power down current in SLEEP mode does not depend on the oscillator type. Power-down current is measured with the part in SLEEP mode, with all I/O pins in hi-impedance state and tied to VDD or Vss.

4: For RC osc configuration, current through Rext is not included. The current through the resistor can be estimated by the formula IR = VDD/2Rext (mA) with Rext in kOhm.

### 11.3 DC CHARACTERISTICS: PIC16C84-04 (Commercial, Industrial) PIC16C84-10 (Commercial, Industrial) PIC16LC84-04 (Commercial, Industrial)

			Standard	One	rating Co	onditio	ns (unless otherwise stated)		
DC Chara	cteristi		Operating				$\leq$ TA $\leq$ +70°C (commercial)		
	All Pins Except						$\leq$ TA $\leq$ +85°C (industrial)		
Power Su	oply Pi	ns	Operating voltage VDD range as described in DC spec						
			Section 1	1-1 ar	nd Sectio	n 11.2.			
Parameter No.	Sym	Characteristic	Min	Typ†	Мах	Units	Conditions		
	•••	Input Low Voltage							
	VIL	I/O ports							
D030	VIL	with TTL buffer	Vss		0.8	V	4.5 ≤ VDD ≤ 5.5V		
D030A			VSS		0.0 0.16Vdd	-	entire range <sup>(4)</sup>		
D030A		with Schmitt Trigger buffer	VSS		0.10VDD 0.2VDD	v	entire range		
D032		MCLR, RA4/T0CKI, OSC1	VSS		0.2VDD	v	entire range		
0032		(RC mode)	V 33		0.2000	v			
D033		OSC1 (XT, HS and LP modes) <sup>(1)</sup>	Vss	_	0.3Vdd	V			
		Input High Voltage							
	VIH	I/O ports							
D040		with TTL buffer	0.36Vdd	_	Vdd	V	$4.5 \leq VDD \leq 5.5V$		
D040A			0.48Vdd				entire range <sup>(4)</sup>		
D041		with Schmitt Trigger buffer	0.45Vdd	_	Vdd		entire range		
D042		MCLR, RA4/T0CKI, OSC1	0.85Vdd	_	Vdd	V	5		
		(RC mode)							
D043		OSC1 (XT, HS and LP modes)(1)	0.7Vdd	_	Vdd	V			
D070	IPURB	PORTB weak pull-up current	50*	250*	400*	μA	VDD = 5V, VPIN = VSS		
		Input Leakage Current <sup>(2,3)</sup>							
D060	lı∟	I/O ports	_	_	±1	μA	$Vss \leq VPIN \leq VDD$ ,		
						•	Pin at hi-impedance		
D061		MCLR, RA4/T0CKI	—	—	±5	μA	$Vss \leq VPIN \leq VDD$		
D063		OSC1/CLKIN	—	—	±5	μA	Vss $\leq$ VPIN $\leq$ VDD, XT, HS and LP		
							osc configuration		
		Output Low Voltage							
D080	Vol	I/O ports	—	—	0.6		IOL = 8.5 mA, VDD = 4.5V		
D083		OSC2/CLKOUT	—	—	0.6	V	IOL = 1.6 mA, VDD = 4.5V		
		(RC osc configuration)							
		Output High Voltage							
D090	Vон	I/O ports <sup>(3)</sup>	Vdd - 0.7		—	V	IOH = -3.0 mA, VDD = 4.5V		
D093		OSC2/CLKOUT	Vdd - 0.7	—	—	V	IOH = -1.3 mA, VDD = 4.5V		
		(RC osc configuration)							

These parameters are characterized but not tested.

† Data in "Typ" column is at 5.0V, 25°C unless otherwise stated. These parameters are for design guidance only and are not tested.

Note 1: In RC oscillator configuration, the OSC1/CLKIN pin is a Schmitt Trigger input. It is not recommended that the PIC16C84 be driven with external clock in RC mode.

2: The leakage current on the MCLR pin is strongly dependent on the applied voltage level. The specified levels represent normal operating conditions. Higher leakage current may be measured at different input voltages.

3: Negative current is defined as coming out of the pin.

4: The user may use better of the two specs.

#### 11.4 DC CHARACTERISTICS: PIC16C84-04 (Commercial, Industrial) PIC16C84-10 (Commercial, Industrial) PIC16LC84-04 (Commercial, Industrial)

No.Capacitive Loading Specs on Output Pins OSC2/CLKOUT pin15pFIn XT, HS and LP mode external clock is used OSC1.D100Cosc2OSC2/CLKOUT pin15pFIn XT, HS and LP mode external clock is used OSC1.D101CloAll I/O pins and OSC250pFD120EDData EEPROM Memory Endurance1M10ME/W25°C at 5VD121VDRWVDD for read/writeVMIN6.0VVMIN = Minimum operation voltageD122TDEWErase/Write cycle time <sup>(1)</sup> 1020*ms	DC Charao All Pins E Power Sup	xcept		Standard Operating Conditions (unless otherwise stated)Operating temperature $0^{\circ}C \leq TA \leq +70^{\circ}C$ (commercial) $-40^{\circ}C \leq TA \leq +85^{\circ}C$ (industrial)Operating voltage VDD range as described in DC spec Section 1and Section 11.2.						
D100Cosc2on Output Pins OSC2/CLKOUT pin15pFIn XT, HS and LP model external clock is used OSC1.D101CioAll I/O pins and OSC2 (RC mode)50pFD120EDData EEPROM Memory Endurance1M10M MINE/W25°C at 5V 		Sym	Characteristic	Min	Тур†	Мах	Units	Conditions		
D101CIOAll I/O pins and OSC2 (RC mode)——50pFexternal clock is used OSC1.D120EDData EEPROM Memory Endurance1M10M—E/W25°C at 5V VMIND121VDRWVDD for read/writeVMIN—6.0VVMIN = Minimum operation voltageD122TDEWErase/Write cycle time <sup>(1)</sup> —1020*ms			on Output Pins							
Image: matrix (RC mode)Image: matrix (RC mode)Image: matrix (RC mode)D120EDData EEPROM Memory Endurance1M10M—E/W25°C at 5VD121VDRWVDD for read/writeVMIN—6.0VVMIN = Minimum operative voltageD122TDEWErase/Write cycle time(1)—1020*ms	D100	Cosc2	OSC2/CLKOUT pin	_	_	15	pF	In XT, HS and LP modes when external clock is used to drive OSC1.		
D120     ED     Endurance     1M     10M     -     E/W     25°C at 5V       D121     VDRW     VDD for read/write     VMIN     -     6.0     V     VMIN = Minimum operative       D122     TDEW     Erase/Write cycle time <sup>(1)</sup> -     10     20*     ms	D101	Сю		_	_	50	pF			
D121     VDRW     VDD for read/write     VMIN     —     6.0     V     VMIN = Minimum operative       D122     TDEW     Erase/Write cycle time <sup>(1)</sup> —     10     20*     ms			Data EEPROM Memory							
D122 TDEW Erase/Write cycle time <sup>(1)</sup> — 10 20* ms	D120	ED	Endurance	1M	10M	-	E/W	25°C at 5V		
	D121	Vdrw	VDD for read/write	VMIN	_	6.0	V	VMIN = Minimum operating voltage		
	D122	TDEW	Erase/Write cycle time <sup>(1)</sup>	-	10	20*	ms			
			Program EEPROM Memory							
D130 EP Endurance 100 1000 — E/W	D130	Eр	Endurance	100	1000	_	E/W			
D131 VPR VDD for read VMIN — 6.0 V VMIN = Minimum operative voltage	D131	Vpr	VDD for read	VMIN	-	6.0	V	Vмın = Minimum operating voltage		
D132 VPEW VDD for erase/write 4.5 — 5.5 V	D132	VPEW	VDD for erase/write	4.5	_	5.5	V			
D133 TPEW Erase/Write cycle time <sup>(1)</sup> — 10 — ms	D133	TPEW	Erase/Write cycle time <sup>(1)</sup>	-	10	-	ms			

† Data in "Typ" column is at 5.0V, 25°C unless otherwise stated. These parameters are for design guidance only and are not tested.

Note 1: The user should use interrupts or poll the EEIF or WR bits to ensure the write cycle has completed.

## TABLE 11-2 TIMING PARAMETER SYMBOLOGY

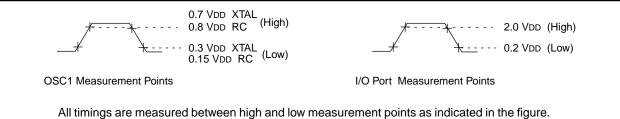
The timing parameter symbols have been created following one of the following formats:

1. TppS2ppS

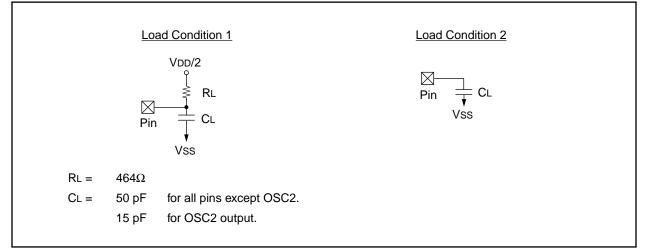
2. TppS

<b>T</b>				
F	Frequency	т	Time	
Lowerc	case symbols (pp) and their meanings:	•		
рр				
2	to	os,osc	OSC1	
ck	CLKOUT	ost	oscillator start-up timer	
су	cycle time	pwrt	power-up timer	
io	I/O port	rbt	RBx pins	
inp	INT pin	tO	TOCKI	
mc	MCLR	wdt	watchdog timer	
Upperc	case symbols and their meanings:			
S				
F	Fall	P	Period	
Н	High	R	Rise	
	Invalid (Hi-impedance)	V	Valid	
L	Low	Z	Hi-impedance	

## FIGURE 11-1: PARAMETER MEASUREMENT INFORMATION



## FIGURE 11-2: LOAD CONDITIONS



## 11.5 Timing Diagrams and Specifications

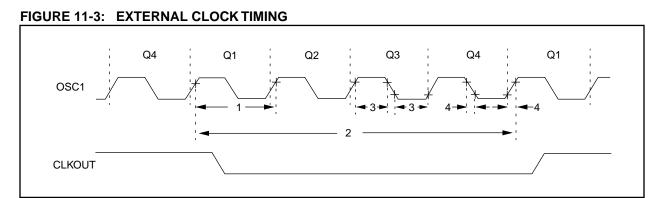


TABLE 11-3	EXTERNAL CLOCK TIMING REQUIREMENTS

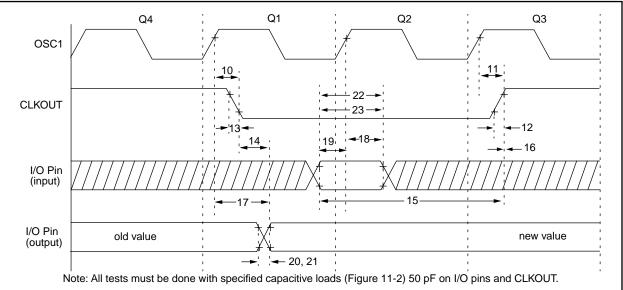
Parameter No.	Sym	Characteristic	Min	Тур†	Max	Units	Conditions	
	Fosc	External CLKIN Frequency <sup>(1)</sup>	DC		2	MHz	XT, RC osc	PIC16LC84-04
			DC	_	4	MHz	XT, RC osc	PIC16C84-04
			DC	_	10	MHz	HS osc	PIC16C84-10
			DC	—	200	kHz	LP osc	PIC16LC84-04
		Oscillator Frequency <sup>(1)</sup>	DC	_	2	MHz	RC osc	PIC16LC84-04
			DC	—	4	MHz	RC osc	PIC16C84-04
			0.1	—	2	MHz	XT osc	PIC16LC84-04
			0.1	—	4	MHz	XT osc	PIC16C84-04
			1	—	10	MHz	HS osc	PIC16C84-10
			DC	—	200	kHz	LP osc	PIC16LC84-04
1	Tosc	External CLKIN Period <sup>(1)</sup>	500	—	-	ns	XT, RC osc	PIC16LC84-04
			250	—	—	ns	XT, RC osc	PIC16C84-04
			100	—	—	ns	HS osc	PIC16C84-10
			5	—	_	μs	LP osc	PIC16LC84-04
		Oscillator Period <sup>(1)</sup>	500	—	_	ns	RC osc	PIC16LC84-04
			250	—	—	ns	RC osc	PIC16C84-04
			500	—	10,000	ns	XT osc	PIC16LC84-04
			250	—	10,000	ns	XT osc	PIC16C84-04
			100	—	1,000	ns	HS osc	PIC16C84-10
			5	—	—	μs	LP osc	PIC16LC84-04
2	Тсү	Instruction Cycle Time <sup>(1)</sup>	0.4	4/Fosc	DC	μs		
3	TosL,	Clock in (OSC1) High or Low	60 *	—	—	ns	XT osc	PIC16LC84-04
	TosH	Time	50 *	—	—	ns	XT osc	PIC16C84-04
			2 *	—	—	μs	LP osc	PIC16LC84-04
			35 *		—	ns	HS osc	PIC16C84-10
4	TosR,	Clock in (OSC1) Rise or Fall Time	25 *	_	-	ns	XT osc	PIC16C84-04
	TosF		50 *	—	—	ns	LP osc	PIC16LC84-04
			15 *	—	_	ns	HS osc	PIC16C84-10

\* These parameters are characterized but not tested.

† Data in "Typ" column is at 5.0V, 25°C unless otherwise stated. These parameters are for design guidance only and are not tested.

Note 1: Instruction cycle period (TCY) equals four times the input oscillator time base period. All specified values are based on characterization data for that particular oscillator type under standard operating conditions with the device executing code. Exceeding these specified limits may result in an unstable oscillator operation and/or higher than expected current consumption. All devices are tested to operate at "min." values with an external clock applied to the OSC1 pin. When an external clock input is used, the "Max." cycle time limit is "DC" (no clock) for all devices.





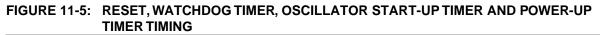
Parameter No.	Sym	Characteristic		Min	Тур†	Max	Units	Conditions
10	TosH2ckL	OSC1 <sup>↑</sup> to CLKOUT↓	PIC16C84		15	30 *	ns	Note 1
10A			PIC16LC84		15	120 *	ns	Note 1
11	TosH2ckH	OSC1 <sup>↑</sup> to CLKOUT <sup>↑</sup>	PIC16C84	_	15	30 *	ns	Note 1
11A			PIC16LC84	_	15	120 *	ns	Note 1
12	TckR	CLKOUT rise time	PIC16C84	—	15	30 *	ns	Note 1
12A			PIC16LC84	_	15	100 *	ns	Note 1
13	TckF	CLKOUT fall time	PIC16C84		15	30 *	ns	Note 1
13A			PIC16LC84	—	15	100 *	ns	Note 1
14	TckL2ioV	CLKOUT ↓ to Port out	valid	_	_	0.5Tcy +20 *	ns	Note 1
15	TioV2ckH	Port in valid before	PIC16C84	0.30Tcy + 30 *	_	_	ns	Note 1
		CLKOUT ↑	PIC16LC84	0.30Tcy + 80 *	_	_	ns	Note 1
16	TckH2iol	Port in hold after CLKOUT 1		0 *	_	_	ns	Note 1
17	TosH2ioV	OSC1 <sup>↑</sup> (Q1 cycle) to	PIC16C84	_	_	125 *	ns	
		Port out valid	PIC16LC84	_	_	250 *	ns	
18	TosH2iol	OSC1↑ (Q2 cycle) to P (I/O in hold time)	Port input invalid	TBD	—	_	ns	
19	TioV2osH	Port input valid to OSC (I/O in setup time)	1↑	TBD	—	_	ns	
20	TioR	Port output rise time	PIC16C84		10	25 *	ns	
20A			PIC16LC84	—	10	60 *	ns	
21	TioF	Port output fall time	PIC16C84	_	10	25 *	ns	
21A			PIC16LC84		10	60 *	ns	
22	Tinp	INT pin high	PIC16C84	20 *	_	—	ns	
22A		or low time	PIC16LC84	55 *	_	_	ns	
23	Trbp	RB7:RB4 change INT	PIC16C84	20 *	_	—	ns	
23A		high or low time	PIC16LC84	55 *	_	_	ns	

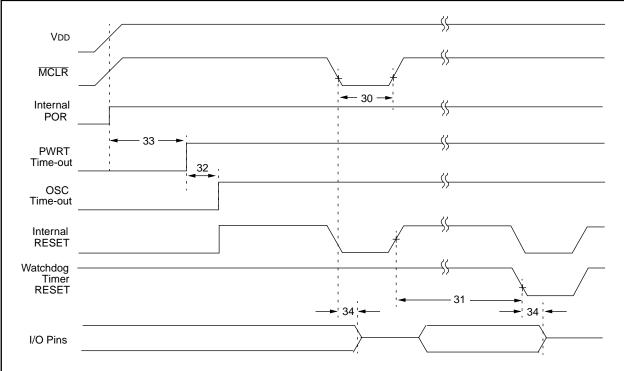
	TABLE 11-4	<b>CLKOUT AND I/O TIMING REQUIREMENTS</b>
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\* These parameters are characterized but not tested.

† Data in "Typ" column is at 5.0V, 25°C unless otherwise stated. These parameters are for design guidance only and are not tested.

Note 1: Measurements are taken in RC Mode where CLKOUT output is 4 x Tosc.





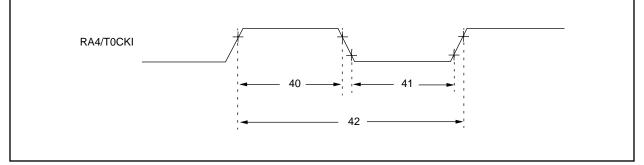
# TABLE 11-5RESET, WATCHDOG TIMER, OSCILLATOR START-UP TIMER AND POWER-UP<br/>TIMER REQUIREMENTS

Parameter No.	Sym	Characteristic	Min	Тур†	Мах	Units	Conditions
30	TmcL	MCLR Pulse Width (low)	350 * 150 *	_	_	ns ns	$\begin{array}{l} 2.0V \leq VDD \leq 3.0V \\ 3.0V \leq VDD \leq 6.0V \end{array}$
31	Twdt	Watchdog Timer Time-out Period (No Prescaler)	7 *	18	33 *	ms	VDD = 5V
32	Tost	Oscillation Start-up Timer Period	_	1024Tosc	—	ms	Tosc = OSC1 period
33	Tpwrt	Power-up Timer Period	28 *	72	132 *	ms	VDD = 5.0V
34	Tıoz	I/O Hi-impedance from MCLR Low or reset	_	—	100 *	ns	

These parameters are characterized but not tested.

† Data in "Typ" column is at 5.0V, 25°C unless otherwise stated. These parameters are for design guidance only and are not tested.

## FIGURE 11-6: TIMER0 CLOCK TIMINGS



## TABLE 11-6 TIMER0 CLOCK REQUIREMENTS

Parameter No.	Sym	Characteristic		Min	Тур†	Max	Units	Conditions
40	Tt0H	T0CKI High Pulse Width	No Prescaler	0.5Tcy + 20 *	—	_	ns	
			With Prescaler	50 * 30 *	_	_	ns ns	$\begin{array}{l} 2.0V \leq V \text{DD} \leq 3.0V \\ 3.0V \leq V \text{DD} \leq 6.0V \end{array}$
41	Tt0L	T0CKI Low Pulse Width	No Prescaler	0.5Tcy + 20 *	—	_	ns	
			With Prescaler	50 * 20 *	_	_	ns ns	$\begin{array}{l} 2.0V \leq V \text{DD} \leq 3.0V \\ 3.0V \leq V \text{DD} \leq 6.0V \end{array}$
42	Tt0P	T0CKI Period		<u>Tcy + 40</u> * N	_	_	ns	N = prescale value (2, 4,, 256)

\* These parameters are characterized but not tested.

† Data in "Typ" column is at 5.0V, 25°C unless otherwise stated. These parameters are for design guidance only and are not tested. NOTES:

# 12.0 DC & AC CHARACTERISTICS GRAPHS/TABLES FOR PIC16C84

The graphs and tables provided in this section are for design guidance and are not tested or guaranteed.

In some graphs or tables, the data presented are **outside specified operating range** (i.e., outside specified VDD range). This is for **information only** and devices are guaranteed to operate properly only within the specified range.

The data presented in this section is a **statistical summary** of data collected on units from different lots over a period of time and matrix samples. 'Typical' represents the mean of the distribution at  $25^{\circ}$ C, while 'max' or 'min' represents (mean +  $3\sigma$ ) and (mean -  $3\sigma$ ) respectively, where  $\sigma$  is standard deviation.

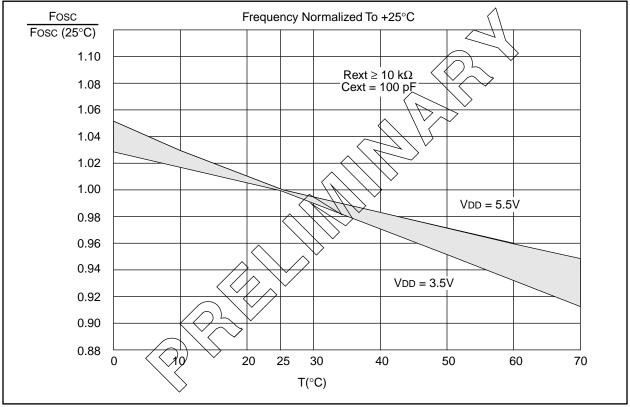


FIGURE 12-1: TYPICAL RC OSCILLATOR FREQUENCY vs. TEMPERATURE

TABLE 12-1	<b>RC OSCILLATOR FREQUENCIES *</b>
------------	------------------------------------

Cext	Rext	Average Fosc @ 5V, 25°C	
20 pF	3.3k	4.68 MHz	± 27%
	5.1k	3.94 MHz	± 25%
	10k	2.34 MHz	± 29%
	100k	250.16 kHz	± 33%
100 pF	3.3k	1.49 MHz	± 25%
	5.1k	1.12 MHz	± 25%
	10k	620.31 kHz	± 30%
	100k	90.25 kHz	± 26%
300 pF	3.3k	524.24 kHz	± 28%
	5.1k	415.52 kHz	± 30%
	10k	270.33 kHz	± 26%
	100k	25.37 kHz	± 25%

\*Measured in PDIP Packages. The percentage variation indicated here is part to part variation due to normal process distribution. The variation indicated is ±3 standard deviation from average value.

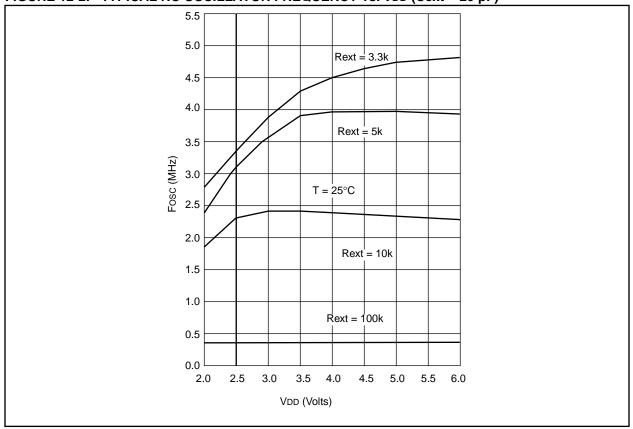


FIGURE 12-2: TYPICAL RC OSCILLATOR FREQUENCY vs. VDD (Cext = 20 pF)

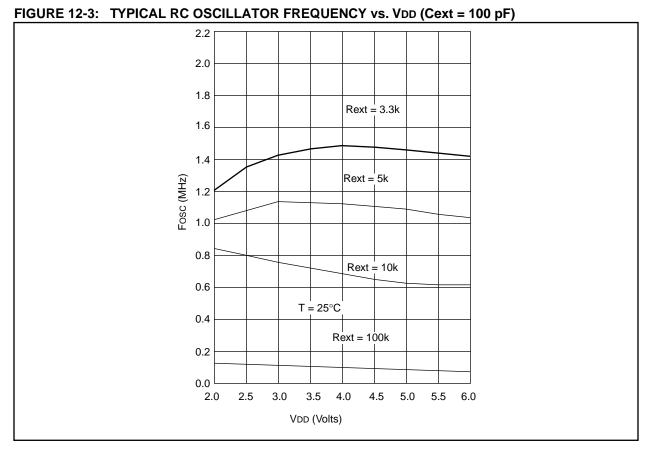
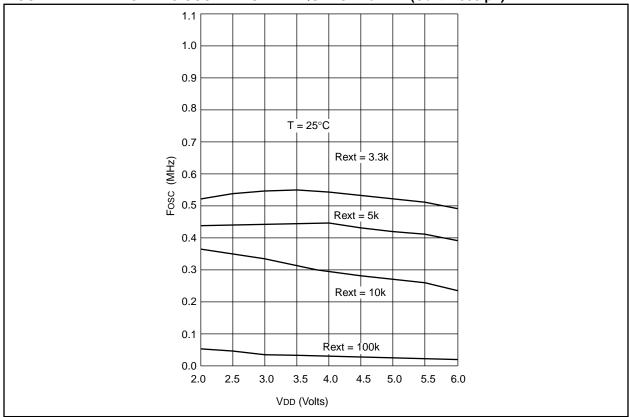


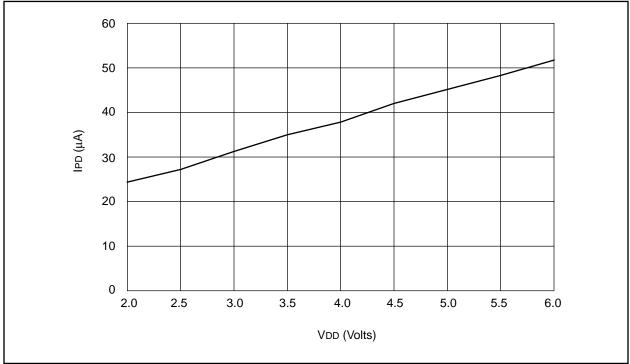
FIGURE 12-4: TYPICAL RC OSCILLATOR FREQUENCY vs. VDD (Cext = 300 pF)



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# **PIC16C84**





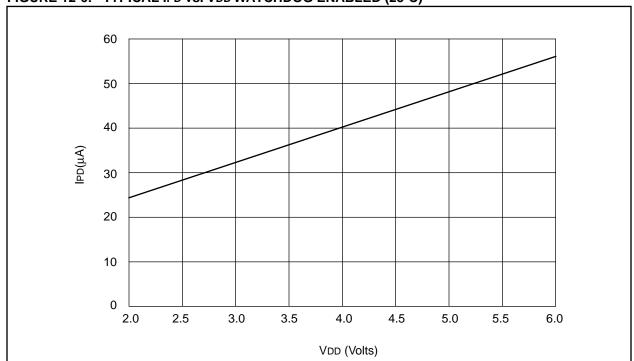
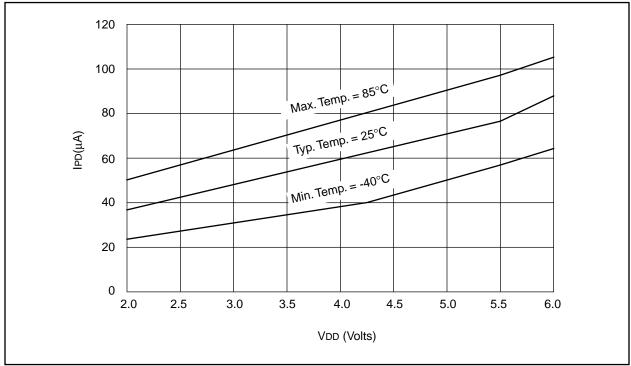
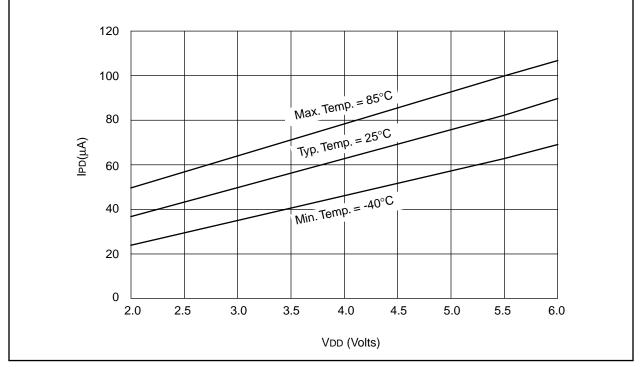




FIGURE 12-7: MAXIMUM IPD vs. VDD WATCHDOG DISABLED







\* IPD, with Watchdog Timer enabled, has two components: The leakage current which increases with higher temperature and the operating current of the Watchdog Timer logic which increases with lower temperature. At -40°C, the latter dominates explaining the apparently anomalous behavior.

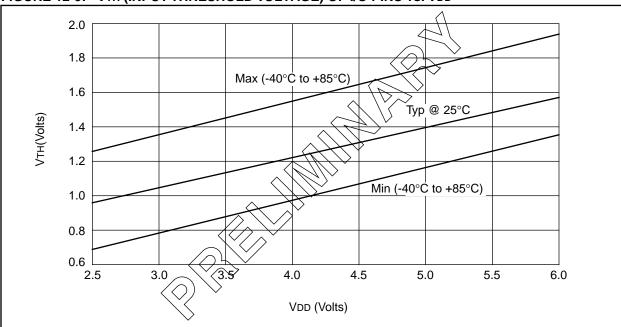
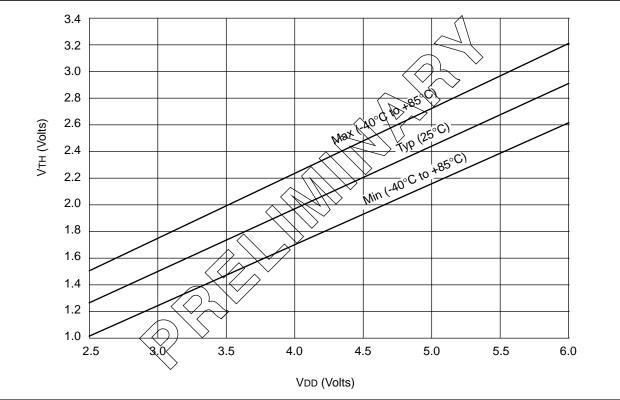


FIGURE 12-9: VTH (INPUT THRESHOLD VOLTAGE) OF I/O PINS vs. VDD





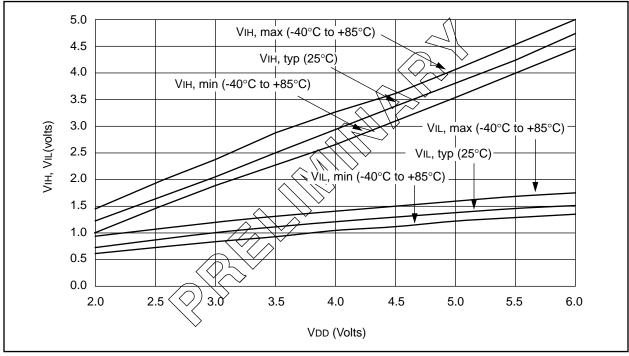
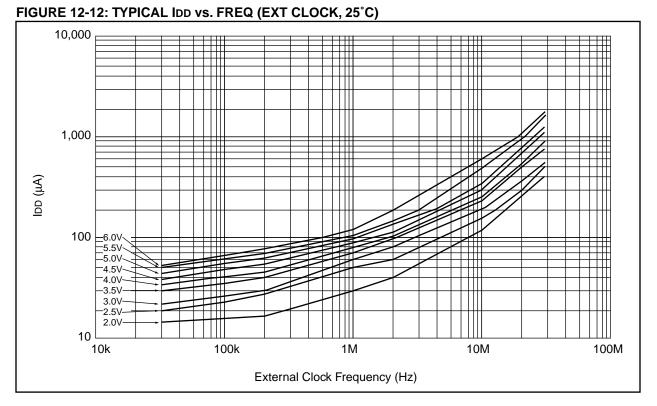


FIGURE 12-11: VIH, VIL OF MCLR, TOCKI and OSC1 (IN RC MODE) vs. VDD



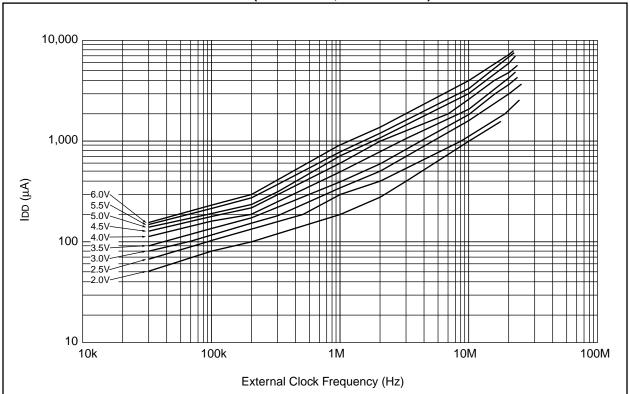
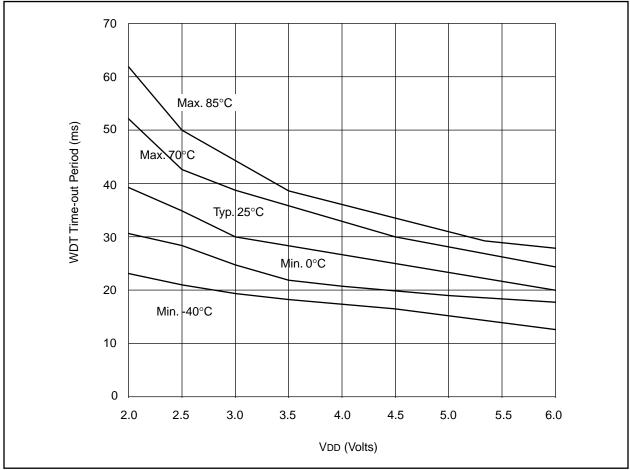


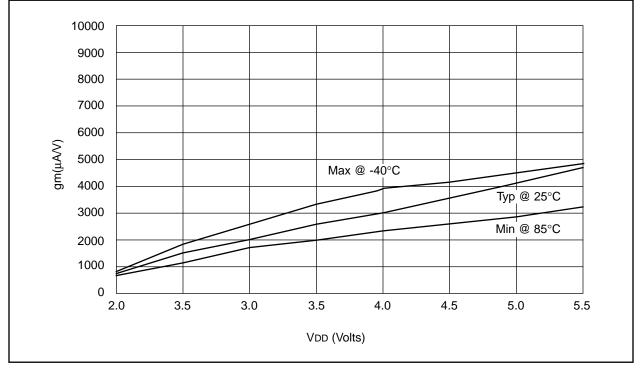
FIGURE 12-13: MAXIMUM IDD vs. FREQ (EXT CLOCK, -40° TO +85°C)

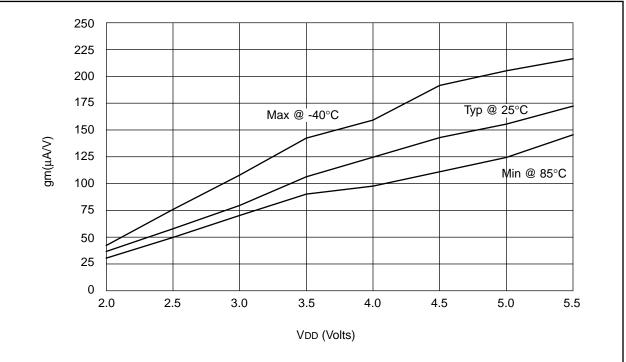
# **PIC16C84**

### FIGURE 12-14: WDT TIME-OUT PERIOD vs. VDD













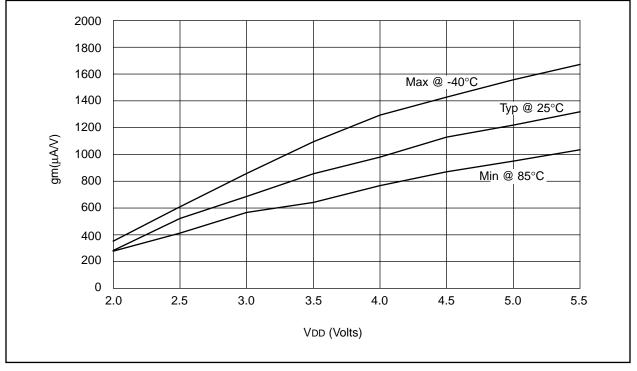


FIGURE 12-18: IOH vs. VOH, VDD = 3V

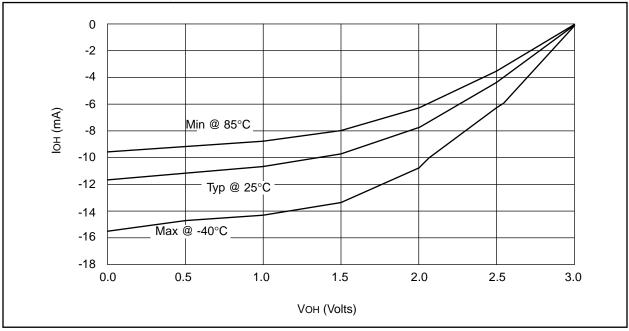
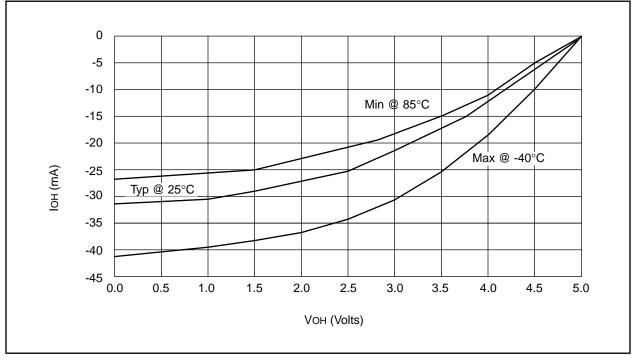


FIGURE 12-19: IOH vs. VOH, VDD = 5V



## FIGURE 12-20: IOL vs. VOL, VDD = 3V

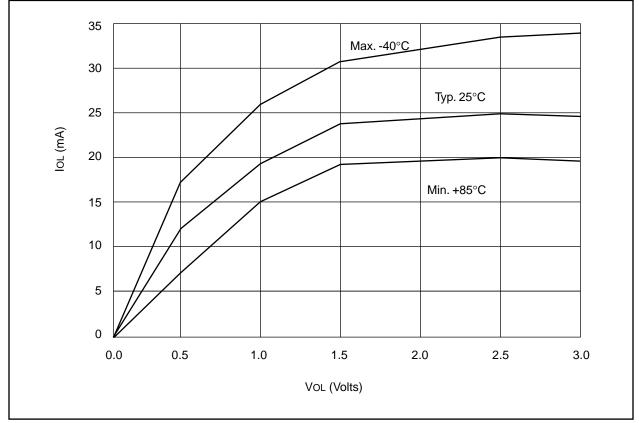
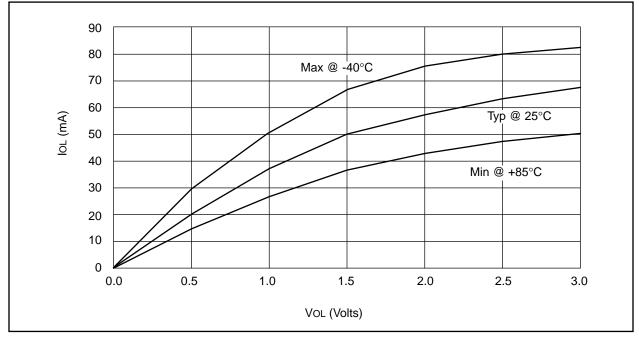
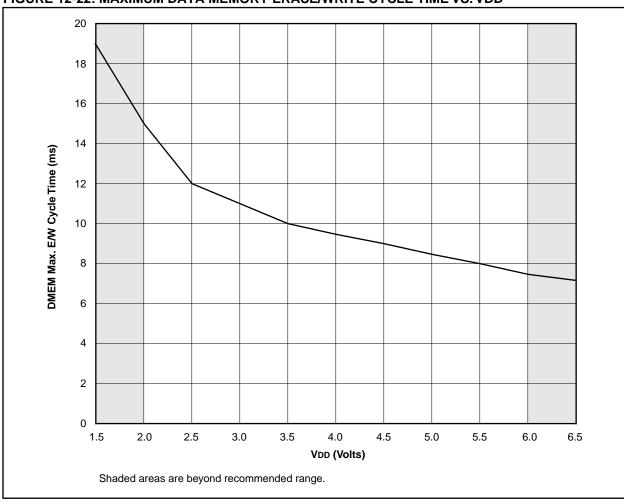


FIGURE 12-21: IOL vs. VOL, VDD = 5V





## FIGURE 12-22: MAXIMUM DATA MEMORY ERASE/WRITE CYCLE TIME VS. VDD

## TABLE 12-2 INPUT CAPACITANCE\*

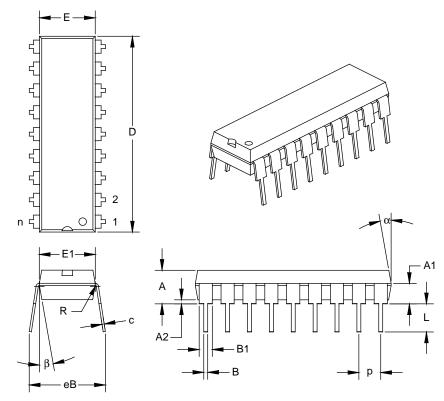
Din Neme	Typical Capacitance (pF)			
Pin Name	18L PDIP	18L SOIC		
PORTA	5.0	4.3		
PORTB	5.0	4.3		
MCLR	17.0	17.0		
OSC1/CLKIN	4.0	3.5		
OSC2/CLKOUT	4.3	3.5		
ТОСКІ	3.2	2.8		

\* All capacitance values are typical at 25°C. A part to part variation of ±25% (three standard deviations) should be taken into account.

NOTES:

# **13.0 PACKAGING INFORMATION**

13.1 K04-007 18-Lead Plastic Dual In-line (P) – 300 mil



Units			INCHES*		Μ	ILLIMETER	S
Dimension Limits		MIN	NOM	MAX	MIN	NOM	MAX
PCB Row Spacing			0.300			7.62	
Number of Pins	n		18			18	
Pitch	р		0.100			2.54	
Lower Lead Width	В	0.013	0.018	0.023	0.33	0.46	0.58
Upper Lead Width	B1 <sup>†</sup>	0.055	0.060	0.065	1.40	1.52	1.65
Shoulder Radius	R	0.000	0.005	0.010	0.00	0.13	0.25
Lead Thickness	С	0.005	0.010	0.015	0.13	0.25	0.38
Top to Seating Plane	A	0.110	0.155	0.155	2.79	3.94	3.94
Top of Lead to Seating Plane	A1	0.075	0.095	0.115	1.91	2.41	2.92
Base to Seating Plane	A2	0.000	0.020	0.020	0.00	0.51	0.51
Tip to Seating Plane	L	0.125	0.130	0.135	3.18	3.30	3.43
Package Length	D‡	0.890	0.895	0.900	22.61	22.73	22.86
Molded Package Width	E‡	0.245	0.255	0.265	6.22	6.48	6.73
Radius to Radius Width	E1	0.230	0.250	0.270	5.84	6.35	6.86
Overall Row Spacing	eB	0.310	0.349	0.387	7.87	8.85	9.83
Mold Draft Angle Top	α	5	10	15	5	10	15
Mold Draft Angle Bottom	β	5	10	15	5	10	15

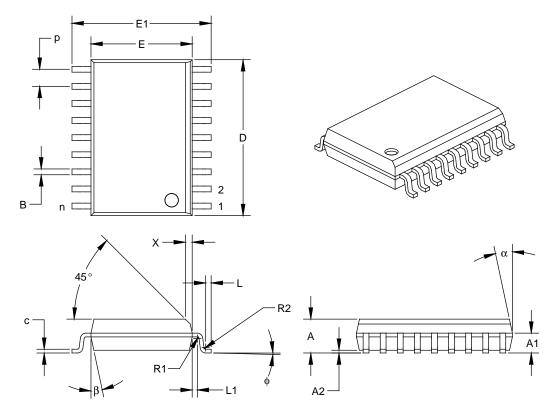
\* Controlling Parameter.

<sup>†</sup> Dimension "B1" does not include dam-bar protrusions. Dam-bar protrusions shall not exceed 0.003" (0.076 mm) per side or 0.006" (0.152 mm) more than dimension "B1."

<sup>‡</sup> Dimensions "D" and "E" do not include mold flash or protrusions. Mold flash or protrusions shall not exceed 0.010" (0.254 mm) per side or 0.020" (0.508 mm) more than dimensions "D" or "E."

# **PIC16C84**

# 13.2 K04-051 18-Lead Plastic Small Outline (SO) – Wide, 300 mil



Units		INCHES*		М	ILLIMETERS	S	
Dimension Limits		MIN	NOM	MAX	MIN	NOM	MAX
Pitch	р		0.050			1.27	
Number of Pins	n		18			18	
Overall Pack. Height	А	0.093	0.099	0.104	2.36	2.50	2.64
Shoulder Height	A1	0.048	0.058	0.068	1.22	1.47	1.73
Standoff	A2	0.004	0.008	0.011	0.10	0.19	0.28
Molded Package Length	D <sup>‡</sup>	0.450	0.456	0.462	11.43	11.58	11.73
Molded Package Width	E‡	0.292	0.296	0.299	7.42	7.51	7.59
Outside Dimension	E1	0.394	0.407	0.419	10.01	10.33	10.64
Chamfer Distance	X	0.010	0.020	0.029	0.25	0.50	0.74
Shoulder Radius	R1	0.005	0.005	0.010	0.13	0.13	0.25
Gull Wing Radius	R2	0.005	0.005	0.010	0.13	0.13	0.25
Foot Length	L	0.011	0.016	0.021	0.28	0.41	0.53
Foot Angle	φ	0	4	8	0	4	8
Radius Centerline	L1	0.010	0.015	0.020	0.25	0.38	0.51
Lead Thickness	С	0.009	0.011	0.012	0.23	0.27	0.30
Lower Lead Width	B <sup>†</sup>	0.014	0.017	0.019	0.36	0.42	0.48
Mold Draft Angle Top	α	0	12	15	0	12	15
Mold Draft Angle Bottom	β	0	12	15	0	12	15

Controlling Parameter.

 Dimension "B" does not include dam-bar protrusions. Dam-bar protrusions shall not exceed 0.003" (0.076 mm) per side or 0.006" (0.152 mm) more than dimension "B."

<sup>‡</sup> Dimensions "D" and "E" do not include mold flash or protrusions. Mold flash or protrusions shall not exceed 0.010" (0.254 mm) per side or 0.020" (0.508 mm) more than dimensions "D" or "E."

# APPENDIX A: FEATURE IMPROVEMENTS -FROM PIC16C5X TO PIC16C84

The following is the list of feature improvements over the PIC16C5X microcontroller family:

- Instruction word length is increased to 14 bits. This allows larger page sizes both in program memory (2K now as opposed to 512 before) and the register file (128 bytes now versus 32 bytes before).
- 2. A PC latch register (PCLATH) is added to handle program memory paging. PA2, PA1 and PA0 bits are removed from the status register and placed in the option register.
- 3. Data memory paging is redefined slightly. The STATUS register is modified.
- 4. Four new instructions have been added: RETURN, RETFIE, ADDLW, and SUBLW. Two instructions, TRIS and OPTION, are being phased out although they are kept for compatibility with PIC16C5X.
- 5. OPTION and TRIS registers are made addressable.
- 6. Interrupt capability is added. Interrupt vector is at 0004h.
- 7. Stack size is increased to 8 deep.
- 8. Reset vector is changed to 0000h.
- Reset of all registers is revisited. Five different reset (and wake-up) types are recognized. Registers are reset differently.
- 10. Wake up from SLEEP through interrupt is added.
- 11. Two separate timers, the Oscillator Start-up Timer (OST) and Power-up Timer (PWRT), are included for more reliable power-up. These timers are invoked selectively to avoid unnecessary delays on power-up and wake-up.
- 12. PORTB has weak pull-ups and interrupt on change features.
- 13. T0CKI pin is also a port pin (RA4/T0CKI).
- 14. FSR is a full 8-bit register.
- 15. "In system programming" is made possible. The user can program PIC16CXX devices using only five pins: VDD, VSS, VPP, RB6 (clock) and RB7 (data in/out).

# APPENDIX B: CODE COMPATIBILITY - FROM PIC16C5X TO PIC16C84

To convert code written for PIC16C5X to PIC16C84, the user should take the following steps:

- 1. Remove any program memory page select operations (PA2, PA1, PA0 bits) for CALL, GOTO.
- 2. Revisit any computed jump operations (write to PC or add to PC, etc.) to make sure page bits are set properly under the new scheme.
- 3. Eliminate any data memory page switching. Redefine data variables for reallocation.
- 4. Verify all writes to STATUS, OPTION, and FSR registers since these have changed.
- 5. Change reset vector to 0000h.

## APPENDIX C: WHAT'S NEW IN THIS DATA SHEET

No new information has been added to this data sheet.

For information on upgrade devices from the PIC16C84, please refer to the PIC16F8X data sheet.

# APPENDIX D: WHAT'S CHANGED IN THIS DATA SHEET

Here's what's changed in this data sheet:

- 1. Some sections have been rearranged for clarity and consistency.
- 2. Errata information has been included.

# APPENDIX E: CONVERSION CONSIDERATIONS - PIC16C84 TO PIC16F83/F84 AND PIC16CR83/CR84

Considerations for converting from the PIC16C84 to the PIC16F84 are listed in the table below. These considerations apply to converting from the PIC16C84 to the PIC16F83 (same as PIC16F84 except for program and data RAM memory sizes) and the PIC16CR84 and PIC16CR83 (ROM versions of Flash devices). Development Systems support is available for all of the PIC16X8X devices.

Difference	PIC16C84	PIC16F84
The polarity of the PWRTE bit has been reversed. Ensure that the pro- grammer has this bit correctly set before programming.	PWRTE	PWRTE
The PIC16F84 (and PIC16CR84) have larger RAM sizes. Ensure that this does not cause an issue with your program.	RAM = 36 bytes	RAM = 68 bytes
The MCLR pin now has an on-chip filter. The input signal on the MCLR pin will require a longer low pulse to generate an interrupt.		$\frac{\text{MCLR}}{\text{Pulse width (low)}}$ = 1000ns; 2.0V $\leq$ VDD $\leq$ 6.0V
Some electrical specifications have been improved (see IPD example). Compare the electrical specifica- tions of the two devices to ensure that this will not cause a compatibil- ity issue.	IPD (typ @ 2V) = 26μA IPD (max @ 4V, WDT disabled) =100μA (PIC16C84) =100μA (PIC16LC84)	IPD (typ @ 2V) < 1μA IPD (max @ 4V, WDT disabled) =14μA (PIC16F84) =7μA (PIC16LF84)
PORTA and crystal oscillator values less than 500kHz	For crystal oscillator configurations operating below 500kHz, the device may generate a spurious internal Q- clock when PORTA<0> switches state.	N/A
RB0/INT pin	TTL	TTL/ST* (* This buffer is a Schmitt Trigger input when configured as the exter- nal interrupt.)
EEADR<7:6> and IDD	It is recommended that the EEADR<7:6> bits be cleared. When either of these bits is set, the maximum IDD for the device is higher than when both are cleared.	N/A
Code Protect	1 CP bit	9 CP bits
Recommended value of REXT for RC oscillator circuits	Rext = 3kΩ - 100kΩ	Rext = 5kΩ - 100kΩ
GIE bit unintentional enable	If an interrupt occurs while the Glo- bal Interrupt Enable (GIE) bit is being cleared, the GIE bit may unin- tentionally be re-enabled by the user's Interrupt Service Routine (the RETFIE instruction).	N/A

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#### ftp.mchip.com/biz/mchip

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#### Internet:

You can telnet or ftp to the Microchip BBS at the address:

#### mchipbbs.microchip.com

#### CompuServe Communications Network:

When using the BBS via the Compuserve Network, in most cases, a local call is your only expense. The Microchip BBS connection does not use CompuServe membership services, therefore you do not need CompuServe membership to join Microchip's BBS. There is no charge for connecting to the Microchip BBS. The procedure to connect will vary slightly from country to country. Please check with your local CompuServe agent for details if you have a problem. CompuServe service allow multiple users various baud rates depending on the local point of access.

The following connect procedure applies in most locations.

- 1. Set your modem to 8-bit, No parity, and One stop (8N1). This is not the normal CompuServe setting which is 7E1.
- 2. Dial your local CompuServe access number.
- 3. Depress the <Enter> key and a garbage string will appear because CompuServe is expecting a 7E1 setting.
- 4. Type +, depress the <Enter> key and "Host Name:" will appear.
- 5. Type MCHIPBBS, depress the <Enter> key and you will be connected to the Microchip BBS.

In the United States, to find the CompuServe phone number closest to you, set your modem to 7E1 and dial (800) 848-4480 for 300-2400 baud or (800) 331-7166 for 9600-14400 baud connection. After the system responds with "Host Name:", type NETWORK, depress the <Enter> key and follow CompuServe's directions.

For voice information (or calling from overseas), you may call (614) 723-1550 for your local CompuServe number.

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The Systems Information and Upgrade Line provides system users a listing of the latest versions of all of Microchip's development systems software products. Plus, this line provides information on how customers can receive any currently available upgrade kits.The Hot Line Numbers are:

1-800-755-2345 for U.S. and most of Canada, and

1-602-786-7302 for the rest of the world.

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# PIC16C84 PRODUCT IDENTIFICATION SYSTEM

To order or obtain information, e.g., on pricing or delivery, refer to the factory or the listed sales office.

PART NO. Device	-XX X /XX XXX Frequency Temperature Package Pattern Range Range	Examples: a) PIC16C84 -04/P 301 = Commercial temp., PDIP package, 4MHz, normal VDD
Device	PIC16C84 <sup>(2)</sup> , PIC16C84T <sup>(3)</sup> PIC16LC84 <sup>(2)</sup> , PIC16LC84T <sup>(3)</sup>	<ul> <li>limitis, QTP pattern #301.</li> <li>b) PIC16LC84 - 04I/SO = Industrial temp., SOIC package, 200kHz, Extended VDD</li> </ul>
Frequency Range	04 = 4 MHz 10 = 10 MHz	limits.
Temperature Range	$b^{(1)} = 0^{\circ}C \text{ to } +70^{\circ}C \text{ (Commercial)}$ I = -40^{\circ}C \text{ to } +85^{\circ}C \text{ (Industrial)}	<ul> <li>Note 1: b = blank</li> <li>2: C = Standard VDD range LC = Extended VDD range</li> <li>3: T = in tape and reel - SOIC, SSOP packages only.</li> </ul>
Package	P = PDIP SO = SOIC (Gull Wing, 300 mil body)	
Pattern	3-digit Pattern Code for QTP, ROM (blank otherwise)	

# SALES AND SUPPORT

#### **Data Sheets**

Products supported by a preliminary Data Sheet may have an errata sheet describing minor operational differences and recommended workarounds. To determine if an errata sheet exists for a particular device, please contact one of the following:

- 1.
- Your local Microchip sales office (see last page) The Microchip Corporate Literature Center U.S. FAX: (602) 786-7277 2.
- The Microchip's Bulletin Board, via your local CompuServe number (CompuServe membership NOT required). 3.
- Please specify which device, revision of silicon and Data Sheet (include Literature #) you are using.

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For the latest version information and upgrade kits for Microchip Development Tools, please call 1-800-755-2345 or 1-602-786-7302. The latest version of Development Tools software can be downloaded from either our Bulletin Board or Worldwide Web Site. (Information on how to connect to our BBS or WWW site can be found in the On-Line Support section of this data sheet.)



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